

FM24C16

16Kb FRAM Serial Memory



Features

16K bit Ferroelectric Nonvolatile RAM

- Organized as 2,048 x 8 bits
- High Endurance 10 Billion (10^{10}) Read/Writes
- 10 Year Data Retention
- NoDelay™ Writes
- Advanced High-Reliability Ferroelectric Process

Fast Two-wire Serial Interface

- Up to 400 kHz maximum bus frequency
- Direct hardware replacement for EEPROM

Low Power Operation

- 5V operation
- 150 μ A Active Current (100 kHz)
- 10 μ A Standby Current

Industry Standard Configuration

- Industrial Temperature -40° C to $+85^{\circ}$ C
- 8-pin SOIC or DIP

Description

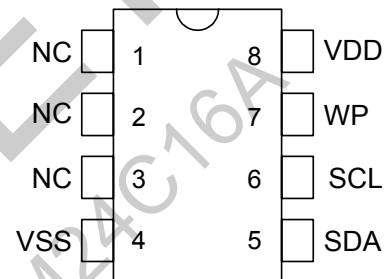
The FM24C16 is a 16-kilobit nonvolatile memory employing an advanced ferroelectric process. A ferroelectric random access memory or FRAM is nonvolatile and performs reads and writes like a RAM. It provides reliable data retention for 10 years while eliminating the complexities, overhead, and system level reliability problems caused by EEPROM and other nonvolatile memories.

Unlike serial EEPROMs, the FM24C16 performs write operations at bus speed. No write delays are incurred. Data is written to the memory array in the cycle after it has been successfully transferred to the device. The next bus cycle may commence immediately. The FM24C16 is capable of supporting 10^{10} read/write cycles, or 10,000 times more write cycles than EEPROM.

These capabilities make the FM24C16 ideal for nonvolatile memory applications requiring frequent or rapid writes. Examples range from data collection where the number of write cycles may be critical, to demanding industrial controls where the long write time of EEPROM can cause data loss. The combination of features allows more frequent data writing with less overhead for the system.

The FM24C16 provides substantial benefits to users of serial EEPROM, yet these benefits are available in a hardware drop-in replacement. The FM24C16 is provided in industry standard 8-pin packages using a two-wire protocol. The specifications are guaranteed over an industrial temperature range of -40° C to $+85^{\circ}$ C.

Pin Configuration



Pin Names	Function
SDA	Serial Data/address
SCL	Serial Clock
WP	Write Protect
VSS	Ground
VDD	Supply Voltage 5V

Ordering Information	
FM24C16-P	8-pin plastic DIP
FM24C16-S	8-pin SOIC

This product conforms to specifications per the terms of the Ramtron standard warranty. Production processing does not necessarily include testing of all parameters.

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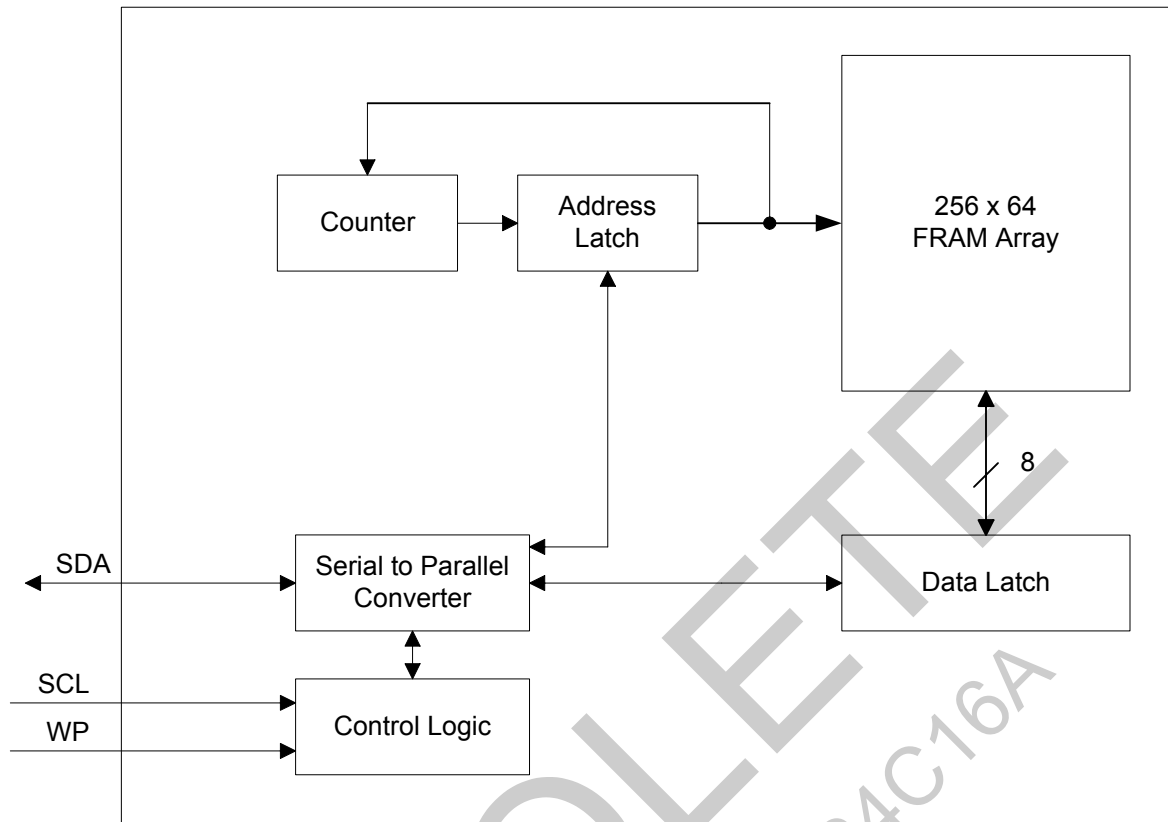


Figure 1. Block Diagram

Pin Description

Pin Name	I/O	Pin Description
SDA	I/O	Serial Data/Address: This is a bi-directional pin used to shift serial data and addresses for the two-wire interface. It employs an open-drain output and is intended to be wire-OR'd with other devices on the two-wire bus. The input buffer incorporates a Schmitt trigger for noise immunity and the output driver includes slope control for falling edges.
SCL	Input	Serial Clock: The serial clock input for the two-wire interface. Data is clocked out of the device on the SCL falling edge, and clocked in on the SCL rising edge.
WP	Input	Write Protect: When WP is high, addresses in the upper half of the logical memory map (A2=1 in the slave address) will be write-protected. Write access to the lower half of the addresses is permitted. When WP is low, all addresses may be written. This pin must not be left floating.
NC	-	No connect
VDD	Supply	Supply Voltage: 5V
VSS	Supply	Ground

Overview

The FM24C16 is a serial FRAM memory. The memory array is logically organized as a 2,048 x 8 memory array and is accessed using an industry standard two-wire interface. Functional operation of the FRAM is similar to serial EEPROMs. The major difference between the FM24C16 and a serial EEPROM with the same pinout relates to its superior write performance.

Memory Architecture

When accessing the FM24C16, the user addresses 2,048 locations each with 8 data bits. These data bits are shifted serially. The 2,048 addresses are accessed using the two-wire protocol, which includes a slave address (to distinguish other non-memory devices), a page address, and a word address. The word address consists of 8-bits that specify one of 256 addresses. The page address is 3-bits and so there are 8 pages each of 256 locations. The complete address of 11-bits specifies each byte address uniquely.

Most functions of the FM24C16 are either controlled by the two-wire interface or are handled automatically by on-board circuitry. The memory is read or written at the speed of the two-wire bus. Unlike an EEPROM, it is not necessary to poll the device for a ready condition since writes occur at bus speed. That is, by the time a new bus transaction can be shifted into the part, a write operation will be complete. This is explained in more detail in the interface section below.

Users expect several obvious system benefits from the FM24C16 due to its fast write cycle and high endurance as compared with EEPROM. However there are less obvious benefits as well. For example in a high noise environment, the fast write operation is less susceptible to corruption than an EEPROM since it is completed quickly. By contrast, an EEPROM requiring milliseconds to write is vulnerable to noise during much of the cycle.

Note that the FM24C16 contains no power management circuits other than a simple internal power-on reset. It is the user's responsibility to ensure that VDD is within data sheet tolerances to prevent incorrect operation.

Two-wire Interface

The FM24C16 employs a bi-directional two-wire bus protocol using few pins and little board space. Figure 2 illustrates a typical system configuration using the FM24C16 in a microcontroller-based system. The industry standard two-wire bus is familiar to many users but is described in this section.

By convention, any device that is sending data onto the bus is the transmitter while the target device for this data is the receiver. The device that is controlling the bus is the master. The master is responsible for generating the clock signal for all operations. Any device on the bus that is being controlled is a slave. The FM24C16 is always a slave device.

The bus protocol is controlled by transition states in the SDA and SCL signals. There are four conditions including start, stop, data bit, or acknowledge. Figure 3 illustrates the signal conditions that specify the four states. Detailed timing diagrams are in the electrical specifications.

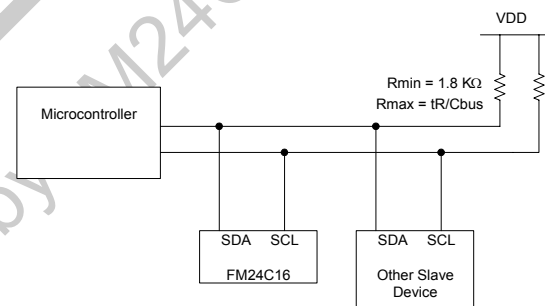


Figure 2. Typical System Configuration

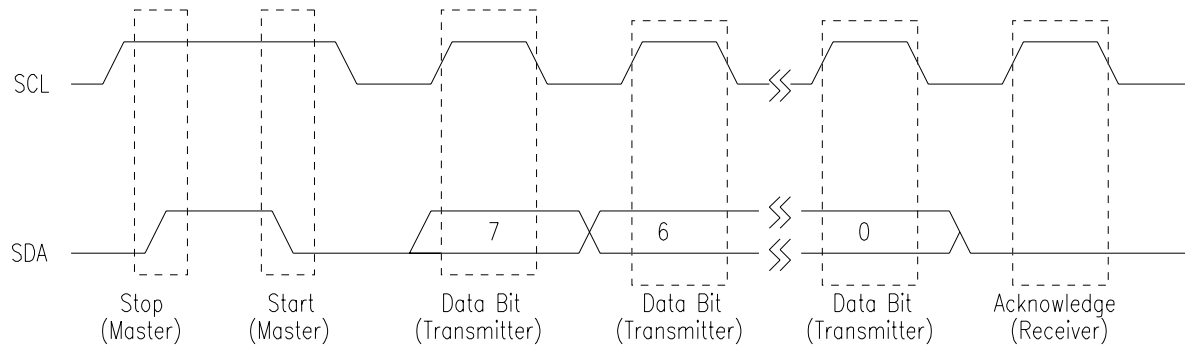


Figure 3. Data Transfer Protocol

Start Condition

A start condition is indicated when the bus master drives SDA from high to low while the SCL signal is high. All commands must be preceded by a start condition. An operation in progress can be aborted by asserting a start condition at any time. Aborting an operation using the start condition will ready the FM24C16 for a new operation.

If during operation the power supply drops below the specified VDD minimum, the system should issue a start condition prior to performing another operation.

Stop Condition

A stop condition is indicated when the bus master drives SDA from low to high while the SCL signal is high. All operations using the FM24C16 should end with a stop condition. If an operation is in progress when a stop is asserted, the operation will be aborted. The master must have control of SDA (not a memory read) in order to assert a stop condition.

Data/Address Transfer

All data transfers (including addresses) take place while the SCL signal is high. Except under the two conditions described above, the SDA signal should not change while SCL is high.

Acknowledge

The acknowledge takes place after the 8th data bit has been transferred in any transaction. During this state the transmitter should release the SDA bus to allow the receiver to drive it. The receiver drives the SDA signal low to acknowledge receipt of the byte. If the receiver does not drive SDA low, the condition is a no-acknowledge and the operation is aborted.

The receiver would fail to acknowledge for two distinct reasons. First is that a byte transfer fails. In this case the no-acknowledge ceases the current operation so that the part can be addressed again. This allows the last byte to be recovered in the event of a communication error.

Second and most common, the receiver does not acknowledge to deliberately end an operation. For example, during a read operation, the FM24C16 will continue to place data onto the bus as long as the receiver sends acknowledges (and clocks). When a read operation is complete and no more data is needed, the receiver must not acknowledge the last byte. If the receiver acknowledges the last byte, this will cause the FM24C16 to attempt to drive the bus on the next clock while the master is sending a new command such as stop.

Slave Address

The first byte that the FM24C16 expects after a start condition is the slave address. As shown in Figure 4, the slave address contains the device type, the page of memory to be accessed, and a bit that specifies if the transaction is a read or a write.

Bits 7-4 are the device type and should be set to 1010b for the FM24C16. The device type allows other types of functions to reside on the 2-wire bus within an identical address range. Bits 3-1 are the page select. They specify the 256-byte block of memory that is targeted for the current operation. Bit 0 is the read/write bit. A 0 indicates a write operation.

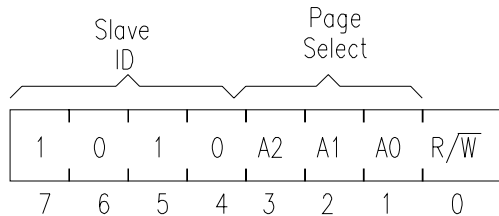


Figure 4. Slave Address

Word Address

After the FM24C16 (as receiver) acknowledges the slave ID, the master will place the word address on the bus for a write operation. The word address is the lower 8-bits of the address to be combined with the 3-bits of the page select to specify exactly the byte to be written. The complete 11-bit address is latched internally.

No word address occurs for a read operation, though the 3-bit page select is latched internally. Reads always use the lower 8-bits that are held internally in the address latch. That is, reads always begin at the address following the previous access. A random read address can be loaded by doing a write operation as explained below.

After transmission of each data byte, just prior to the acknowledge, the FM24C16 increments the internal address latch. This allows the next sequential byte to be accessed with no additional addressing. After the last address (7FFh) is reached, the address latch will roll over to 000h. There is no limit to the number of bytes that can be accessed with a single read or write operation.

Data Transfer

After all address information has been transmitted, data transfer between the bus master and the FM24C16 can begin. For a read operation the FM24C16 will place 8 data bits on the bus then wait for an acknowledge. If the acknowledge occurs, the next sequential byte will be transferred. If the acknowledge is not sent, the read operation is concluded. For a write operation, the FM24C16 will accept 8 data bits from the master then send an acknowledge. All data transfer occurs MSB (most significant bit) first.

Memory Operation

The FM24C16 is designed to operate in a manner very similar to other 2-wire interface memory products. The major differences result from the higher performance write capability of FRAM technology. These improvements result in some differences between the FM24C16 and a similar configuration EEPROM during writes. The complete operation for both writes and reads is explained below.

Write Operation

All writes begin with a slave ID then a word address as mentioned above. The bus master indicates a write operation by setting the LSB of the Slave ID to a 0. After addressing, the bus master sends each byte of data to the memory and the memory generates an acknowledge condition. Any number of sequential bytes may be written. If the end of the address range is reached internally, the address counter will wrap from 7FFh to 000h.

Unlike other nonvolatile memory technologies, there is no effective write delay with FRAM. Since the read and write access times of the underlying memory are the same, the user experiences no delay through the bus. The entire memory cycle occurs in less time than a single bus clock. Therefore any operation including read or write can occur immediately following a write. Acknowledge polling, a technique used with EEPROMs to determine if a write is complete is unnecessary and will always return a done condition.

An actual memory array write occurs after the 8th data bit is transferred. It will be complete before the acknowledge is sent. Therefore, if the user desires to abort a write without altering the memory contents, this should be done using start or stop condition prior to the 8th data bit. The FM24C16 needs no page buffering.

Portions of the memory array can be write protected using the WP pin. Setting the WP pin to a high condition (VDD) will write-protect addresses from 400h to 7FFh. The FM24C16 will not acknowledge data bytes that are written to protected addresses. In addition, the address counter will not increment if writes are attempted to these addresses. Setting WP to a low state (VSS) will deactivate this feature. WP should not be left floating.

Figure 5 below illustrates both a single-byte and multiple-write.

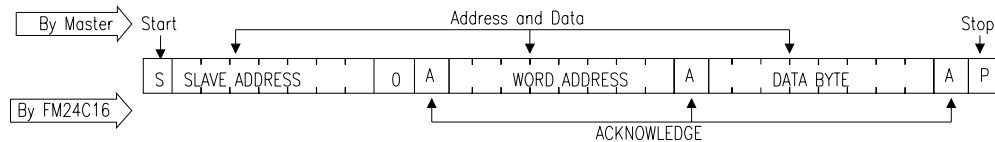


Figure 5 Byte Write

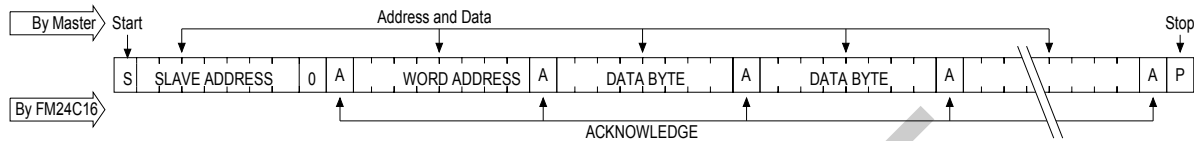


Figure 6 Multiple Byte Write

Read Operation

There are two basic types of read operations. They are current address read and selective address read. In a current address read, the FM24C16 uses the internal address latch to supply the lower 8 address bits. In a selective read, the user performs a procedure to set these lower address bits to a specific value.

Current Address & Sequential Read

As mentioned above the FM24C16 uses an internal latch to supply the lower 8 address bits for a read operation. A current address read uses the existing value in the address latch as a starting place for the read operation. This is the address immediately following that of the last operation.

To perform a current address read, the bus master supplies a slave address with the LSB set to 1. This indicates that a read operation is requested. The 3 page select bits in the slave ID specify the block of memory that is used for the read operation. On the next clock, the FM24C16 will begin shifting out data from the current address. The current address is the 3 bits from the slave ID combined with the 8 bits that were in the internal address latch.

Beginning with the current address, the bus master can read any number of bytes. Thus, a sequential read is simply a current address read with multiple byte transfers. After each byte the internal address counter will be incremented. Each time the bus master acknowledges a byte, this indicates that the FM24C16 should read out the next sequential byte. There are four ways to properly terminate a read operation. Failing to properly terminate the read will most likely create a bus contention as the FM24C16

attempts to read out additional data onto the bus. The four valid methods are as follows:

1. The bus master issues a no-acknowledge in the 9th clock cycle and a stop in the 10th clock cycle. This is illustrated in the diagrams below. This is the preferred method.
2. The bus master issues a no-acknowledge in the 9th clock cycle and a start in the 10th.
3. The bus master issues a stop in the 9th clock cycle. Bus contention may result.
4. The bus master issues a start in the 9th clock cycle. Bus contention may result.

If the internal address reaches 7FFh it will wrap around to 000h on the next read cycle. Figures 7 and 8 below show the proper operation for current address reads.

Selective (Random) Read

There is a simple technique that allows a user to select a random address location as the starting point for a read operation. This involves using the first two bytes of a write operation to set the internal address byte followed by subsequent read operations.

To perform a selective read, the bus master sends out the slave ID with the LSB set to 0. This specifies a write operation. According to the write protocol, the bus master then sends the word address byte that is loaded into the internal address latch. After the FM24C16 acknowledges the word address, the bus master issues a start condition. This simultaneously aborts the write operation and allows the read command to be issued with the slave ID LSB set to a 1. The operation is now a current address read. This operation is illustrated in Figure 9.

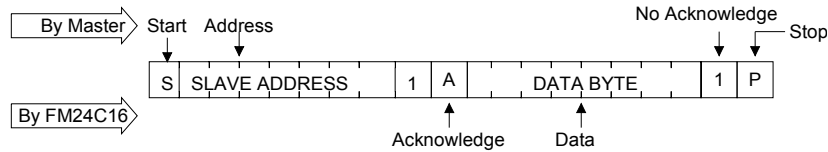


Figure 7 Current Address Read

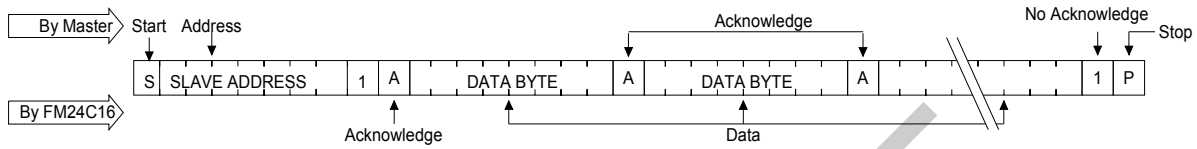


Figure 8 Sequential Read

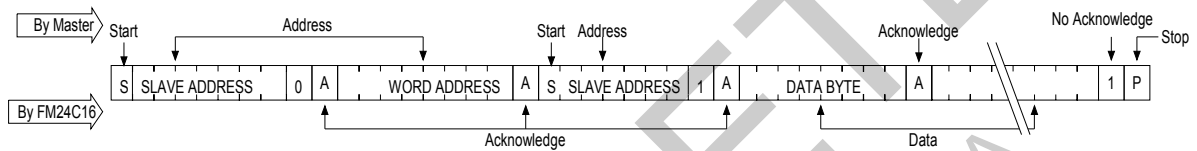


Figure 9 Selective (Random) Read

Endurance

Internally, a FRAM operates with a read and restore mechanism. Therefore, endurance cycles are applied for each read or write cycle. The FRAM architecture is based on an array of rows and columns. Rows are defined by A10-A3. Each access causes an endurance cycle for a row. Endurance can be optimized by

ensuring frequently accessed data is located in different rows. Regardless, FRAM read and write endurance is effectively unlimited at the 400kHz two-wire speed. Even at 30 accesses per second to the same row, 10 years time will elapse before 10 billion endurance cycles occur.

Applications

The versatility of FRAM technology fits into many diverse applications. Clearly the strength of higher write endurance and faster writes make FRAM superior to EEPROM in all but one-time programmable applications. The advantage is most obvious in data collection environments where writes are frequent and data must be nonvolatile.

The attributes of fast writes and high write endurance combine in many innovative ways. A short list of ideas is provided here.

1. Data collection. In applications where data is collected and saved, FRAM provides a superior alternative to other solutions. It is more cost effective than battery backup for SRAM and provides better write attributes than EEPROM.

2. Configuration. Any nonvolatile memory can retain a configuration. However, if the configuration changes and power failure is a possibility, the higher write endurance of FRAM allows changes to be recorded without restriction. Any time the system state is altered, the change can be written. This avoids writing to memory on power down when the available time is short and power scarce.

3. High noise environments. Writing to EEPROM in a noisy environment can be challenging. When severe noise or power fluctuations are present, the long write time of EEPROM creates a window of vulnerability during which the write can be corrupted. The fast write of FRAM is complete within a microsecond. This time is typically too short for noise or power fluctuation to disturb it.

4. Time to market. In a complex system, multiple software routines may need to access the nonvolatile memory. In this environment the time delay associated with programming EEPROM adds undue complexity to the software development. Each software routine must wait for complete programming before allowing access to the next routine. When time to market is critical, FRAM can eliminate this simple obstacle. As soon as a write is issued to the FM24C16, it is effectively done -- no waiting.

5. RF/ID. In the area of contactless memory, FRAM provides an ideal solution. Since RF/ID memory is powered by an RF field, the long programming time and high current consumption needed to write EEPROM is unattractive. FRAM provides a superior solution. The FM24C16 is suitable for multi-chip RF/ID products.

6. Maintenance tracking. In sophisticated systems, the operating history and system state during a failure is important knowledge. Maintenance can be expedited when this information has been recorded. Due to the high write endurance, FRAM makes an ideal system log. In addition, the convenient 2-wire interface of the FM24C16 allows memory to be distributed throughout the system using minimal additional resources.

Electrical Specifications

Absolute Maximum Ratings

Symbol	Description	Ratings
V_{DD}	Power Supply Voltage with respect to V_{SS}	-1.0V to +7.0V
V_{IN}	Voltage on any signal pin with respect to V_{SS}	-1.0V to +7.0V and $V_{IN} < V_{DD} + 1.0V$
T_{STG}	Storage temperature	-40°C to +85°C
T_{LEAD}	Lead temperature (Soldering, 10 seconds)	300° C

Stresses above those listed under Absolute Maximum Ratings may cause permanent damage to the device. This is a stress rating only, and the functional operation of the device at these or any other conditions above those listed in the operational section of this specification is not implied. Exposure to absolute maximum ratings conditions for extended periods may affect device reliability.

DC Operating Conditions ($T_A = -40^\circ\text{C}$ to $+85^\circ\text{C}$, $V_{DD} = 4.5\text{V}$ to 5.5V unless otherwise specified)

Symbol	Parameter	Min	Typ	Max	Units	Notes
V_{DD}	Main Power Supply	4.5	5.0	5.5	V	
I_{DD}	VDD Supply Current @ SCL = 100 kHz @ SCL = 400 kHz		115 400	150 500	μA μA	1
I_{SB}	Standby Current		1	10	μA	2
I_{LI}	Input Leakage Current			10	μA	3
I_{LO}	Output Leakage Current			10	μA	3
V_{IL}	Input Low Voltage	-0.3		$0.3 V_{DD}$	V	4
V_{IH}	Input High Voltage	$0.7 V_{DD}$		$V_{DD} + 0.5$	V	4
V_{OL}	Output Low Voltage @ $I_{OL} = 3\text{ mA}$ @ $I_{OL} = 6\text{ mA}$			0.4 0.6	V V	
V_{HYS}	Input Hysteresis	$0.05 V_{DD}$			V	4

Notes

1. SCL toggling between $V_{DD} - 0.3\text{V}$ and V_{SS} , other inputs V_{SS} or $V_{DD} - 0.3\text{V}$
2. SCL = SDA = V_{DD} . All inputs V_{SS} or V_{DD} . Stop command issued.
3. V_{IN} or $V_{OUT} = V_{SS}$ to V_{DD}
4. This parameter is periodically sampled and not 100% tested.

AC Parameters ($T_A = -40^\circ\text{C}$ to $+85^\circ\text{C}$, $V_{DD} = 4.5\text{V}$ to 5.5V unless otherwise specified)

Symbol	Parameter	Standard Mode		Fast Mode		Units	Notes
		Min	Max	Min	Max		
f_{SCL}	SCL Clock Frequency	0	100	0	400	kHz	
t_{SP}	Noise Suppression Time Constant on SCL, SDA		50		50	ns	
t_{AA}	SCL Low to SDA Data Out Valid		3		0.9	μs	
t_{BUF}	Bus Free Before New Transmission	4.7		1.3		μs	
$t_{HD:STA}$	Start Condition Hold Time	4.0		0.6		μs	
t_{LOW}	Clock Low Period	4.7		1.3		μs	
t_{HIGH}	Clock High Period	4.0		0.6		μs	
$t_{SU:STA}$	Start Condition Setup for Repeated Start	4.7		0.6		μs	
$t_{HD:DAT}$	Data In Hold	0		0		ns	
$t_{SU:DAT}$	Data In Setup	250		100		ns	
t_{RISE}	SDA and SCL Rise Time		1000	$20+0.1 C_b$	300	ns	1,2
t_{FALL}	SDA and SCL Fall Time		300	$20+0.1 C_b$	300	ns	1,2
$t_{SU:STO}$	Stop Condition Setup	4.0		0.6		μs	
t_{DH}	Data Output Hold (from SCL @ V_{IL})	0		0		ns	
t_{OF}	Output Fall Time (V_{IH} min to V_{IL} Max)		250	$20+0.1 C_b$	250	ns	1,2

Notes : All SCL specifications as well as start and stop conditions apply to both read and write operations.

- 1 This parameter is periodically sampled and not 100% tested.
- 2 C_b = Total capacitance of one bus line in pF.

Capacitance ($T_A = 25^\circ\text{C}$, $f=1.0\text{MHz}$, $V_{DD} = 5\text{V}$)

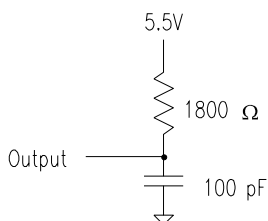
Symbol	Parameter	Max	Units	Notes
$C_{I/O}$	Input/output capacitance (SDA)	8	pF	1
C_{IN}	Input capacitance	6	pF	1

Notes

- 1 This parameter is periodically sampled and not 100% tested.

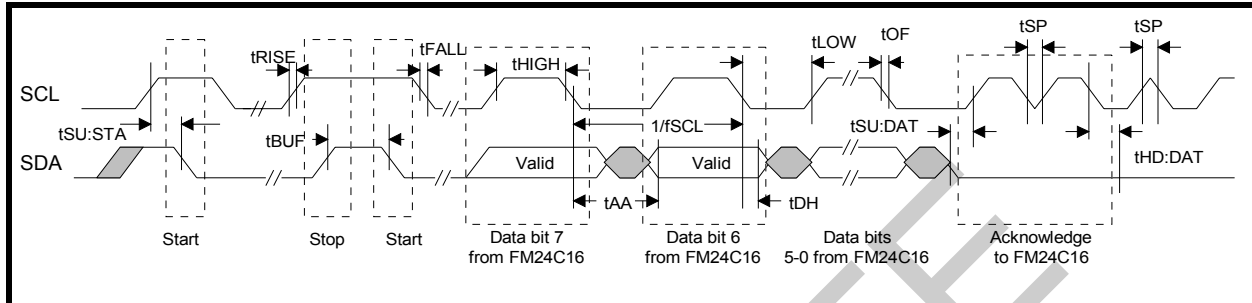
AC Test Conditions

Input Pulse Levels $0.1 V_{DD}$ to $0.9 V_{DD}$
 Input rise and fall times 10 ns
 Input and output timing levels $0.5 V_{DD}$

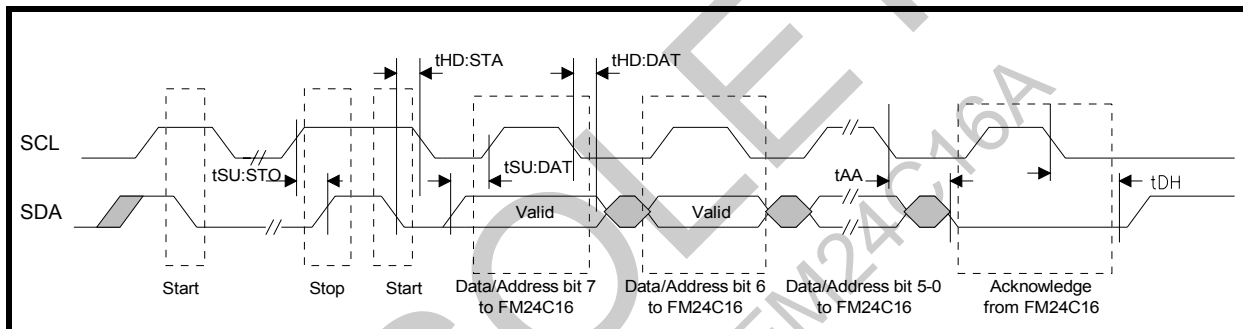
Equivalent AC Load Circuit**Diagram Notes**

All start and stop timing parameters apply to both read and write cycles. Clock specifications are identical for read and write cycles. Write timing parameters apply to slave address, word address, and write data bits. Functional relationships are illustrated in the relevant data sheet sections. These diagrams illustrate the timing parameters only.

Read Bus Timing



Write Bus Timing



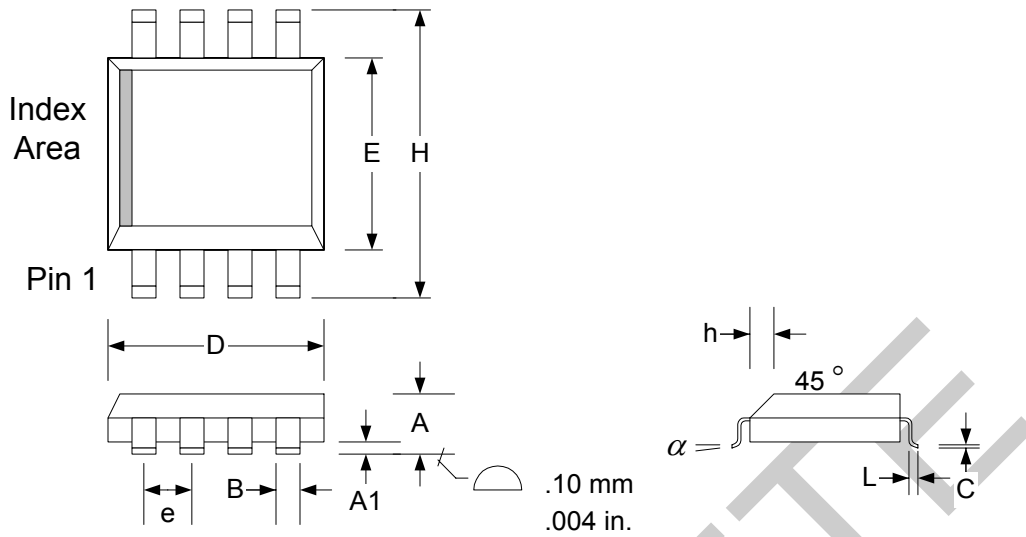
Data Retention ($V_{DD} = 4.5V$ to $5.5V$ unless otherwise specified)

Parameter	Min	Units	Notes
Data Retention	10	Years	1

Notes

1. The relationship between retention, temperature, and the associated reliability level is characterized separately.

8-pin SOIC (JEDEC Standard MS-012 variation AA)

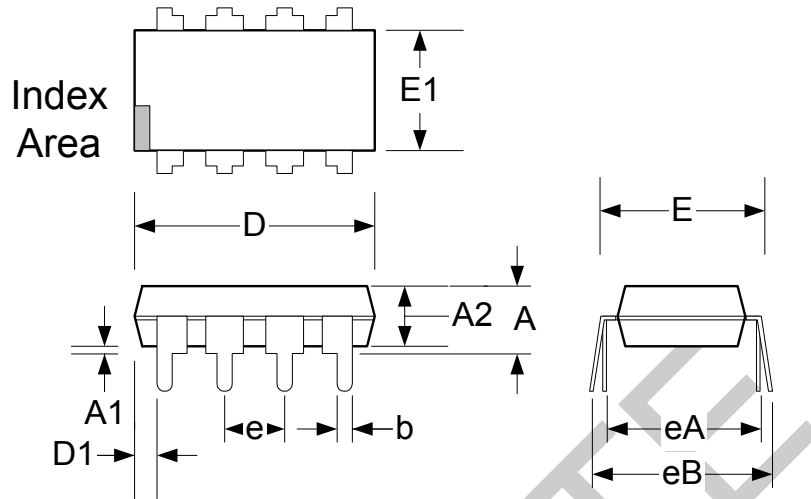


Selected Dimensions

Refer to JEDEC MS-012 for complete dimensions and notes.
 Controlling dimensions in millimeters.
 Conversions to inches are not exact.

Symbol	Dim	Min	Nom.	Max
A	mm in.	1.35 0.053		1.75 0.069
A1	mm in.	0.10 0.004		0.25 0.010
B	mm in.	0.33 0.013		0.51 0.020
C	mm in.	0.19 0.007		0.25 0.010
D	mm in.	4.80 0.189		5.00 0.197
E	mm in.	3.80 0.150		4.00 0.157
e	mm in.		1.27 BSC 0.050 BSC	
H	mm in.	5.80 0.228		6.20 0.244
h	mm in.	0.25 0.010		0.50 0.197
L	mm in.	0.40 0.016		1.27 0.050
α		0°		8°

8-pin DIP JEDEC MS-001



Selected Dimensions

Refer to JEDEC MS-001 for complete dimensions and notes.
 Controlling dimensions in inches.
 Conversions to millimeters are not exact.

Symbol	Dim	Min	Nom.	Max
A	in. mm			0.210 5.33
A1	in. mm	0.015 0.381		
A2	in. mm	0.115 2.92	0.130 3.30	0.195 4.95
b	in. mm	0.014 0.356	0.018 0.457	0.022 0.508
D	in. mm	0.355 9.02	0.365 9.27	0.400 10.2
D1	in. mm	0.005 0.127		
E	in. mm	0.300 7.62	0.310 7.87	0.325 8.26
E1	in. mm	0.240 6.10	0.250 6.35	0.280 7.11
e	in. mm		0.100 BSC 2.54 BSC	
eA	in. mm		0.300 BSC 7.62 BSC	
eB	in. mm			0.430 10.92
L	in. mm	0.115 2.92	0.130 3.30	0.150 3.81