

SYNCHRONOUS DRAM

MT48LC1M16A1 S - 512K x 16 x 2 banks

FEATURES

- PC100-compliant functionality
- Fully synchronous; all signals registered on positive edge of system clock
- Internal pipelined operation; column address can be changed every clock cycle
- Internal banks for hiding row access/precharge:
1 Meg x 16 - 512K x 16 x 2 banks architecture with 11 row, 8 column addresses per bank
- Programmable burst lengths: 1, 2, 4, 8 or full page
- Auto Precharge Mode
- Self Refresh and Adaptable Auto Refresh Modes
 - 32ms, 2,048-cycle refresh (15.6µs/row), or
 - 64ms, 2,048-cycle refresh (31.2µs/row), or
 - 64ms, 4,096-cycle refresh (15.6µs/row)
- LVTTTL-compatible inputs and outputs
- Single +3.3V ±0.3V power supply

OPTIONS

- Architecture
1 Meg x 16 (512K x 16 x 2 banks) **1M16**
- Plastic Package - OCPL
50-pin TSOP (400 mil) **TG**
- Timing
 - 8ns cycle time (≤ 125 MHz clock rate) **-8A**
 - 10ns cycle time (≤ 100 MHz clock rate) **-10**
 - 12ns cycle time (≤ 83 MHz clock rate) **-12**
- Refresh
2K or 4K with Self Refresh Mode at 64ms **S**
- Part Number Example: MT48LC1M16A1TG-8A S

MARKING

KEY TIMING PARAMETERS

| SPEED GRADE | CLOCK FREQUENCY | ACCESS TIME | | SETUP TIME | HOLD TIME |
|-------------|-----------------|-------------|---------|------------|-----------|
| | | *CL = 2 | *CL = 3 | | |
| -8A | 125 MHz | 9ns | 6ns | 2ns | 1ns |
| -10 | 100 MHz | 9ns | 7.5ns | 3ns | 1ns |
| -12 | 83 MHz | 9ns | 9ns | 3ns | 1ns |

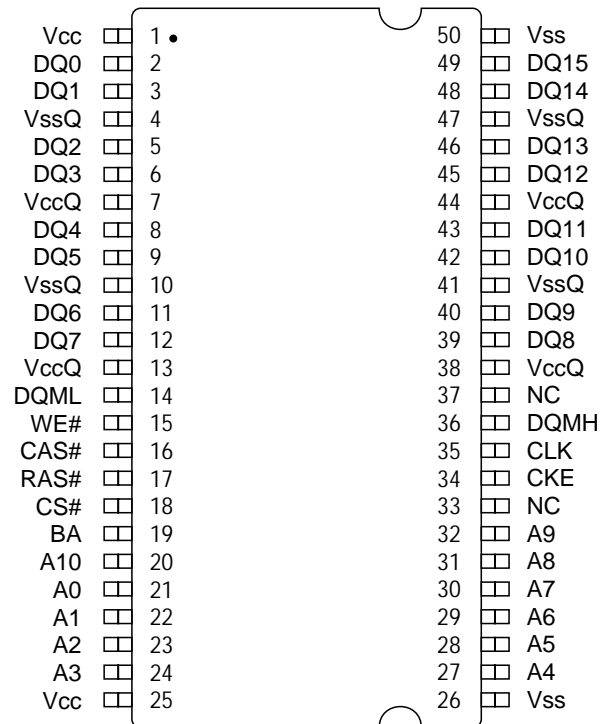
*CL = CAS (READ) Latency

16Mb (x16) SDRAM PART NUMBER

| PART NUMBER | ARCHITECTURE |
|------------------|----------------------------------|
| MT48LC1M16A1TG S | 1 Meg x 16 (512K x 16 x 2 banks) |

PIN ASSIGNMENT (Top View)

50-Pin TSOP



Note: The # symbol indicates signal is active LOW.

| | |
|-------------------|---------------------|
| Configuration | 512K x 16 x 2 banks |
| Refresh Count | 2K or 4K |
| Row Addressing | 2K (A0-A10) |
| Bank Addressing | 1 (BA) |
| Column Addressing | 256 (A0-A7) |

GENERAL DESCRIPTION

The 16Mb SDRAM is a high-speed CMOS dynamic random access memory containing 16,777,216 bits. It is internally configured as a dual 512K x 16 DRAM with a synchronous interface (all signals are registered on the positive edge of the clock signal, CLK). Each of the 512K x 16 bit banks is organized as 2,048 rows by 256 columns by 16 bits. Read and write accesses to the SDRAM are burst oriented; accesses start at a selected location and continue for a programmed

GENERAL DESCRIPTION (continued)

number of locations in a programmed sequence. Accesses begin with the registration of an ACTIVE command, which is then followed by a READ or WRITE command. The address bits registered coincident with the ACTIVE command are used to select the bank and row to be accessed (BA selects the bank, A0-A10 select the row). The address bits registered coincident with the READ or WRITE command are used to select the starting column location for the burst access.

The SDRAM provides for programmable READ or WRITE burst lengths of 1, 2, 4 or 8 locations, or the full page, with a BURST TERMINATE option. An AUTO PRECHARGE function may be enabled to provide a self-timed row precharge that is initiated at the end of the burst sequence.

The 1 Meg x 16 SDRAM uses an internal pipelined architecture to achieve high-speed operation. This architec-

ture is compatible with the 2n rule of prefetch architectures, but it also allows the column address to be changed on every clock cycle to achieve a high-speed fully random access. Precharging one bank while accessing the alternate bank will hide the PRECHARGE cycles and provide seamless, high-speed, random-access operation.

The 1 Meg x 16 SDRAM is designed to operate in 3.3V, low-power memory systems. An auto refresh mode is provided, along with a power-saving power-down mode. All inputs and outputs are LVTTTL-compatible.

SDRAMs offer substantial advances in DRAM operating performance, including the ability to synchronously burst data at a high data rate with automatic column-address generation, the ability to interleave between internal banks in order to hide precharge time, and the capability to randomly change column addresses on each clock cycle during a burst access.

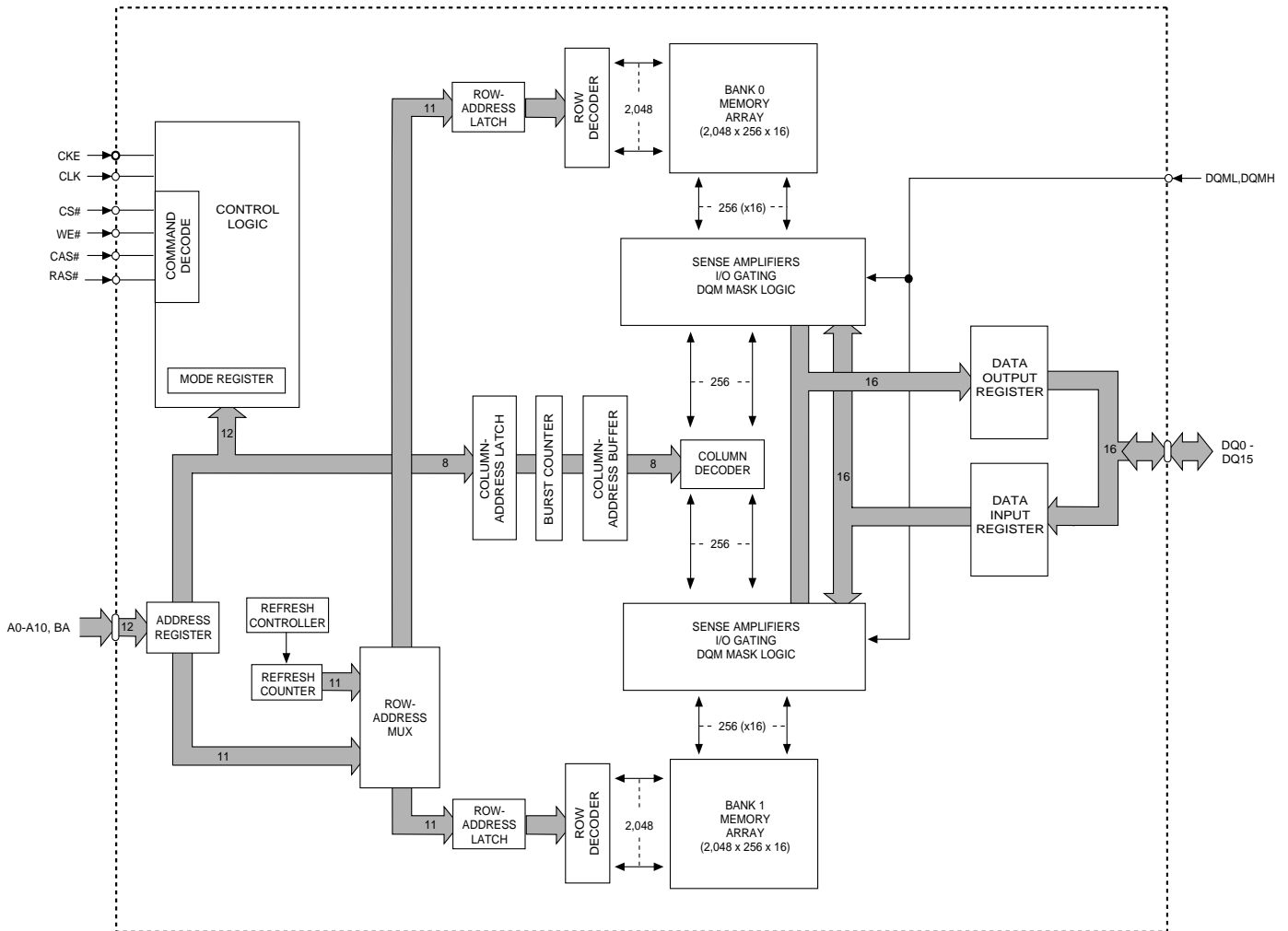
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FUNCTIONAL BLOCK DIAGRAM
1 Meg x 16 SDRAM



PIN DESCRIPTIONS

| TSOP PIN NUMBERS | SYMBOL | TYPE | DESCRIPTION |
|--|-------------------|------------------|---|
| 35 | CLK | Input | Clock: CLK is driven by the system clock. All SDRAM input signals are sampled on the positive edge of CLK. CLK also increments the internal burst counter and controls the output registers. |
| 34 | CKE | Input | Clock Enable: CKE activates (HIGH) and deactivates (LOW) the CLK signal. Deactivating the clock provides PRECHARGE POWER-DOWN and SELF REFRESH operations (all banks idle), ACTIVE POWER-DOWN (row ACTIVE in either bank) or CLOCK SUSPEND operation (burst/access in progress). CKE is synchronous except after the device enters power-down and self refresh modes, where CKE becomes asynchronous until after exiting the same mode. The input buffers, including CLK, are disabled during power-down and self refresh modes, providing low standby power. CKE may be tied HIGH. |
| 18 | CS# | Input | Chip Select: CS# enables (registered LOW) and disables (registered HIGH) the command decoder. All commands are masked when CS# is registered HIGH. CS# provides for external bank selection on systems with multiple banks. CS# is considered part of the command code. |
| 17, 15, 16 | RAS#, WE# CAS# | Input | Command Inputs: RAS#, CAS#, and WE# (along with CS#) define the command being entered. |
| 14, 36 | DQML, DQMH | Input | Input/Output Mask: DQM is an input mask signal for write accesses and an output enable signal for read accesses. Input data is masked when DQM is sampled HIGH during a WRITE cycle. The output buffers are placed in a High-Z state (two-clock latency) when DQM is sampled HIGH during a READ cycle. DQML corresponds to DQ0-DQ7; DQMH corresponds to DQ8-DQ15. DQML and DQMH are considered same state when referenced as DQM. |
| 19 | BA | Input | Bank Address Inputs: BA defines to which bank the ACTIVE, READ, WRITE or PRECHARGE command is being applied. BA is also used to program the 12th bit of the Mode Register. |
| 21-24, 27-32, 20 | A0-A10 | Input | Address Inputs: A0-A10 are sampled during the ACTIVE command (row address A0-A10) and READ/WRITE command (column address A0-A7, with A10 defining AUTO PRECHARGE) to select one location out of the 512K available in the respective bank. A10 is sampled during a PRECHARGE command to determine if all banks are to be precharged (A10 HIGH). The address inputs also provide the op-code during a LOAD MODE REGISTER command. |
| 2, 3, 5, 6, 8, 9, 11, 12, 39, 40, 42, 43, 45, 46, 48, 49 | DQ0- DQ15 | Input/ Output | Data I/O: Data bus. |
| 33, 37 | NC | — | No Connect: These pins should be left unconnected. |
| 7, 13, 38, 44 | VccQ | Supply | DQ Power: Provide isolated power to DQs for improved noise immunity. |
| 4, 10, 41, 47 | VssQ | Supply | DQ Ground: Provide isolated ground to DQs for improved noise immunity. |
| 1, 25 | Vcc | Supply | Power Supply: +3.3V ±0.3V. |
| 26, 50 | Vss | Supply | Ground. |

FUNCTIONAL DESCRIPTION

In general, the SDRAM is a dual 512K x 16 DRAM that operates at 3.3V and includes a synchronous interface (all signals are registered on the positive edge of the clock signal, CLK). Each of the 512K x 16 bit banks is organized as 2,048 rows by 256 columns by 16 bits.

Read and write accesses to the SDRAM are burst oriented; accesses start at a selected location and continue for a programmed number of locations in a programmed sequence. Accesses begin with the registration of an ACTIVE command, which is then followed by a READ or WRITE command. The address bits registered coincident with the ACTIVE command are used to select the bank and row to be accessed (BA selects the bank, A0-A10 select the row). The address bits (A0-A7) registered coincident with the READ or WRITE command are used to select the starting column location for the burst access.

Prior to normal operation, the SDRAM must be initialized. The following sections provide detailed information covering device initialization, register definition, command descriptions and device operation.

INITIALIZATION

SDRAMs must be powered up and initialized in a pre-defined manner. Operational procedures other than those specified may result in undefined operation. Once power is applied to V_{cc} and V_{ccQ} (simultaneously) and the clock is stable, the SDRAM requires a 100μs delay prior to applying an executable command. Starting at some point during this 100μs period and continuing at least through the end of this period, COMMAND INHIBIT or NOP commands should be applied.

Once the 100μs delay has been satisfied, with at least one COMMAND INHIBIT or NOP command having been applied, a PRECHARGE command should be applied. All banks must be precharged, thereby placing the device in the All Banks Idle state.

Once in the Idle state, two AUTO REFRESH cycles must be performed. Once the AUTO REFRESH cycles are complete, the SDRAM is ready for Mode Register programming. Because the Mode Register will power up in an unknown state, it should be loaded prior to applying any operational command.

REGISTER DEFINITION

MODE REGISTER

The Mode Register is used to define the specific mode of operation of the SDRAM. This definition includes the selection of a burst length, a burst type, a CAS latency, an operating mode, and a write burst mode, as shown in

Figure 1. The Mode Register is programmed via the LOAD MODE REGISTER command and will retain the stored information until it is programmed again or the device loses power.

Mode Register bits M0-M2 specify the burst length, M3 specifies the type of burst (sequential or interleaved), M4-M6 specify the CAS latency, M7 and M8 specify the operating mode, M9 specifies the write burst mode, and M10 and M11 are reserved for future use.

The Mode Register must be loaded when all banks are idle, and the controller must wait the specified time before initiating the subsequent operation. Violating either of these requirements will result in unspecified operation.

Burst Length

Read and write accesses to the SDRAM are burst oriented, with the burst length being programmable, as shown in Figure 1. The burst length determines the maximum number of column locations that can be accessed for a given READ or WRITE command. Burst lengths of 1, 2, 4 or 8 locations are available for both the sequential and the interleaved burst types, and a full-page burst is available for the sequential type. The full-page burst is used in conjunction with the BURST TERMINATE command to generate arbitrary burst lengths.

Reserved states should not be used, as unknown operation or incompatibility with future versions may result.

When a READ or WRITE command is issued, a block of columns equal to the burst length is effectively selected. All accesses for that burst take place within this block, meaning that the burst will wrap within the block if a boundary is reached. The block is uniquely selected by A1-A7 when the burst length is set to two, by A2-A7 when the burst length is set to four and by A3-A7 when the burst length is set to eight. The remaining (least significant) address bit(s) are used to select the starting location within the block. Full-page bursts wrap within the page if the boundary is reached.

Burst Type

Accesses within a given burst may be programmed to be either sequential or interleaved; this is referred to as the burst type and is selected via bit M3.

The ordering of accesses within a burst is determined by the burst length, the burst type and the starting column address, as shown in Table 1.

Read Latency

The READ latency is the delay, in clock cycles, between the registration of a READ command and the availability of the first piece of output data. The latency can be set to 1, 2 or 3 clocks.

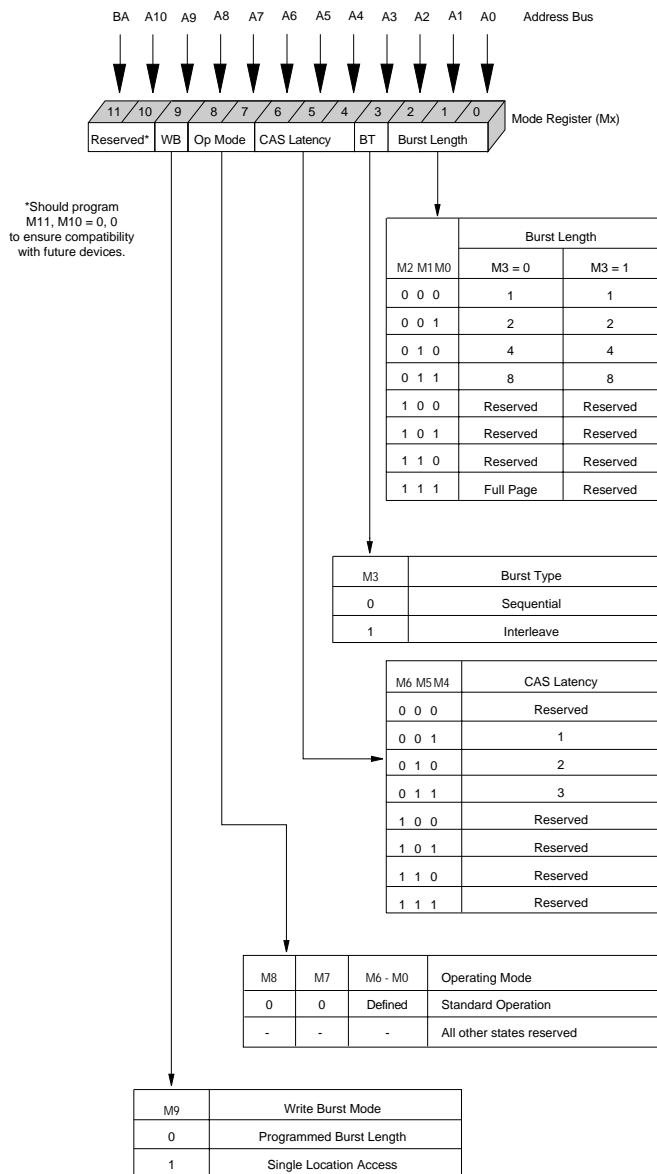


Figure 1
MODE REGISTER DEFINITION

Table 1
BURST DEFINITION

| Burst Length | Starting Column Address: | Order of Accesses Within a Burst | |
|-----------------|----------------------------|--|--------------------|
| | | Type = Sequential | Type = Interleaved |
| 2 | A0 | | |
| | 0 | 0-1 | 0-1 |
| | 1 | 1-0 | 1-0 |
| 4 | A1 A0 | | |
| | 0 0 | 0-1-2-3 | 0-1-2-3 |
| | 0 1 | 1-2-3-0 | 1-0-3-2 |
| | 1 0 | 2-3-0-1 | 2-3-0-1 |
| | 1 1 | 3-0-1-2 | 3-2-1-0 |
| 8 | A2 A1 A0 | | |
| | 0 0 0 | 0-1-2-3-4-5-6-7 | 0-1-2-3-4-5-6-7 |
| | 0 0 1 | 1-2-3-4-5-6-7-0 | 1-0-3-2-5-4-7-6 |
| | 0 1 0 | 2-3-4-5-6-7-0-1 | 2-3-0-1-6-7-4-5 |
| | 0 1 1 | 3-4-5-6-7-0-1-2 | 3-2-1-0-7-6-5-4 |
| | 1 0 0 | 4-5-6-7-0-1-2-3 | 4-5-6-7-0-1-2-3 |
| | 1 0 1 | 5-6-7-0-1-2-3-4 | 5-4-7-6-1-0-3-2 |
| | 1 1 0 | 6-7-0-1-2-3-4-5 | 6-7-4-5-2-3-0-1 |
| 1 1 1 | 7-0-1-2-3-4-5-6 | 7-6-5-4-3-2-1-0 | |
| Full Page (256) | n = A0-A7 (location 0-255) | Cn, Cn+1, Cn+2 Cn+3, Cn+4... ...Cn-1, Cn... | Not supported |

- NOTE:**
1. For a burst length of two, A1-A7 select the block of two burst; A0 selects the starting column within the block.
 2. For a burst length of four, A2-A7 select the block of four burst; A0-A1 select the starting column within the block.
 3. For a burst length of eight, A3-A7 select the block of eight burst; A0-A2 select the starting column within the block.
 4. For a full-page burst, the full row is selected and A0-A7 select the starting column.
 5. Whenever a boundary of the block is reached within a given sequence above, the following access wraps within the block.
 6. For a burst length of one, A0-A7 select the unique column to be accessed, and Mode Register bit M3 is ignored.

If a READ command is registered at clock edge n , and the latency is m clocks, the data will be available by clock edge $n + m$. The DQs will start driving as a result of the clock edge one cycle earlier ($n + m - 1$), and provided that the relevant access times are met, the data will be valid by clock edge $n + m$. For example, assuming that the clock cycle time is such that all relevant access times are met, if a READ command is registered at T0, and the latency is programmed to two clocks, the DQs will start driving after T1 and the data will be valid by T2, as shown in Figure 2. Table 2 below indicates the operating frequencies at which each CAS latency setting can be used.

Reserved states should not be used, as unknown operation or incompatibility with future versions may result.

Operating Mode

The normal operating mode is selected by setting M7 and M8 to zero; the other combinations of values for M7 and M8 are reserved for future use and/or test modes. The programmed burst length applies to both READ and WRITE bursts.

Test modes and reserved states should not be used because unknown operation or incompatibility with future versions may result.

Write Burst Mode

When M9 = 0, the burst length programmed via M0-M2 applies to both READ and WRITE bursts; when M9 = 1, the programmed burst length applies to READ bursts, but write accesses are single-location (nonburst) accesses.

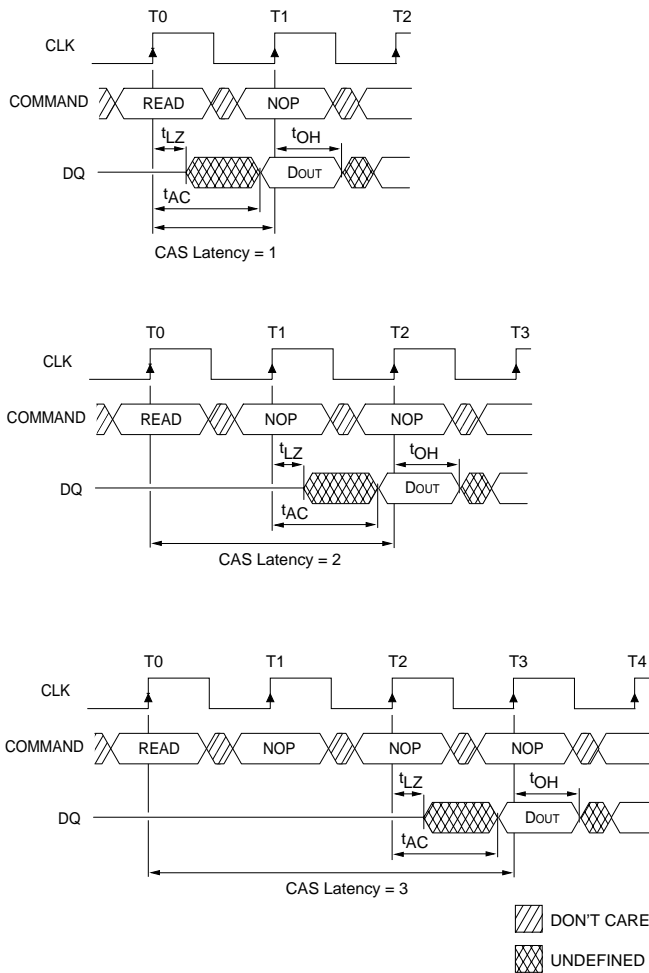


Figure 2
CAS LATENCY

Table 2
CAS LATENCY

| SPEED | ALLOWABLE OPERATING FREQUENCY (MHz) | | |
|-------|-------------------------------------|-----------------|-----------------|
| | CAS LATENCY = 1 | CAS LATENCY = 2 | CAS LATENCY = 3 |
| -8A | ≤ 33 | ≤ 76 | ≤ 125 |
| -10 | ≤ 33 | ≤ 66 | ≤ 100 |
| -12 | ≤ 33 | ≤ 66 | ≤ 83 |

COMMANDS

Truth Table 1 provides a quick reference of available commands. This is followed by a verbal description of each command. Two additional Truth Tables appear following

the Operation section; these tables provide current state/next state information.

TRUTH TABLE 1 – Commands and DQM Operation

(Notes: 1)

| NAME (FUNCTION) | CS# | RAS# | CAS# | WE# | DQM | ADDR | DQs | NOTES |
|--|-----|------|------|-----|-----|----------|--------|-------|
| COMMAND INHIBIT (NOP) | H | X | X | X | X | X | X | |
| NO OPERATION (NOP) | L | H | H | H | X | X | X | |
| ACTIVE (select bank and activate row) | L | L | H | H | X | Bank/Row | X | 3 |
| READ (select bank and column and start READ burst) | L | H | L | H | X | Bank/Col | X | 4 |
| WRITE (select bank and column and start WRITE burst) | L | H | L | L | X | Bank/Col | Valid | 4 |
| BURST TERMINATE | L | H | H | L | X | X | Active | |
| PRECHARGE (deactivate row in bank or banks) | L | L | H | L | X | Code | X | 5 |
| AUTO REFRESH or SELF REFRESH (enter self refresh mode) | L | L | L | H | X | X | X | 6, 7 |
| LOAD MODE REGISTER | L | L | L | L | X | OpCode | X | 2 |
| Write Enable/Output Enable | - | - | - | - | L | - | Active | 8 |
| Write Inhibit/Output High-Z | - | - | - | - | H | - | High-Z | 8 |

- NOTE:**
1. CKE is HIGH for all commands shown except SELF REFRESH.
 2. A0-A10 and BA define the op-code written to the Mode Register.
 3. A0-A10 provide row address, and BA determines which bank is made active.
 4. A0-A7 provide column address; A10 HIGH enables the AUTO PRECHARGE feature (nonpersistent), while A10 LOW disables the AUTO PRECHARGE feature; BA determines which bank is being read from or written to.
 5. For A10 LOW, BA determines bank being precharged; for A10 HIGH, all banks are precharged and BA is a "Don't Care."
 6. This command is AUTO REFRESH if CKE is HIGH, SELF REFRESH if CKE is LOW.
 7. Internal refresh counter controls row addressing; all inputs and I/Os are "Don't Care" except for CKE.
 8. Activates or deactivates the DQs during WRITES (zero-clock delay) and READS (two-clock delay).

COMMAND INHIBIT

The COMMAND INHIBIT function prevents new commands from being executed by the SDRAM, regardless of whether the CLK signal is enabled. The SDRAM is effectively deselected. Operations already in progress are not affected.

NO OPERATION (NOP)

The NO OPERATION (NOP) command is used to perform a NOP to an SDRAM which is selected (CS# is LOW). This prevents unwanted commands from being registered during idle or wait states. Operations already in progress are not affected.

LOAD MODE REGISTER

The Mode Register is loaded via inputs A0-A10 and BA. See Mode Register heading in Register Definition section. The LOAD MODE REGISTER command can only be issued when all banks are idle, and a subsequent executable command cannot be issued until ^tMTC is met.

ACTIVE

The ACTIVE command is used to open (or activate) a row in a particular bank for a subsequent access. The value on the BA input selects the bank, and the address provided on inputs A0-A10 selects the row. This row remains active (or open) for accesses until a PRECHARGE command is issued to that bank. A PRECHARGE command must be issued before opening a different row in the same bank.

READ

The READ command is used to initiate a burst read access to an active row. The value on the BA input selects the bank, and the address provided on inputs A0-A7 selects the starting column location. The value on input A10 determines whether or not AUTO PRECHARGE is used. If AUTO PRECHARGE is selected, the row being accessed will be precharged at the end of the READ burst; if AUTO PRECHARGE is not selected, the row will remain open for subsequent accesses. Read data appears on the DQs, subject to the logic level on the DQM inputs two clocks earlier. If a given DQM signal was registered HIGH, the corresponding DQs will be High-Z two clocks later; if the DQM signal was registered LOW, the DQs will provide valid data.

WRITE

The WRITE command is used to initiate a burst write access to an active row. The value on the BA input selects the bank, and the address provided on inputs A0-A7 selects the starting column location. The value on input A10 determines whether or not AUTO PRECHARGE is used. If

AUTO PRECHARGE is selected, the row being accessed will be precharged at the end of the WRITE burst; if AUTO PRECHARGE is not selected, the row will remain open for subsequent accesses. Input data appearing on the DQs is written to the memory array subject to the DQM input logic level appearing coincident with the data. If a given DQM signal is registered LOW, the corresponding data will be written to memory; if the DQM signal is registered HIGH, the corresponding data inputs will be ignored, and a WRITE will not be executed to that byte/column location.

PRECHARGE

The PRECHARGE command is used to deactivate the open row in a particular bank or the open row in all banks. The bank(s) will be available for a subsequent row access a specified time (^tRP) after the PRECHARGE command is issued. Input A10 determines whether one or all banks are to be precharged, and in the case where only one bank is to be precharged, input BA selects the bank. Otherwise BA is treated as "Don't Care." Once a bank has been precharged, it is in the Idle state and must be activated prior to any READ or WRITE commands being issued to that bank.

AUTO PRECHARGE

AUTO PRECHARGE is a feature which performs the same individual-bank PRECHARGE function described above, but without requiring an explicit command. This is accomplished by using A10 to enable AUTO PRECHARGE in conjunction with a specific READ or WRITE command. A precharge of the bank/row that is addressed with the READ or WRITE command is automatically performed upon completion of the READ or WRITE burst, except in the full-page burst mode, where AUTO PRECHARGE does not apply. AUTO PRECHARGE is nonpersistent in that it is either enabled or disabled for each individual READ or WRITE command.

AUTO PRECHARGE ensures that the PRECHARGE is initiated at the earliest valid stage within a burst. The user must not issue another command to the same bank until the precharge time (^tRP) is completed. This is determined as if an explicit PRECHARGE command was issued at the earliest possible time, as described for each burst type in the Operation section of this data sheet.

BURST TERMINATE

The BURST TERMINATE command is used to truncate either fixed-length or full-page bursts. The most recently registered READ or WRITE command prior to the BURST TERMINATE command will be truncated as shown in the Operation section of this data sheet.

AUTO REFRESH

AUTO REFRESH is used during normal operation of the SDRAM and is analogous to CAS#-BEFORE-RAS# (CBR) REFRESH in conventional DRAMs. This command is non-persistent, so it must be issued each time a refresh is required.

The addressing is generated by the internal refresh controller. This makes the address bits a "Don't Care" during an AUTO REFRESH command. The 1 Meg x 16 SDRAM requires 2,048 AUTO REFRESH cycles every 64ms (t_{REF}). Providing a distributed AUTO REFRESH command every 31.25 μ s will meet the refresh requirement and ensure that each row is refreshed. Providing 2,048 AUTO REFRESH cycles every 32ms (15.6 μ s per AUTO REFRESH) will meet the refresh requirement and ensure that each row is refreshed. Alternatively, 2,048 AUTO REFRESH commands can be issued in a burst at the minimum cycle rate (t_{RC}) once every 32ms or 64ms.

The refresh scheme is compatible with the 4K refresh requirements at the standard rate of 15.625 μ s per row for 64ms period. Thus, the Micron 1 Meg x 16 SDRAM can be used in either a 2K or a 4K refresh memory. Of the three types of refresh options, utilizing the 2,048 cycles every 64ms (31.25 μ s per refresh) provides the maximum power savings.

SELF REFRESH

The SELF REFRESH command can be used to retain data in the SDRAM, even if the rest of the system is powered down. When in the self refresh mode, the SDRAM retains data without external clocking. The SELF REFRESH command is initiated like an AUTO REFRESH command except CKE is disabled (LOW). Once the SELF REFRESH command is registered, all the inputs to the SDRAM become "Don't Care," with the exception of CKE, which must remain LOW.

Once self refresh mode is engaged, the SDRAM provides its own internal clocking, causing it to perform its own auto refresh cycles. The SDRAM must remain in self refresh mode for a minimum period equal to t_{RAS} , and may remain in self refresh mode for an indefinite period beyond that.

The procedure for exiting self refresh requires a sequence of commands. First, CLK must be stable prior to CKE going back HIGH. Once CKE is HIGH, the SDRAM must have NOP commands issued (a minimum of two clocks) for t_{XSR} , because time is required for the completion of any internal refresh in progress.

If, during normal operation, AUTO REFRESH cycles are issued in bursts (as opposed to being evenly distributed), a burst of 2,048 AUTO REFRESH cycles should be completed just prior to entering and just after exiting the self refresh mode.

OPERATION

BANK/ROW ACTIVATION

Before any READ or WRITE commands can be issued to a bank within the SDRAM, a row in that bank must be "opened." This is accomplished via the ACTIVE command, which selects both the bank and the row to be activated.

After opening a row (issuing an ACTIVE command) a READ or WRITE command may be issued to that row, subject to the t_{RCD} specification. t_{RCD} (MIN) should be divided by the clock period and rounded up to the next whole number to determine the earliest clock edge after the ACTIVE command on which a READ or WRITE command can be entered. For example, a t_{RCD} specification of 30ns with a 90 MHz clock (11.1ns period) results in 2.7 clocks rounded to 3. This is reflected in Figure 4, which covers any case where $2 < t_{RCD} \text{ (MIN)} / t_{CK} \leq 3$. (The same procedure is used to convert other specification limits from time units to clock cycles.)

A subsequent ACTIVE command to a different row in the same bank can only be issued after the previous active row has been "closed" (precharged). The minimum time interval between successive ACTIVE commands to the same bank is defined by t_{RC} .

A subsequent ACTIVE command to another bank can be issued while the first bank is being accessed, which results in a reduction of total row access overhead. The minimum time interval between successive ACTIVE commands to different banks is defined by t_{RRD} .

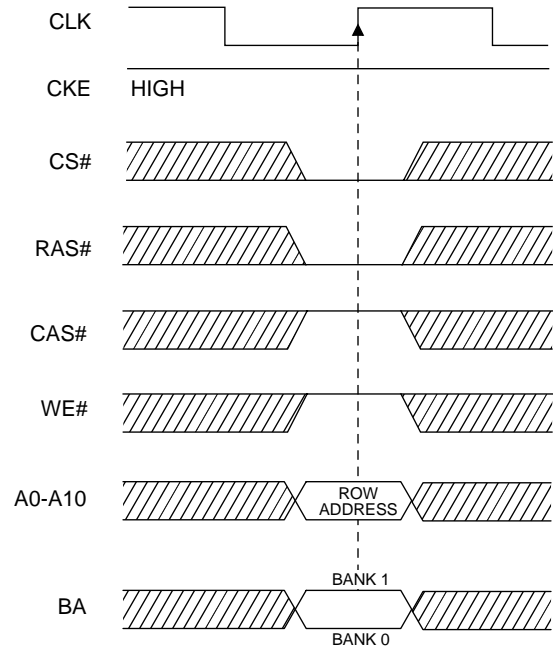


Figure 3
ACTIVATING A SPECIFIC ROW IN A
SPECIFIC BANK

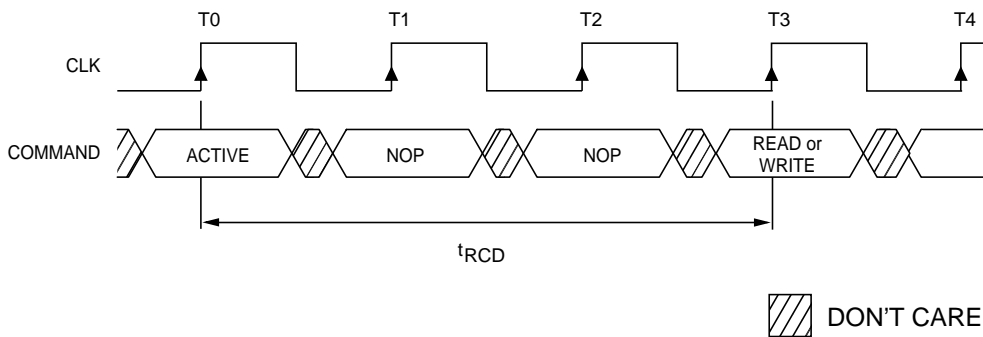


Figure 4
EXAMPLE: MEETING t_{RCD} (MIN) WHEN $2 < t_{RCD} \text{ (MIN)} / t_{CK} \leq 3$

READS

READ bursts are initiated with a READ command, as shown in Figure 5.

The starting column and bank addresses are provided with the READ command and AUTO PRECHARGE is either enabled or disabled for that burst access. If AUTO PRECHARGE is enabled, the row being accessed is precharged at the completion of the burst. For the generic READ commands used in the following illustrations, AUTO PRECHARGE is disabled.

During READ bursts, the valid data-out element from the starting column address will be available following the CAS latency after the READ command. Each subsequent data-out element will be valid by the next positive clock edge. Figure 6 shows general timing for each possible CAS latency setting.

Upon completion of a burst, assuming no other commands have been initiated, the DQs will go High-Z. A full-page burst will continue until terminated (at the end of the page it will wrap to column 0 and continue).

Data from any READ burst may be truncated with a subsequent READ command, and data from a fixed-length READ burst may be immediately followed by data from a subsequent READ command. In either case, a continuous flow of data can be maintained. The first data element from the new burst follows either the last element of a completed burst, or the last desired data element of a longer burst which is being truncated. The new READ command should be issued x cycles before the clock edge at which the last desired data element is valid, where x equals the CAS latency minus one. This is shown in Figure 7 for READ

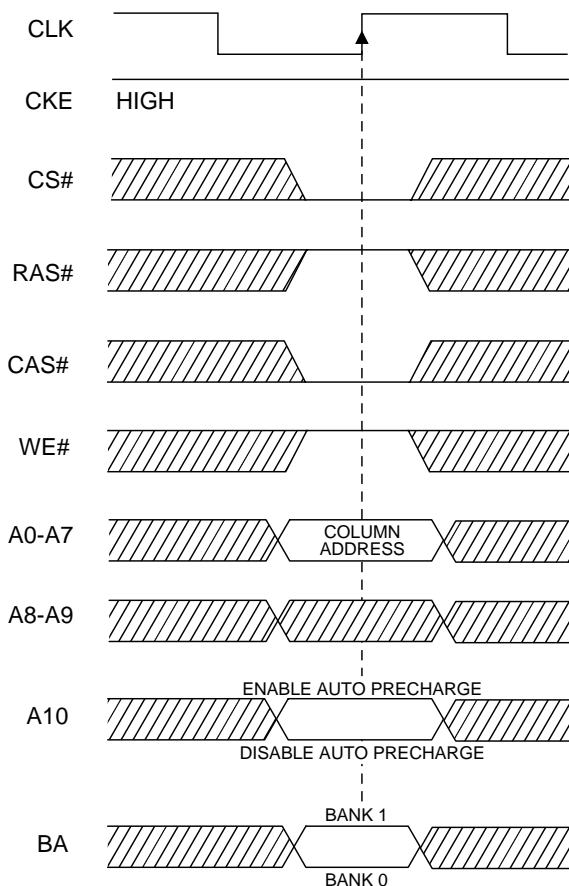


Figure 5
READ COMMAND

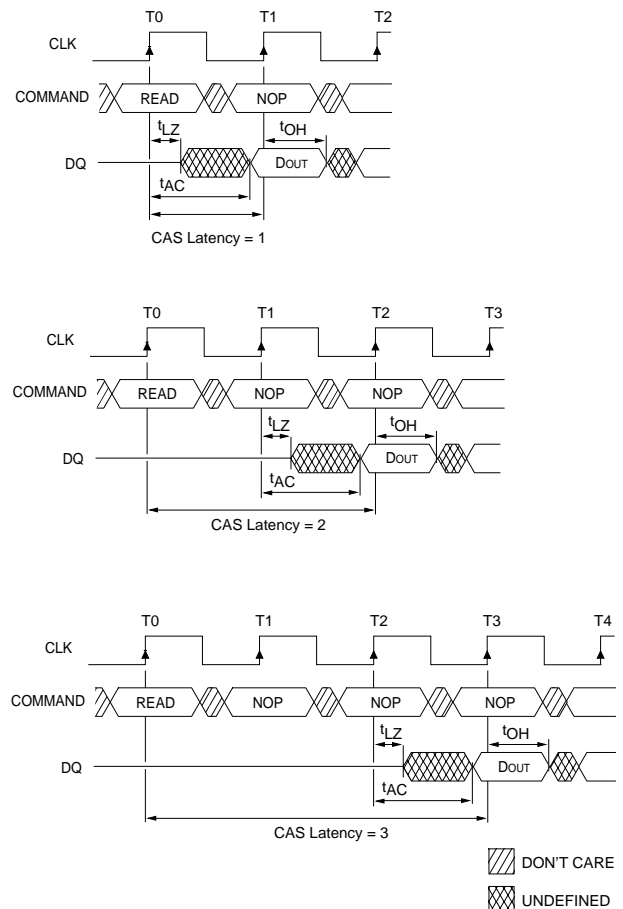
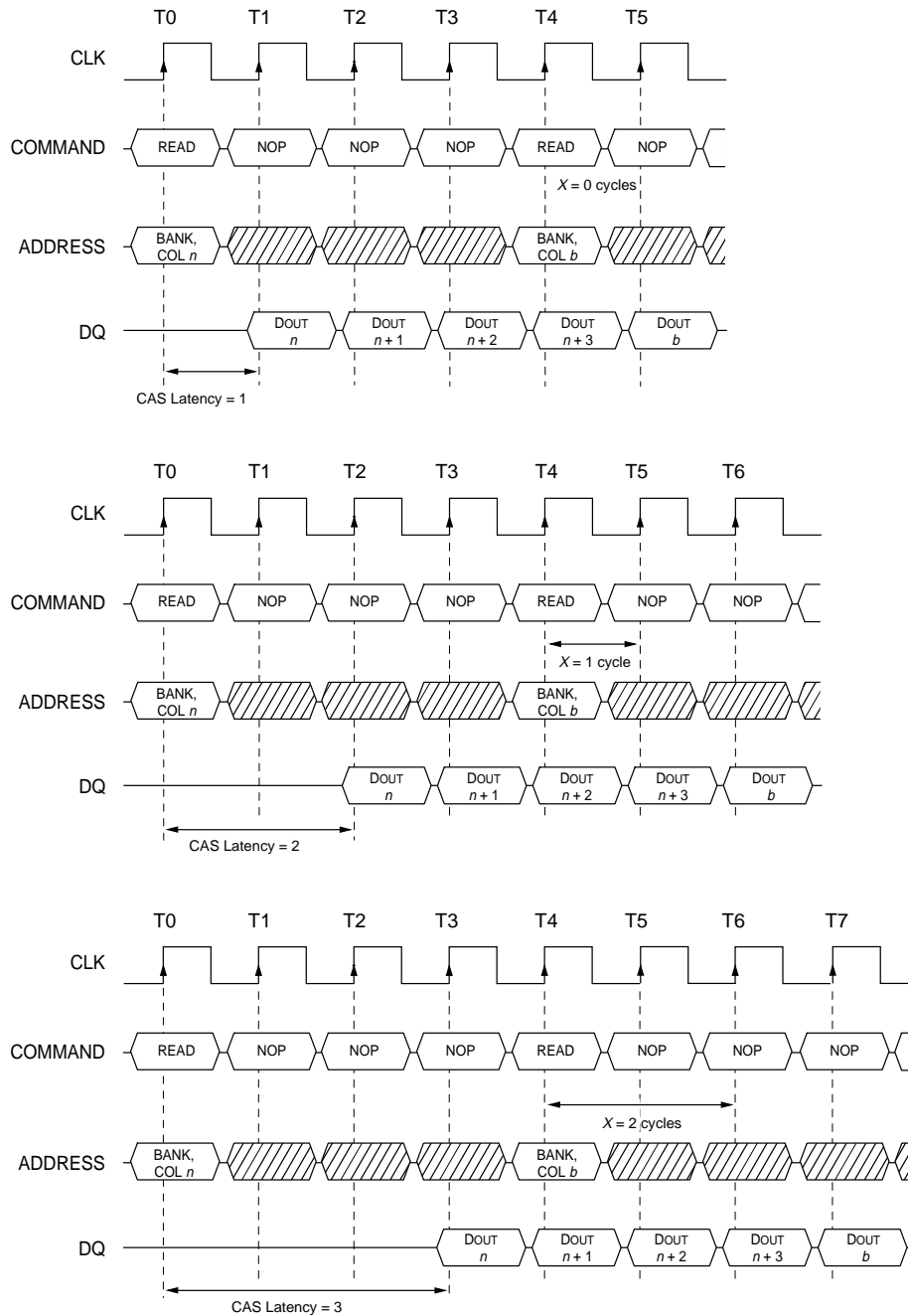


Figure 6
CAS LATENCY

latencies of one, two and three; data element $n + 3$ is either the last of a burst of four or the last desired of a longer burst. The 1 Meg x 16 SDRAM uses a pipelined architecture and therefore does not require the $2n$ rule associated with a

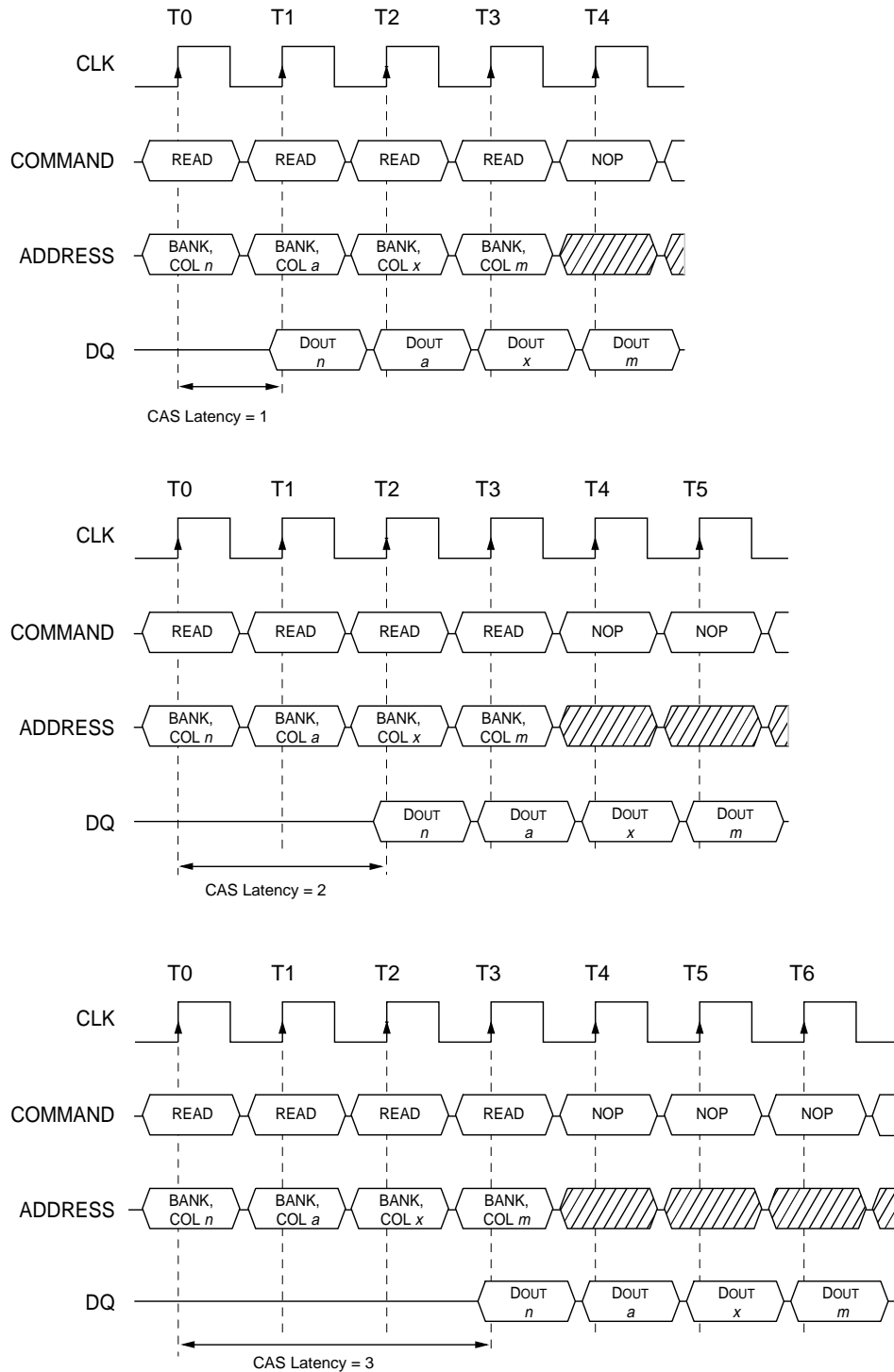
prefetch architecture. A READ command can be initiated on any clock cycle following a previous READ command. Full-speed, random read accesses within a page can be performed as shown in Figure 8.



NOTE: Each READ command may be to either bank. DQM is LOW.

DON'T CARE

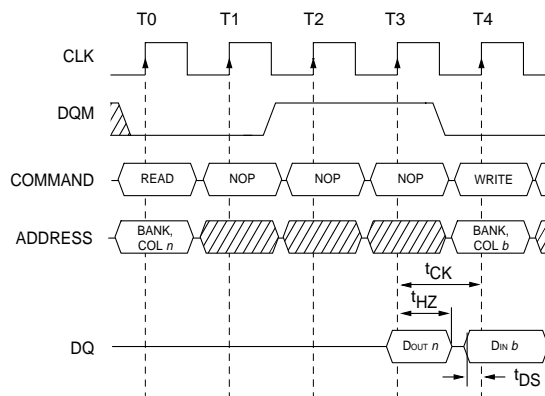
Figure 7
CONSECUTIVE READ BURSTS



NOTE: Each READ command may be to either bank. DQM is LOW. / / DON'T CARE

Figure 8
RANDOM READ ACCESSES WITHIN A PAGE

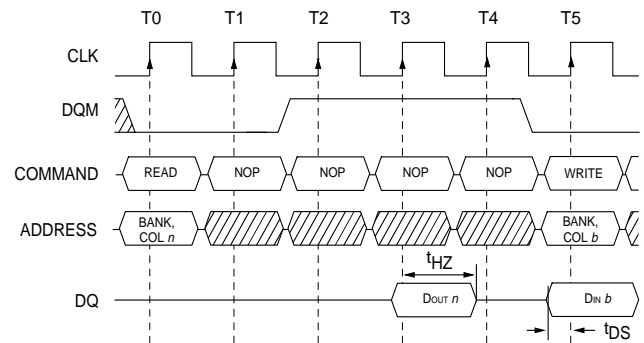
Data from any READ burst may be truncated with a subsequent WRITE command, and data from a fixed-length READ burst may be immediately followed by data from a subsequent WRITE command (subject to bus turnaround limitations). The WRITE burst may be initiated on the clock edge immediately following the last (or last desired) data element from the READ burst, provided that I/O contention can be avoided. In a given system design, there may be the possibility that the device driving the input data would go Low-Z before the SDRAM DQs go High-Z. In this case, at least a single-cycle delay should occur between the last read data and the WRITE command.



NOTE: A CAS latency of three is used for illustration. The READ command may be to any bank, and the WRITE command may be to any bank. If a CAS latency of one is used, then DQM is not required.

Figure 9
READ TO WRITE

The DQM input is used to avoid I/O contention as shown in Figures 9 and 10. The DQM signal must be asserted (HIGH) at least two clocks (DQM latency is two clocks for output buffers) prior to the WRITE command to suppress data-out from the READ. Once the WRITE command is registered, the DQs will go High-Z (or remain High-Z) regardless of the state of the DQM signal. The DQM signal must be de-asserted (DQM latency is zero clocks for input buffers) prior to the WRITE command to ensure that the written data is not masked. Figure 9 shows the case where the clock frequency allows for bus contention to be avoided without adding a NOP cycle, and Figure 10 shows the case where the additional NOP is needed.



NOTE: A CAS latency of three is used for illustration. The READ command may be to any bank, and the WRITE command may be to any bank.

 DON'T CARE

Figure 10
READ TO WRITE WITH EXTRA CLOCK CYCLE

A fixed-length READ burst may be followed by, or truncated with, a PRECHARGE command to the same bank (provided that AUTO PRECHARGE was not activated) and a full-page burst may be truncated with a PRECHARGE command to the same bank. The PRECHARGE command should be issued x cycles before the clock edge at which the last desired data element is valid, where x equals the CAS latency minus one. This is shown in Figure 11 for each possible CAS latency; data element $n + 3$ is either the last of

a burst of four or the last desired of a longer burst. Following the PRECHARGE command, a subsequent command to the same bank cannot be issued until t_{RP} is met. Note that part of the row precharge time is hidden during the access of the last data element(s).

In the case of a fixed-length burst being executed to completion, a PRECHARGE command issued at the optimum time (as described above) provides the same operation that would result from the same fixed-length burst

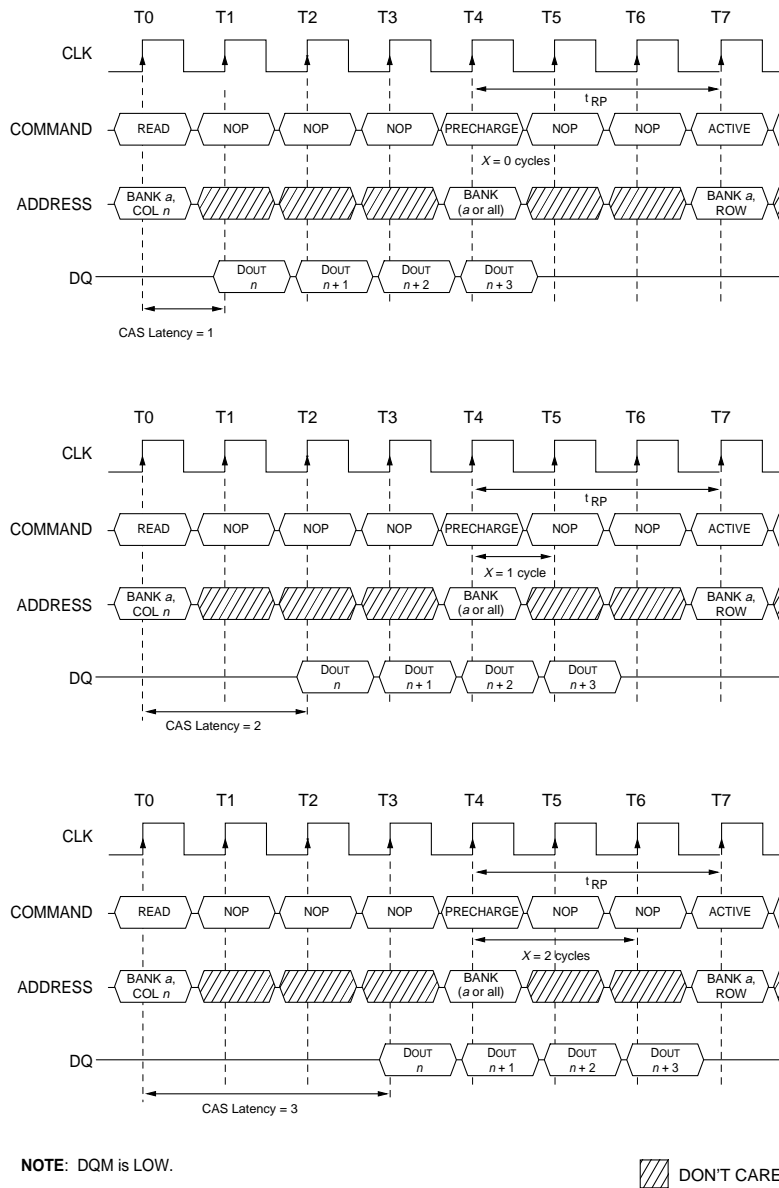


Figure 11
READ TO PRECHARGE

with AUTO PRECHARGE. The disadvantage of the PRECHARGE command is that it requires that the command and address buses be available at the appropriate time to issue the command, but the advantage of the PRECHARGE command is that it can be used to truncate fixed-length or full-page bursts.

Full-page READ bursts can be truncated with the BURST TERMINATE command, and fixed-length READ bursts

may be truncated with a BURST TERMINATE command, provided that AUTO PRECHARGE was not activated. The BURST TERMINATE command should be issued x cycles before the clock edge at which the last desired data element is valid, where x equals the CAS latency minus one. This is shown in Figure 12 for each possible CAS latency; data element $n + 3$ is the last desired data element of a longer burst.

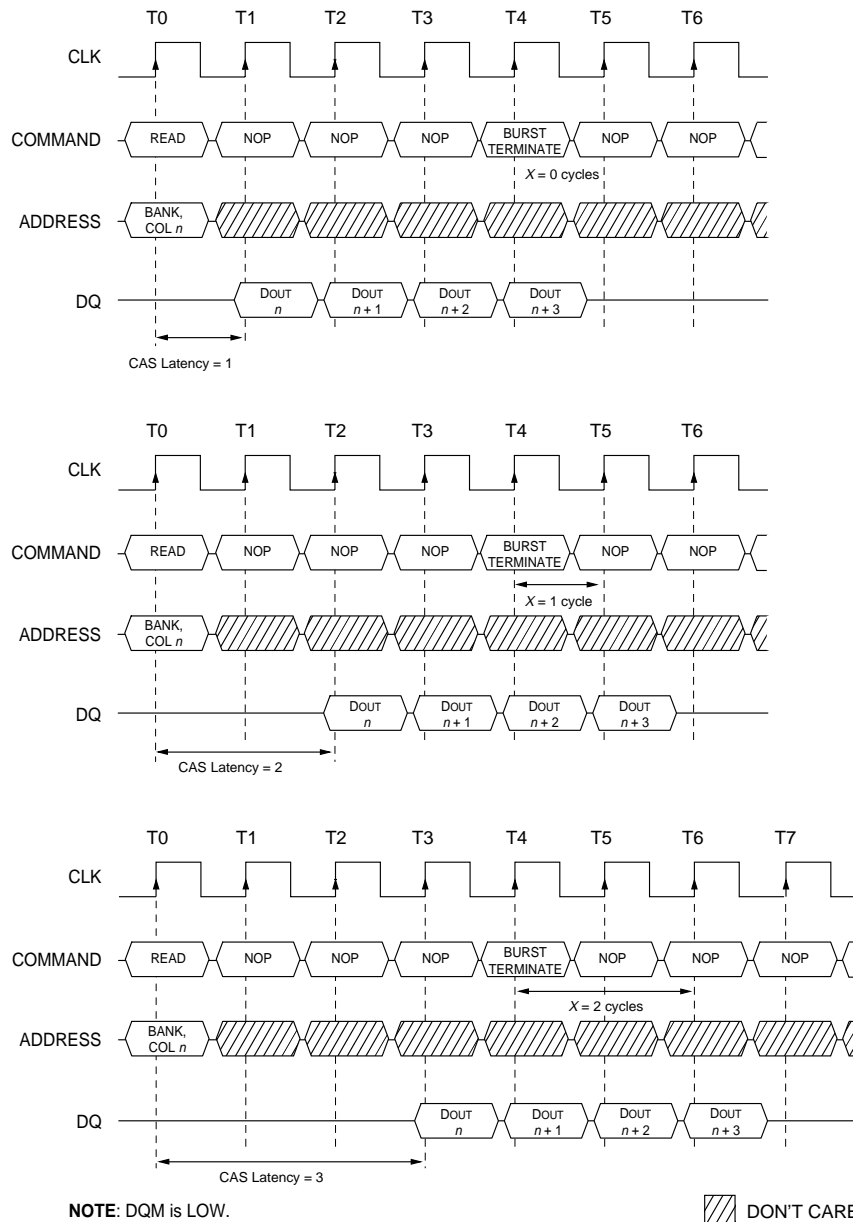


Figure 12
TERMINATING A READ BURST

WRITES

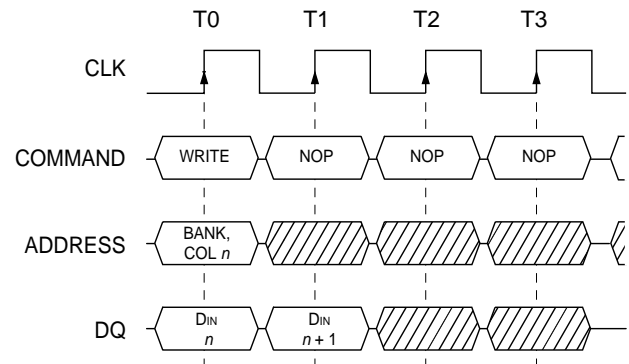
WRITE bursts are initiated with a WRITE command, as shown in Figure 13.

The starting column and bank addresses are provided with the WRITE command and AUTO PRECHARGE is either enabled or disabled for that access. If AUTO PRECHARGE is enabled, the row being accessed is precharged at the completion of the burst. For the generic WRITE commands used in the following illustrations, AUTO PRECHARGE is disabled.

During WRITE bursts, the first valid data-in element will be registered coincident with the WRITE command. Subsequent data elements will be registered on each successive positive clock edge. Upon completion of a fixed-length burst, assuming no other commands have been initiated, the DQs will remain High-Z, and any additional input data will be ignored (see Figure 14). A full-page burst will continue until terminated (at the end of the page it will wrap to column 0 and continue).

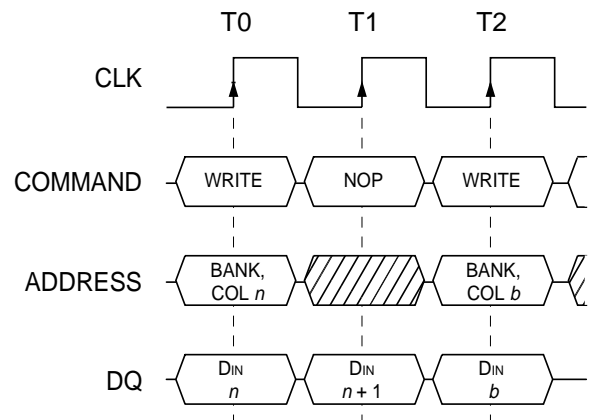
Data for any WRITE burst may be truncated with a subsequent WRITE command, and data for a fixed-length WRITE burst may be immediately followed by data for a subsequent WRITE command. The new WRITE command can be issued on any clock following the previous WRITE command, and the data provided coincident with the new

command applies to the new command. An example is shown in Figure 15. Data $n + 1$ is either the last of a burst of two, or the last desired of a longer burst. The 1 Meg x 16 SDRAM uses a pipelined architecture and therefore does not require the $2n$ rule associated with a prefetch architecture. A WRITE command can be initiated on any clock cycle following a previous WRITE command. Full-speed, ran-



NOTE: Burst length = 2. DQM is LOW.

Figure 14
WRITE BURST



NOTE: DQM is LOW. Each WRITE command may be to any bank.

 DON'T CARE

Figure 15
WRITE TO WRITE

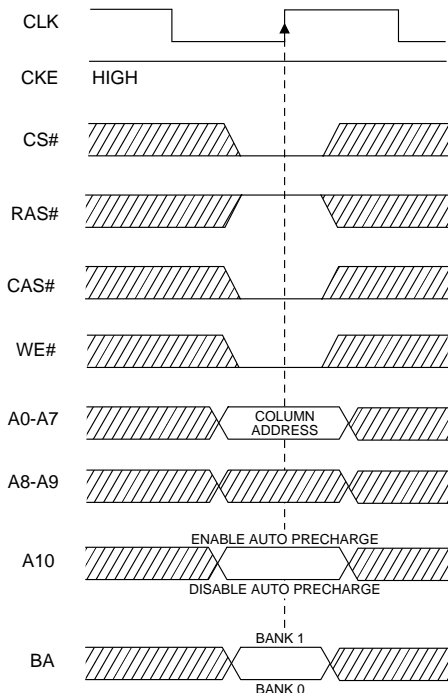


Figure 13
WRITE COMMAND

dom write accesses within a page can be performed as shown in Figure 16.

Data for any WRITE burst may be truncated with a subsequent READ command, and data for a fixed-length WRITE burst may be immediately followed by a subsequent READ command. Once the READ command is registered, the data inputs will be ignored, and WRITES will not be executed. An example is shown in Figure 17. Data $n + 1$ is either the last of a burst of two, or the last desired of a longer burst.

Data for a fixed-length WRITE burst may be followed by, or truncated with, a PRECHARGE command to the same bank (provided that AUTO PRECHARGE was not acti-

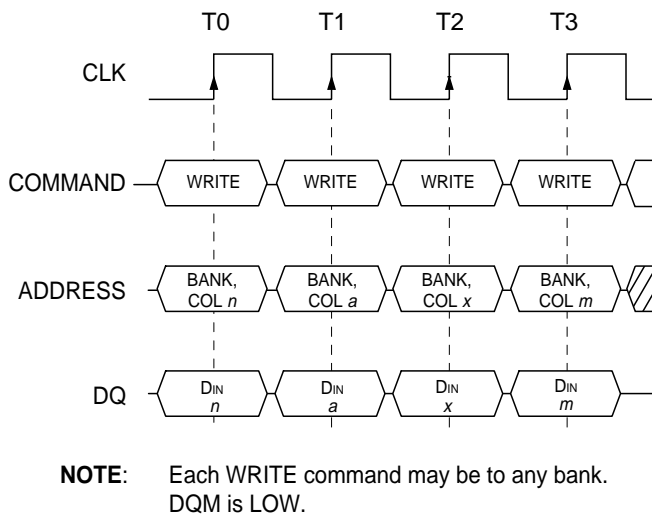


Figure 16
RANDOM WRITE CYCLES WITHIN A PAGE

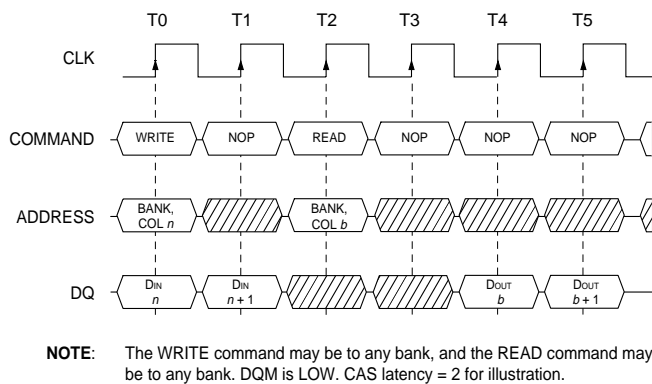


Figure 17
WRITE TO READ

ated), and a full-page WRITE burst may be truncated with a PRECHARGE command to the same bank. The PRECHARGE command should be issued t_{WR} after the clock edge at which the last desired input data element is registered. In addition, when truncating a WRITE burst, the DQM signal must be used to mask input data for the clock edge prior to, and the clock edge coincident with, the PRECHARGE command. An example is shown in Figure 18. Data $n + 1$ is either the last of a burst of two, or the last desired of a longer burst. Following the PRECHARGE command, a subsequent command to the same bank cannot be issued until t_{RP} is met.

In the case of a fixed-length burst being executed to completion, a PRECHARGE command issued at the optimum time (as described above) provides the same operation that would result from the same fixed-length burst with AUTO PRECHARGE. The disadvantage of the PRECHARGE command is that it requires that the command and address buses be available at the appropriate time to issue the command, but the advantage of the PRECHARGE command is that it can be used to truncate fixed-length or full-page bursts.

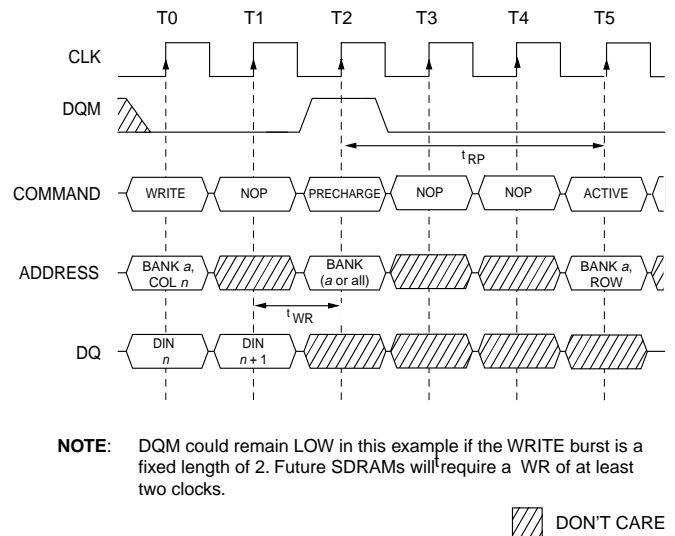


Figure 18
WRITE TO PRECHARGE

Fixed-length or full-page WRITE bursts can be truncated with the BURST TERMINATE command. When truncating a WRITE burst, the input data applied coincident with the BURST TERMINATE command will be ignored. The last data written (provided that DQM is LOW at that time) will be the input data applied one clock previous to the BURST TERMINATE command. This is shown in Figure 19, where data *n* is the last desired data element of a longer burst.

PRECHARGE

The PRECHARGE command is used to deactivate the open row in a particular bank or the open row in all banks. The bank(s) will be available for a subsequent row access some specified time (^tRP) after the PRECHARGE command is issued. Input A10 determines whether one or all banks are to be precharged, and in the case where only one bank is to be precharged, input BA selects the bank. When all banks are to be precharged, input BA is treated as “Don’t Care.” Once a bank has been precharged, it is in the Idle state and must be activated prior to any READ or WRITE commands being issued to that bank.

POWER-DOWN

POWER-DOWN occurs if CKE is registered LOW coincident with a NOP or COMMAND INHIBIT, when no accesses are in progress. If POWER-DOWN occurs when all banks are idle, this mode is referred to as precharge power-down; if power-down occurs when there is a row active in either bank, this mode is referred to as active power-down. Entering power-down deactivates the input and output buffers, excluding CKE, for maximum power savings while in standby. The device may not remain in the power-down state longer than the refresh period (64ms) since no refresh operations are performed in this mode.

The power-down state is exited by registering a NOP or COMMAND INHIBIT and CKE HIGH at the desired clock edge (meeting ^tCKS).

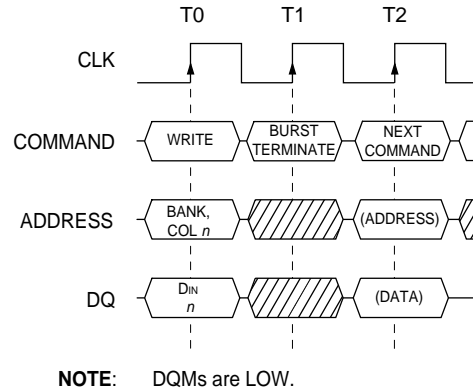


Figure 19
TERMINATING A WRITE BURST

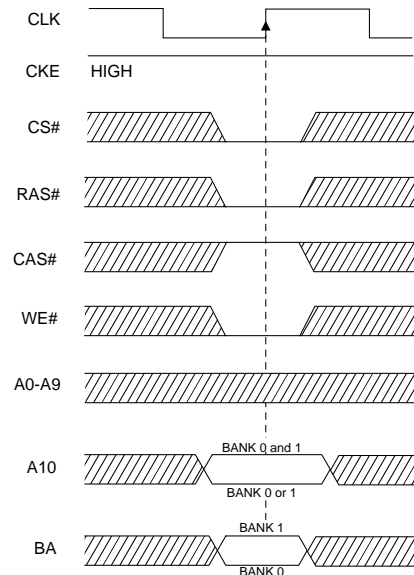


Figure 20
PRECHARGE COMMAND

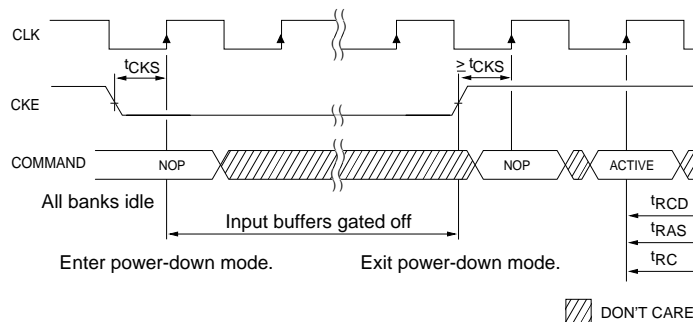
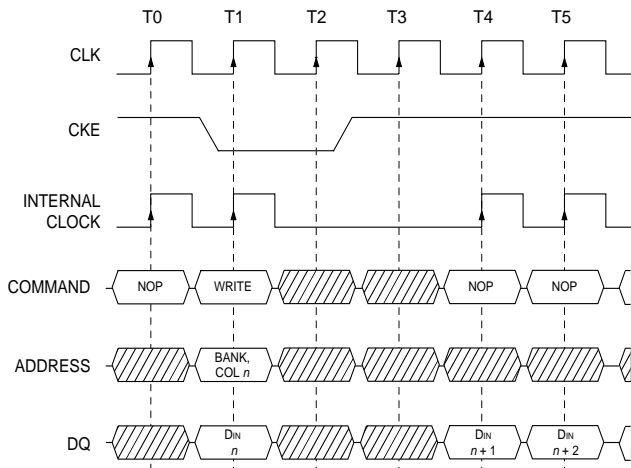


Figure 21
POWER-DOWN

CLOCK SUSPEND

The clock suspend mode occurs when a column access/burst is in progress and CKE is registered LOW. In the clock suspend mode, the internal clock is deactivated, “freezing” the synchronous logic.

For each positive clock edge on which CKE is sampled LOW, the next internal positive clock edge is suspended. Any command or data present on the input pins at the time of a suspended internal clock edge are ignored; any data present on the DQ pins will remain driven; and burst counters are not incremented as long as the clock is suspended (see examples in Figures 22 and 23).



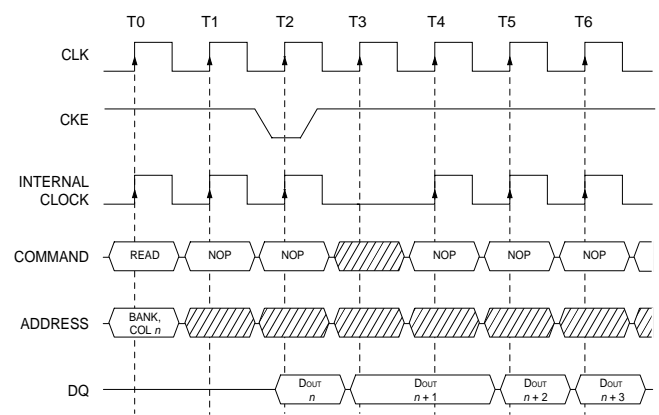
NOTE: For this example, burst length = 4 or greater, and DQM is LOW.

Figure 22
CLOCK SUSPEND DURING WRITE BURST

Clock suspend mode is exited by registering CKE HIGH; the internal clock and related operation will resume on the subsequent positive clock edge.

BURST READ/SINGLE WRITE

The burst read/single write mode is entered by programming the write burst mode bit (M9) in the Mode Register to a logic 1. In this mode, all WRITE commands result in the access of a single column location (burst of 1) regardless of the programmed burst length. READ commands access columns according to the programmed burst length and sequence, just as in the normal mode of operation (M9 = 0).



NOTE: For this example, CAS latency = 2, burst length = 4 or greater, and DQM is LOW.

DON'T CARE

Figure 23
CLOCK SUSPEND DURING READ BURST

TRUTH TABLE 2 – CKE

(Notes: 1-4)

| CKE _{n-1} | CKE _n | CURRENT STATE | COMMAND _n | ACTION _n | NOTES |
|--------------------|------------------|--------------------|------------------------|------------------------|-------|
| L | L | Power-Down | X | Maintain Power-Down | |
| | | Self Refresh | X | Maintain Self Refresh | |
| | | Clock Suspend | X | Maintain Clock Suspend | |
| L | H | Power-Down | COMMAND INHIBIT or NOP | Exit Power-Down | 5 |
| | | Self Refresh | COMMAND INHIBIT or NOP | Exit Self Refresh | 6 |
| | | Clock Suspend | X | Exit Clock Suspend | 7 |
| H | L | All Banks Idle | COMMAND INHIBIT or NOP | Power-Down Entry | |
| | | All Banks Idle | AUTO REFRESH | Self Refresh Entry | |
| | | Reading or Writing | VALID | Clock Suspend Entry | |
| H | H | | See Truth Table 3 | | |

- NOTE:**
1. CKE_n is the logic state of CKE at clock edge *n*; CKE_{n-1} was the state of CKE at the previous clock edge.
 2. Current state is the state of the SDRAM immediately prior to clock edge *n*.
 3. COMMAND_n is the command registered at clock edge *n* and ACTION_n is a result of COMMAND_n.
 4. All states and sequences not shown are illegal or reserved.
 5. Exiting power-down at clock edge *n* will put the device in the All Banks Idle state in time for clock edge *n + 1* (provided that ^tCKS is met).
 6. Exiting SELF REFRESH at clock edge *n* will put the device in the All Banks Idle state once ^tXSR is met. COMMAND INHIBIT or NOP commands should be issued on any clock edges occurring during the ^tXSR period. A minimum of two NOP commands must be provided during ^tXSR period.
 7. After exiting clock suspend at clock edge *n*, the device will resume operation and recognize the next command at clock edge *n + 1*.

TRUTH TABLE 3 – Current State Bank n - Command to Bank n

(Notes: 1-6; notes appear below and on next page)

| CURRENT STATE | CS# | RAS# | CAS# | WE# | COMMAND/ACTION | NOTES |
|------------------------------------|-----|------|------|-----|---|-------|
| Any | H | X | X | X | COMMAND INHIBIT (NOP/continue previous operation) | |
| | L | H | H | H | NO OPERATION (NOP/continue previous operation) | |
| Idle | L | L | H | H | ACTIVE (select and activate row) | |
| | L | L | L | H | AUTO REFRESH | 7 |
| | L | L | L | L | LOAD MODE REGISTER | 7 |
| Row Active | L | H | L | H | READ (select column and start READ burst) | 10 |
| | L | H | L | L | WRITE (select column and start WRITE burst) | 10 |
| | L | L | H | L | PRECHARGE (deactivate row in bank or banks) | 8 |
| Read (Auto-Precharge Disabled) | L | H | L | H | READ (select column and start new READ burst) | 10 |
| | L | H | L | L | WRITE (select column and start WRITE burst) | 10 |
| | L | L | H | L | PRECHARGE (truncate READ burst, start PRECHARGE) | 8 |
| | L | H | H | L | BURST TERMINATE | 9 |
| Write (Auto-Precharge Disabled) | L | H | L | H | READ (select column and start READ burst) | 10 |
| | L | H | L | L | WRITE (select column and start new WRITE burst) | 10 |
| | L | L | H | L | PRECHARGE (truncate WRITE burst, start PRECHARGE) | 8 |
| | L | H | H | L | BURST TERMINATE | 9 |

NOTES:

- This table applies when CKE_{n-1} was HIGH and CKE_n is HIGH (see Truth Table 2) and after t_{XSR} has been met (if the previous state was self refresh).
- This table is bank-specific, except where noted; i.e., the current state is for a specific bank and the commands shown are those allowed to be issued to that bank when in that state. Exceptions are covered in the notes below.
- Current state definitions:
 - Idle: The bank has been precharged and t_{RP} has been met.
 - Row Active: A row in the bank has been activated and t_{RCD} has been met. No data bursts/accesses and no register accesses are in progress.
 - Read: A READ burst has been initiated, with AUTO PRECHARGE disabled, and has not yet terminated or been terminated.
 - Write: A WRITE burst has been initiated, with AUTO PRECHARGE disabled, and has not yet terminated or been terminated.
- The following states must not be interrupted by a command issued to the same bank. COMMAND INHIBIT or NOP commands, or allowable commands to the other bank, should be issued on any clock edge occurring during these states. Allowable commands to the other bank are determined by its current state and Truth Table 3, and according to Truth Table 4.
 - Precharging: Starts with registration of a PRECHARGE command and ends when t_{RP} is met. Once t_{RP} is met, the bank will be in the Idle state.
 - Row Activating: Starts with registration of an ACTIVE command and ends when t_{RCD} is met. Once t_{RCD} is met, the bank will be in the Row Active state.
 - Read w/Auto-Precharge Enabled: Starts with registration of a READ command with AUTO PRECHARGE enabled and ends when t_{RP} has been met. Once t_{RP} is met, the bank will be in the Idle state.
 - Write w/Auto-Precharge Enabled: Starts with registration of a WRITE command with AUTO PRECHARGE enabled and ends when t_{RP} has been met. Once t_{RP} is met, the bank will be in the Idle state.

NOTES (continued):

5. The following states must not be interrupted by any executable command; COMMAND INHIBIT or NOP commands must be applied on each positive clock edge during these states.
 - Refreshing: Starts with registration of an AUTO REFRESH command and ends when t^1RC is met. Once t^1RC is met, the SDRAM will be in the All Banks Idle state.
 - Accessing Mode
 - Register: Starts with registration of a LOAD MODE REGISTER command and ends when t^1MTC has been met. Once t^1MTC is met, the SDRAM will be in the All Banks Idle state.
 - Precharging All: Starts with registration of a PRECHARGE ALL command and ends when t^1RP is met. Once t^1RP is met, all banks will be in the Idle state.
6. All states and sequences not shown are illegal or reserved.
7. Not bank-specific; requires that all banks are idle.
8. May or may not be bank-specific; if all banks are to be precharged, all must be in a valid state for precharging.
9. Not bank-specific; BURST TERMINATE affects the most recent READ or WRITE burst, regardless of bank.
10. READs or WRITEs listed in the Command/Action column include READs or WRITEs with AUTO PRECHARGE enabled and READs or WRITEs with AUTO PRECHARGE disabled.

TRUTH TABLE 4 – Current State Bank n - Command to Bank m

(Notes: 1-6, 8; notes appear below and on next page)

| CURRENT STATE | CS# | RAS# | CAS# | WE# | COMMAND/ACTION | NOTES |
|---|-----|------|------|-----|---|-------|
| Any | H | X | X | X | COMMAND INHIBIT (NOP/continue previous operation) | |
| | L | H | H | H | NO OPERATION (NOP/continue previous operation) | |
| Idle | X | X | X | X | Any command otherwise allowed to bank m | |
| Row Activating, Active or Precharging | L | L | H | H | ACTIVE (select and activate row) | |
| | L | H | L | H | READ (select column and start READ burst) | 7 |
| | L | H | L | L | WRITE (select column and start WRITE burst) | 7 |
| | L | L | H | L | PRECHARGE (READ burst, start PRECHARGE) | |
| Read (Auto- Precharge Disabled) | L | L | H | H | ACTIVE (select and activate row) | |
| | L | H | L | H | READ (select column and start new READ burst) | 7 |
| | L | H | L | L | WRITE (select column and start WRITE burst) | 7 |
| | L | L | H | L | PRECHARGE (READ burst, start PRECHARGE) | |
| Write (Auto- Precharge Disabled) | L | L | H | H | ACTIVE (select and activate row) | |
| | L | H | L | H | READ (select column and start READ burst) | 7 |
| | L | H | L | L | WRITE (select column and start new WRITE burst) | 7 |
| | L | L | H | L | PRECHARGE (WRITE burst, start PRECHARGE) | |
| Read (With Auto- Precharge) | L | L | H | H | ACTIVE (select and activate row) | |
| | L | H | L | H | READ (select column and start new READ burst) | 7 |
| | L | H | L | L | WRITE (select column and start WRITE burst) | 7 |
| | L | L | H | L | PRECHARGE (READ burst, start PRECHARGE) | |
| Write (With Auto- Precharge) | L | L | H | H | ACTIVE (select and activate row) | |
| | L | H | L | H | READ (select column and start READ burst) | 7 |
| | L | H | L | L | WRITE (select column and start new WRITE burst) | 7 |
| | L | L | H | L | PRECHARGE (WRITE burst, start PRECHARGE) | |

NOTES:

1. This table applies when CKE_{n-1} was HIGH and CKE_n is HIGH (see Truth Table 2) and after t_{XSR} has been met (if the previous state was self refresh).
2. This table describes alternate bank operation, except where noted, i.e., the current state is for bank n and the commands shown are those allowed to be issued to bank m (assuming that bank m is in such a state that the given command is allowable). Exceptions are covered in the notes below.

NOTES (continued):

3. Current state definitions:
 - Idle: The bank has been precharged and t_{RP} has been met.
 - Row Active: A row in the bank has been activated and t_{RCD} has been met. No data bursts/ accesses and no register accesses are in progress.
 - Read: A READ burst has been initiated, with AUTO PRECHARGE disabled, and has not yet terminated or been terminated.
 - Write: A WRITE burst has been initiated, with AUTO PRECHARGE disabled, and has not yet terminated or been terminated.
 - Read w/Auto-Precharge Enabled: Starts with registration of a READ command with AUTO PRECHARGE enabled and ends when t_{RP} has been met. Once t_{RP} is met, the bank will be in the Idle state.
 - Write w/Auto-Precharge Enabled: Starts with registration of a WRITE command with AUTO PRECHARGE enabled and ends when t_{RP} has been met. Once t_{RP} is met, the bank will be in the Idle state.
4. AUTO REFRESH, LOAD MODE REGISTER and PRECHARGE ALL commands may only be issued when all banks are idle.
5. A BURST TERMINATE command cannot be issued to another bank; it applies to the bank represented by the current state only.
6. All states and sequences not shown are illegal or reserved.
7. READs or WRITEs listed in the Command/Action column include READs or WRITEs with AUTO PRECHARGE enabled and READs or WRITEs with AUTO PRECHARGE disabled.
8. Truth Tables 3 and 4 differentiate from initial x4 and x8 designs; interrupting an AUTO PRECHARGE is now allowed.

ABSOLUTE MAXIMUM RATINGS*

| | |
|---|-----------------|
| Voltage on Vcc/VccQ Supply | |
| Relative to Vss | -1V to +4.6V |
| Voltage on Inputs, NC or I/O Pins | |
| Relative to Vss | -1V to +4.6V |
| Operating Temperature, T _A (ambient) | 0°C to +70°C |
| Storage Temperature (plastic) | -55°C to +150°C |
| Power Dissipation | 1W |
| Short Circuit Output Current | 50mA |

*Stresses greater than those listed under “Absolute Maximum Ratings” may cause permanent damage to the device. This is a stress rating only and functional operation of the device at these or any other conditions above those indicated in the operational sections of this specification is not implied. Exposure to absolute maximum rating conditions for extended periods may affect reliability.

ELECTRICAL CHARACTERISTICS AND RECOMMENDED DC OPERATING CONDITIONS

(Notes: 1, 6) (0°C ≤ T_A ≤ 70°C; Vcc/VccQ = +3.3V ±0.3V)

| PARAMETER/CONDITION | SYMBOL | MIN | MAX | UNITS | NOTES |
|---|-----------------|------|----------|-------|-------|
| Supply Voltage | Vcc/VccQ | 3.0 | 3.6 | V | |
| Input High (Logic 1) Voltage, all inputs | V _{IH} | 2.0 | Vcc +0.3 | V | |
| Input Low (Logic 0) Voltage, all inputs | V _{IL} | -0.5 | 0.8 | V | |
| INPUT LEAKAGE CURRENT Any input 0V ≤ V _{IN} ≤ Vcc (All other pins not under test = 0V) | I _I | -5 | 5 | μA | |
| OUTPUT LEAKAGE CURRENT (DQs are disabled; 0V ≤ V _{OUT} ≤ VccQ) | I _{OZ} | -10 | 10 | μA | |
| OUTPUT LEVELS Output High Voltage (I _{OUT} = -4mA) | V _{OH} | 2.4 | | V | |
| Output Low Voltage (I _{OUT} = 4mA) | V _{OL} | | 0.4 | V | |

I_{CC} SPECIFICATIONS AND CONDITIONS

(Notes: 1, 6, 13) (0°C ≤ T_A ≤ 70°C; Vcc/VccQ = +3.3V ±0.3V)

| PARAMETER/CONDITION | SYMBOL | MAX | | | UNITS | NOTES |
|--|------------------|-----|-----|-----|-------|--------------|
| | | -8A | -10 | -12 | | |
| OPERATING CURRENT: Active Mode, Burst = 2, READ or WRITE, t _{RC} ≥ t _{RC} (MIN), CAS latency = 3 | I _{CC1} | 135 | 130 | 115 | mA | 3, 18, 19 |
| STANDBY CURRENT: Power-Down Mode, t _{CK} = 15ns (10ns for -8), CKE ≤ V _{IL} (MAX), All banks idle | I _{CC2} | 3 | 2 | 3 | mA | |
| STANDBY CURRENT: CS# ≥ V _{IH} (MIN), CKE ≥ V _{IH} (MIN), t _{CK} = 15ns (10ns for -8), All banks idle | I _{CC3} | 45 | 30 | 30 | mA | 12, 19 |
| STANDBY CURRENT: CS# ≥ V _{IH} (MIN), CKE ≥ V _{IH} (MIN), t _{CK} = 15ns (10ns for -8), All banks active after t _{RCD} met, No accesses in progress | I _{CC4} | 45 | 40 | 40 | mA | 12, 19 |
| OPERATING CURRENT: Burst Mode, Continuous burst, READ or WRITE, t _{CK} = 15ns (10ns for -8), All banks active, Addresses transition once per clock cycle, CAS latency = 3 | I _{CC5} | 195 | 130 | 110 | mA | 3, 18, 19 |
| AUTO REFRESH CURRENT: t _{RC} = 15.625μS, CL = 3 | I _{CC6} | 45 | 30 | 30 | mA | 19 |
| SELF REFRESH CURRENT: CKE ≤ 0.2V | I _{CC7} | 1 | 1 | 1 | mA | 4 |

CAPACITANCE

| PARAMETER | SYMBOL | MAX | UNITS | NOTES |
|--|-----------------|-----|-------|-------|
| Input Capacitance: CLK | C _{I1} | 5.0 | pF | 2 |
| Input Capacitance: All other input-only pins | C _{I2} | 5.0 | pF | 2 |
| Input/Output Capacitance: DQs | C _{I0} | 6.5 | pF | 2 |

ELECTRICAL CHARACTERISTICS AND RECOMMENDED AC OPERATING CONDITIONS

(Notes: 5, 6, 8, 9, 11) (0°C ≤ T_A ≤ +70°C)

| AC CHARACTERISTICS | | | -8A | | -10 | | -12 | | UNITS | NOTES |
|--|--------|------------------|-----|------|-----|------|-----|------|-----------------|-------|
| PARAMETER | | SYM | MIN | MAX | MIN | MAX | MIN | MAX | | |
| Access time from CLK (pos. edge) | CL = 3 | t _{AC} | | 6 | | 7.5 | | 9 | ns | |
| | CL = 2 | t _{AC} | | 9 | | 9 | | 9 | ns | |
| | CL = 1 | t _{AC} | | 27 | | 27 | | 27 | ns | |
| Address hold time | | t _{AH} | 1 | | 1 | | 1 | | ns | |
| Address setup time | | t _{AS} | 2 | | 3 | | 3 | | ns | |
| CLK high level width | | t _{CH} | 3 | | 3.5 | | 3.5 | | ns | |
| CLK low level width | | t _{CL} | 3 | | 3.5 | | 3.5 | | ns | |
| Clock cycle time | CL = 3 | t _{CK} | 8 | | 10 | | 12 | | ns | |
| | CL = 2 | t _{CK} | 13 | | 15 | | 15 | | ns | |
| | CL = 1 | t _{CK} | 30 | | 30 | | 30 | | ns | |
| CKE hold time | | t _{CKH} | 1 | | 1 | | 1 | | ns | |
| CKE setup time | | t _{CKS} | 2 | | 3 | | 3 | | ns | |
| CS#, RAS#, CAS#, WE#, DQM hold time | | t _{CMH} | 1 | | 1 | | 1 | | ns | |
| CS#, RAS#, CAS#, WE#, DQM setup time | | t _{CMS} | 2 | | 3 | | 3 | | ns | |
| Data-in hold time | | t _{DH} | 1 | | 1 | | 1 | | ns | |
| Data-in setup time | | t _{DS} | 2 | | 3 | | 3 | | ns | |
| Data-out high-impedance time | CL = 3 | t _{HZ} | | 6 | | 8 | | 9 | ns | 10 |
| | CL = 2 | t _{HZ} | | 7 | | 10 | | 10 | ns | 10 |
| | CL = 1 | t _{HZ} | | 15 | | 15 | | 15 | ns | 10 |
| Data-out low-impedance time | | t _{LZ} | 1 | | 2 | | 2 | | ns | |
| Data-out hold time | | t _{OH} | 3 | | 3 | | 3 | | ns | |
| ACTIVE to PRECHARGE command | | t _{RAS} | 50 | 120K | 60 | 120K | 72 | 120K | ns | |
| AUTO REFRESH, ACTIVE command period | | t _{RC} | 80 | | 90 | | 105 | | ns | |
| ACTIVE to READ or WRITE delay | | t _{RCD} | 30 | | 30 | | 30 | | ns | |
| Refresh period - 2,048 or 4,096 rows | | t _{REF} | | 64 | | 64 | | 64 | ms | |
| PRECHARGE command period | | t _{RP} | 30 | | 30 | | 36 | | ns | 21 |
| ACTIVE bank A to ACTIVE bank B command | | t _{RRD} | 20 | | 20 | | 20 | | ns | |
| Transition time | | t _T | 0.3 | 1.2 | 0.3 | 10 | 1 | 20 | ns | 7 |
| WRITE recovery time | | t _{WR} | 1 | | 1 | | 1 | | t _{CK} | 21 |
| Exit SELF REFRESH to ACTIVE command | | t _{XSR} | 80 | | 96 | | 105 | | ns | 20 |

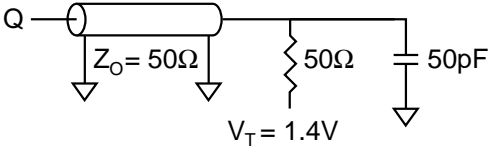
AC FUNCTIONAL CHARACTERISTICS

(Notes: 5, 6, 7, 8, 9, 11) ($0^{\circ}\text{C} \leq T_A \leq +70^{\circ}\text{C}$)

| PARAMETER | | SYM | -8A | -10 | -12 | UNITS | NOTES |
|---|--------|-------------------|-----|-----|-----|-----------------|--------|
| READ/WRITE command to READ/WRITE command | | t_{CCD} | 1 | 1 | 1 | t_{CK} | 17 |
| CKE to clock disable or power-down entry mode | | t_{CKED} | 1 | 1 | 1 | t_{CK} | 14 |
| CKE to clock enable or power-down exit setup mode | | t_{PED} | 1 | 1 | 1 | t_{CK} | 14 |
| DQM to input data delay | | t_{DQD} | 0 | 0 | 0 | t_{CK} | 17 |
| DQM to data mask during WRITES | | t_{DQM} | 0 | 0 | 0 | t_{CK} | 17 |
| DQM to data high-impedance during READs | | t_{DQZ} | 2 | 2 | 2 | t_{CK} | 17 |
| WRITE command to input data delay | | t_{DWD} | 0 | 0 | 0 | t_{CK} | 17 |
| Data-in to ACTIVE command | CL = 3 | t_{DAL} | 5 | 4 | 4 | t_{CK} | 15, 22 |
| | CL = 2 | t_{DAL} | 4 | 3 | 4 | t_{CK} | 15, 22 |
| Data-in to PRECHARGE | | t_{DPL} | 1 | 1 | 1 | t_{CK} | 16 |
| Last data-in to burst STOP command | | t_{BDL} | 0 | 0 | 0 | t_{CK} | 17 |
| Last data-in to new READ/WRITE command | | t_{CDL} | 1 | 1 | 1 | t_{CK} | 17 |
| Last data-in to PRECHARGE command | | t_{RDL} | 1 | 1 | 1 | t_{CK} | 16 |
| LOAD MODE REGISTER command to command | | t_{MRD} | 2 | 2 | 2 | t_{CK} | 17 |
| Data-out to high-impedance from PRECHARGE command | CL = 3 | t_{ROH} | 3 | 3 | 3 | t_{CK} | 17 |
| | CL = 2 | t_{ROH} | 2 | 2 | 2 | t_{CK} | 17 |
| | CL = 1 | t_{ROH} | 1 | 1 | 1 | t_{CK} | 17 |

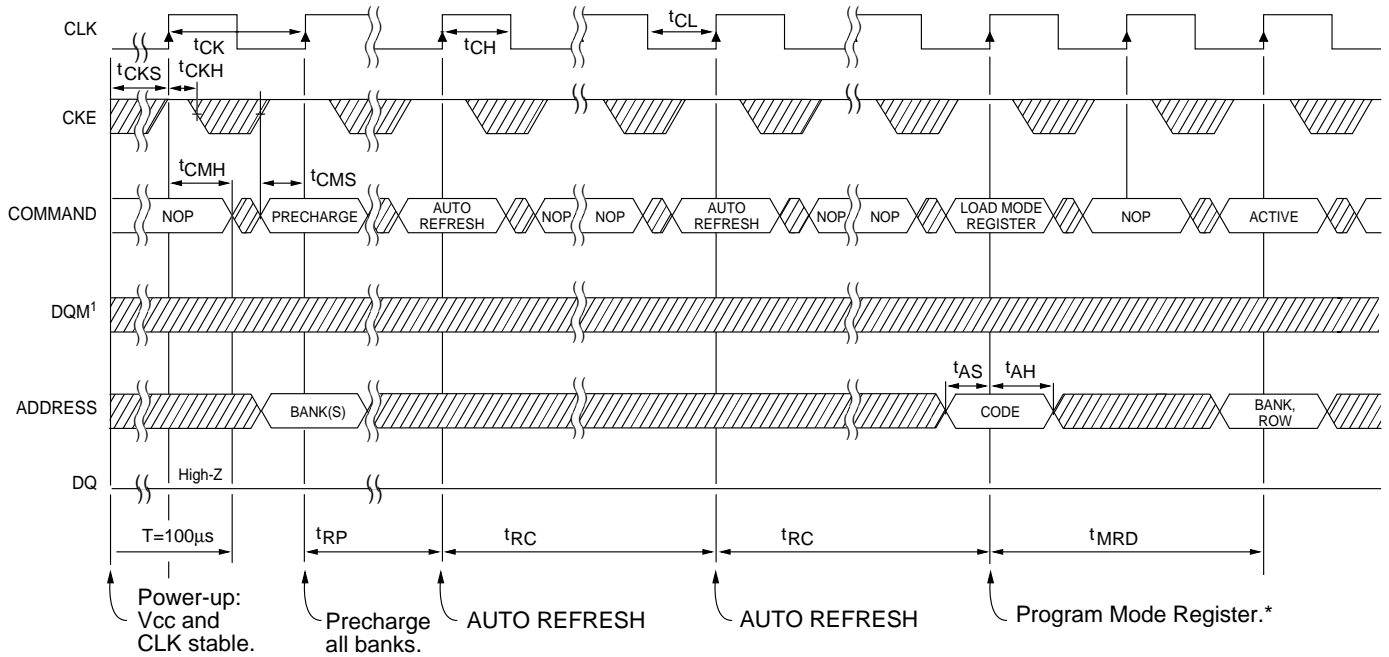
NOTES

1. All voltages referenced to V_{SS} .
2. This parameter is sampled. V_{CC} , $V_{CCQ} = +3.3V \pm 0.3V$; $f = 1\text{ MHz}$, $t_A = 25^\circ\text{C}$.
3. I_{CC} is dependent on output loading and cycle rates. Specified values are obtained with minimum cycle time and the outputs open.
4. Enables on-chip refresh and address counters.
5. The minimum specifications are used only to indicate cycle time at which proper operation over the full temperature range ($0^\circ\text{C} \leq T_A \leq 70^\circ\text{C}$) is ensured.
6. An initial pause of $100\mu\text{s}$ is required after power-up, followed by two AUTO REFRESH commands, before proper device operation is ensured. (V_{CC} and V_{CCQ} must be powered up simultaneously. V_{SS} and V_{SSQ} must be at same potential.) The two AUTO REFRESH command wake-ups should be repeated any time the t_{REF} refresh requirement is exceeded.
7. AC characteristics assume $t_T = 1\text{ ns}$.
8. In addition to meeting the transition rate specification, the clock and CKE must transit between V_{IH} and V_{IL} (or between V_{IL} and V_{IH}) in a monotonic manner.
9. Outputs measured at $1.4V$ with equivalent load:



$Z_O = 50\Omega$
 50Ω
 50pF
 $V_T = 1.4V$
10. t_{HZ} defines the time at which the output achieves the open circuit condition; it is not a reference to V_{OH} or V_{OL} . The last valid data element will meet t_{OH} before going High-Z.
11. AC timing and I_{CC} tests have $V_{IL} = 0V$ and $V_{IH} = 3V$ with timing referenced to $1.4V$ crossover point.
12. Other input signals are allowed to transition no more than once in any 30ns period (20ns on -8) and are otherwise at valid V_{IH} or V_{IL} levels.
13. I_{CC} specifications are tested after the device is properly initialized.
14. Timing actually specified by t_{CKS} ; clock(s) specified as a reference only at minimum cycle rate.
15. Timing actually specified by t_{WR} plus t_{RP} ; clock(s) specified as a reference only at minimum cycle rate.
16. Timing actually specified by t_{WR} .
17. Required clocks are specified by JEDEC functionality and not dependent on any timing parameter.
18. The I_{CC} current will decrease as the CAS latency is reduced. This is due to the fact that the maximum cycle rate is slower as the CAS latency is reduced.
19. Address transitions average one transition every 30ns (20ns on -8).
20. CLK must be toggled a minimum of two times during this period.
21. The device is designed to trade off t_{RP} for faster t_{WR} , i.e., 30ns and one clock, respectively. However, system designs should support a t_{WR} of two clocks to anticipate future SDRAM design requirements.
22. Based on $t_{CK} = 100\text{ MHz}$ for -8 and 66 MHz for -10.

INITIALIZE AND LOAD MODE REGISTER



*The Mode Register may be loaded prior to the AUTO REFRESH cycles if desired.

▨ DON'T CARE
▩ UNDEFINED

TIMING PARAMETERS

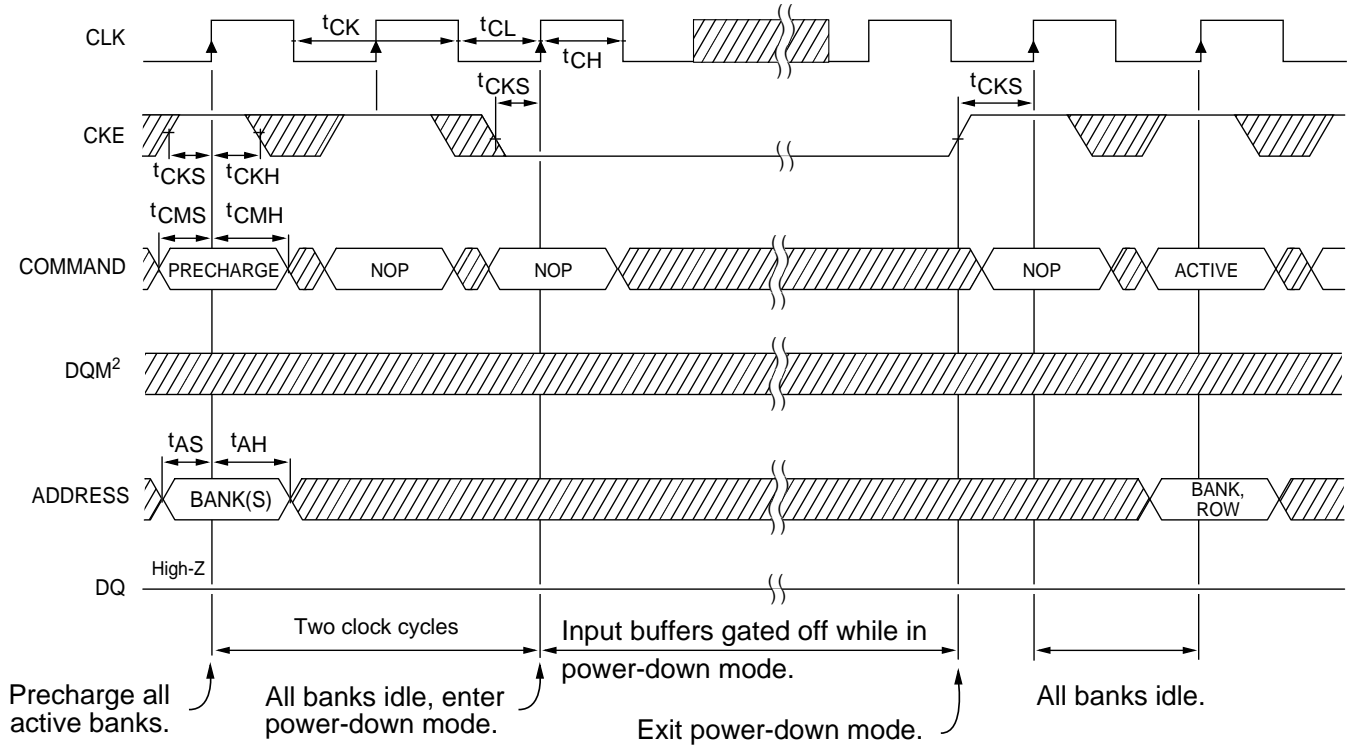
| SYMBOL** | -8A | | -10 | | -12 | | UNITS |
|---------------------|-----|-----|-----|-----|-----|-----|-------|
| | MIN | MAX | MIN | MAX | MIN | MAX | |
| t _{AH} | 1 | | 1 | | 1 | | ns |
| t _{AS} | 2 | | 3 | | 3 | | ns |
| t _{CH} | 3 | | 3.5 | | 3.5 | | ns |
| t _{CL} | 3 | | 3.5 | | 3.5 | | ns |
| t _{CK} (3) | 8 | | 10 | | 12 | | ns |
| t _{CK} (2) | 13 | | 15 | | 15 | | ns |
| t _{CK} (1) | 30 | | 30 | | 30 | | ns |

| SYMBOL** | -8A | | -10 | | -12 | | UNITS |
|------------------|-----|-----|-----|-----|-----|-----|-----------------|
| | MIN | MAX | MIN | MAX | MIN | MAX | |
| t _{CKH} | 1 | | 1 | | 1 | | ns |
| t _{CKS} | 2 | | 3 | | 3 | | ns |
| t _{CMH} | 1 | | 1 | | 1 | | ns |
| t _{CMS} | 2 | | 3 | | 3 | | ns |
| t _{MRD} | 2 | | 2 | | 2 | | t _{CK} |
| t _{RC} | 80 | | 90 | | 105 | | ns |
| t _{RP} | 30 | | 30 | | 36 | | ns |

**CAS latency indicated in parentheses.

NOTE: 1. DQM represents DQML and DQMH. DQML controls the lower byte, and DQMH controls the upper byte.

POWER-DOWN MODE ¹



▨ DON'T CARE
▩ UNDEFINED

TIMING PARAMETERS

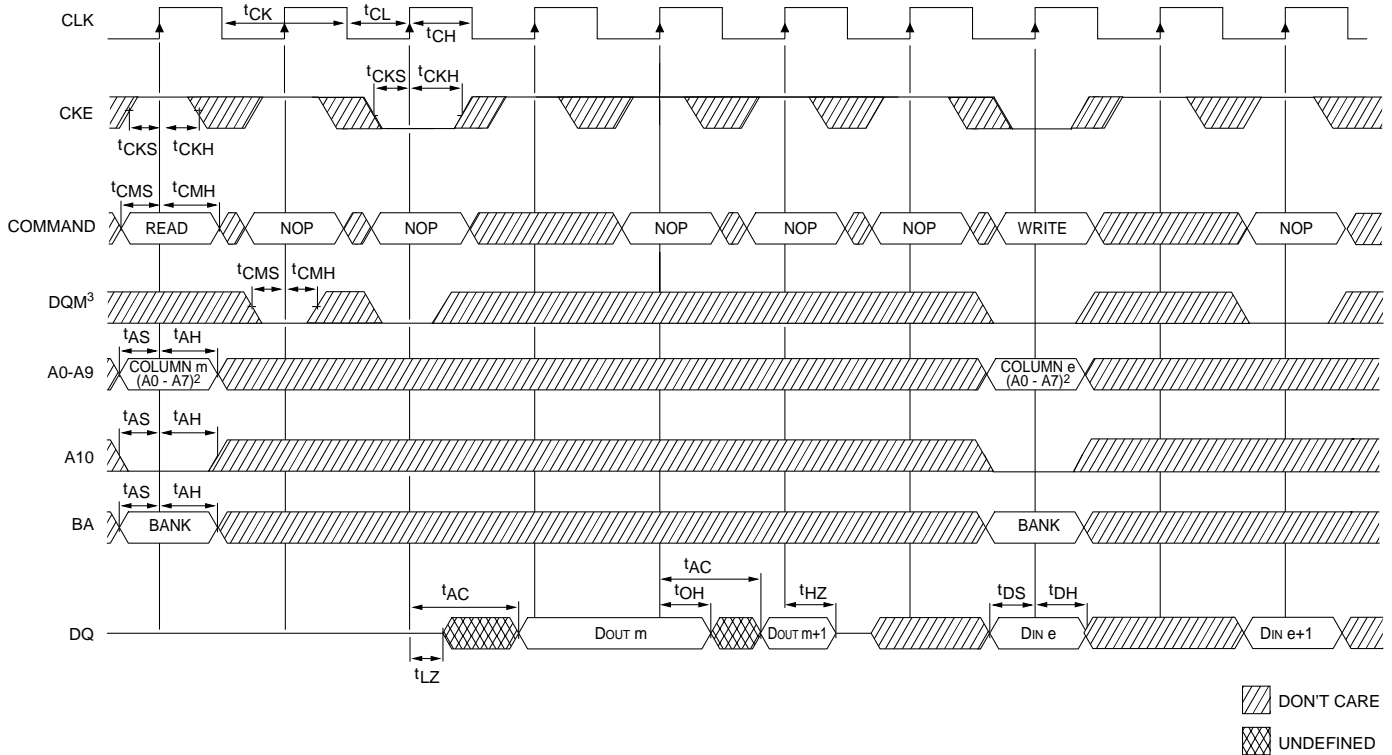
| SYMBOL* | -8A | | -10 | | -12 | | UNITS |
|--------------|-----|-----|-----|-----|-----|-----|-------|
| | MIN | MAX | MIN | MAX | MIN | MAX | |
| t_{AH} | 1 | | 1 | | 1 | | ns |
| t_{AS} | 2 | | 3 | | 3 | | ns |
| t_{CH} | 3 | | 3.5 | | 3.5 | | ns |
| t_{CL} | 3 | | 3.5 | | 3.5 | | ns |
| $t_{CK} (3)$ | 8 | | 10 | | 12 | | ns |
| $t_{CK} (2)$ | 13 | | 15 | | 15 | | ns |

| SYMBOL* | -8A | | -10 | | -12 | | UNITS |
|--------------|-----|-----|-----|-----|-----|-----|-------|
| | MIN | MAX | MIN | MAX | MIN | MAX | |
| $t_{CK} (1)$ | 30 | | 30 | | 30 | | ns |
| t_{CKH} | 1 | | 1 | | 1 | | ns |
| t_{CKS} | 2 | | 3 | | 3 | | ns |
| t_{CMH} | 1 | | 1 | | 1 | | ns |
| t_{CMS} | 2 | | 3 | | 3 | | ns |

*CAS latency indicated in parentheses.

NOTE: 1. Violating refresh requirements during power-down may result in loss of data.
2. DQM represents DQML and DQMH. DQML controls the lower byte, and DQMH controls the upper byte.

CLOCK SUSPEND MODE 1



TIMING PARAMETERS

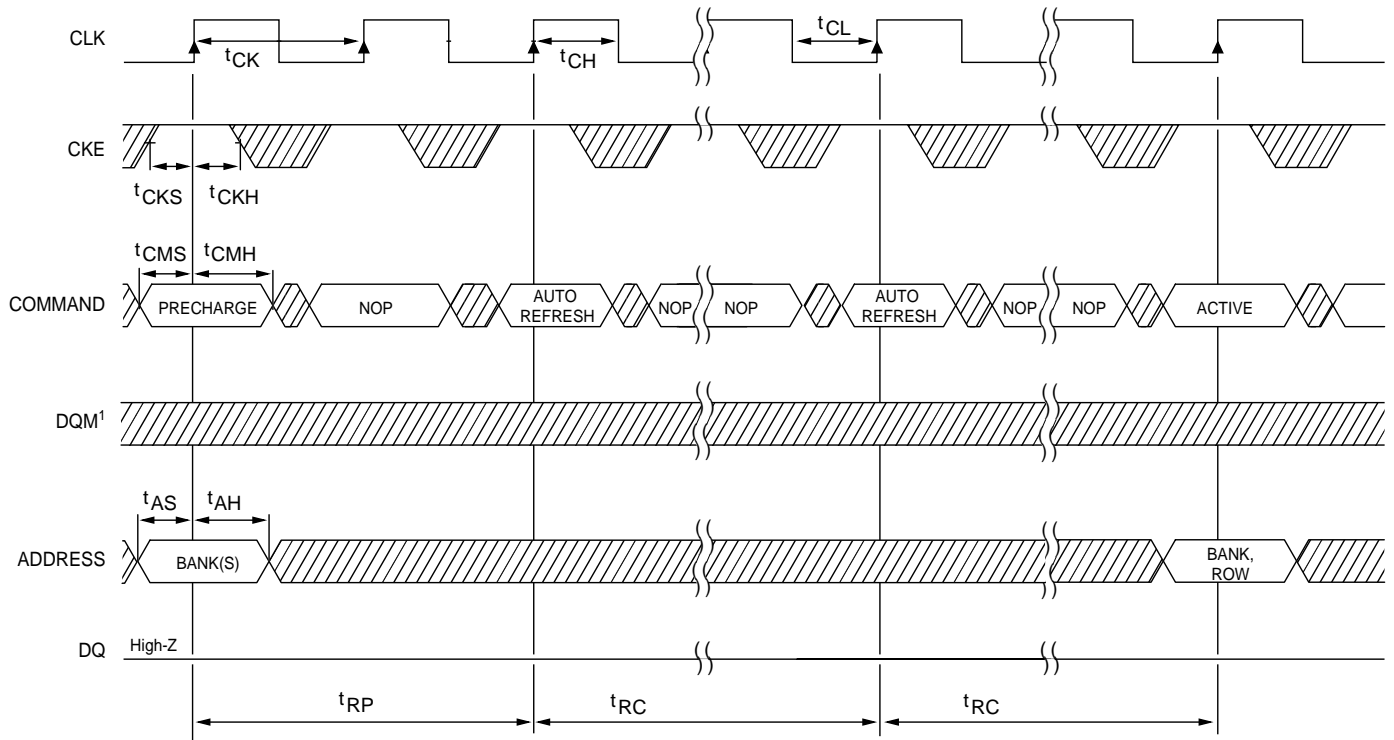
| SYMBOL* | -8A | | -10 | | -12 | | UNITS |
|-------------|-----|-----|-----|-----|-----|-----|-------|
| | MIN | MAX | MIN | MAX | MIN | MAX | |
| $t_{AC}(3)$ | | 6 | | 7.5 | | 9 | ns |
| $t_{AC}(2)$ | | 9 | | 9 | | 9 | ns |
| $t_{AC}(1)$ | | 27 | | 27 | | 27 | ns |
| t_{AH} | 1 | | 1 | | 1 | | ns |
| t_{AS} | 2 | | 3 | | 3 | | ns |
| t_{CH} | 3 | | 3.5 | | 3.5 | | ns |
| t_{CL} | 3 | | 3.5 | | 3.5 | | ns |
| $t_{CK}(3)$ | 8 | | 10 | | 12 | | ns |
| $t_{CK}(2)$ | 13 | | 15 | | 15 | | ns |
| $t_{CK}(1)$ | 30 | | 30 | | 30 | | ns |
| t_{CKH} | 1 | | 1 | | 1 | | ns |

| SYMBOL* | -8A | | -10 | | -12 | | UNITS |
|-------------|-----|-----|-----|-----|-----|-----|-------|
| | MIN | MAX | MIN | MAX | MIN | MAX | |
| t_{CKS} | 2 | | 3 | | 3 | | ns |
| t_{CMH} | 1 | | 1 | | 1 | | ns |
| t_{CMS} | 2 | | 3 | | 3 | | ns |
| t_{DH} | 1 | | 1 | | 1 | | ns |
| t_{DS} | 2 | | 3 | | 3 | | ns |
| $t_{HZ}(3)$ | | 6 | | 8 | | 9 | ns |
| $t_{HZ}(2)$ | | 7 | | 10 | | 10 | ns |
| $t_{HZ}(1)$ | | 15 | | 15 | | 15 | ns |
| t_{LZ} | 1 | | 2 | | 2 | | ns |
| t_{OH} | 3 | | 3 | | 3 | | ns |

*CAS latency indicated in parentheses.

- NOTE:**
1. For this example, the burst length = 2, the CAS latency = 3 and AUTO PRECHARGE is disabled.
 2. A8 and A9 = "Don't Care."
 3. DQM represents DQML and DQMH. DQML controls the lower byte, and DQMH controls the upper byte.

AUTO REFRESH MODE



Precharge all active banks.

DON'T CARE
 UNDEFINED

TIMING PARAMETERS

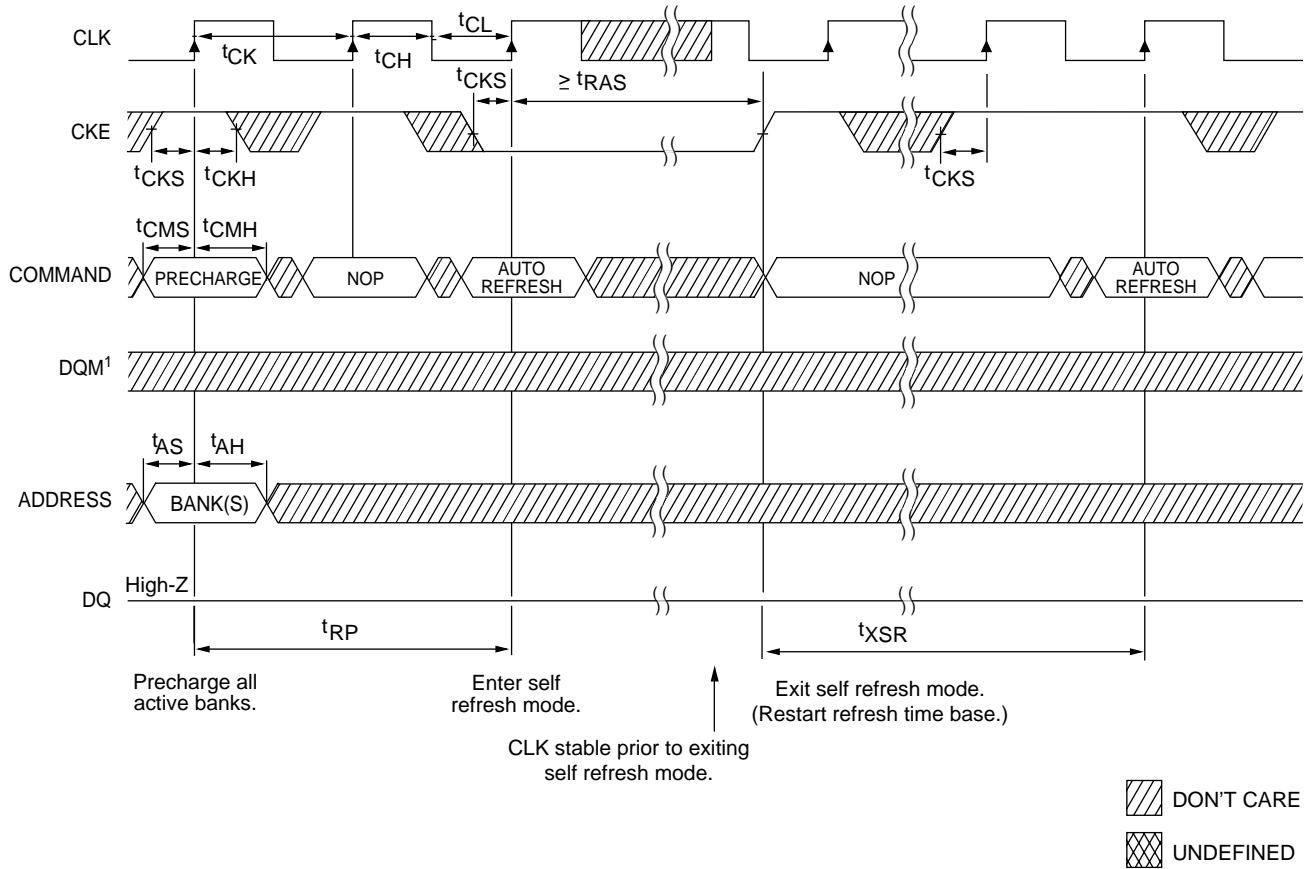
| SYMBOL* | -8A | | -10 | | -12 | | UNITS |
|---------------------|-----|-----|-----|-----|-----|-----|-------|
| | MIN | MAX | MIN | MAX | MIN | MAX | |
| t _{AH} | 1 | | 1 | | 1 | | ns |
| t _{AS} | 2 | | 3 | | 3 | | ns |
| t _{CH} | 3 | | 3.5 | | 3.5 | | ns |
| t _{CL} | 3 | | 3.5 | | 3.5 | | ns |
| t _{CK (3)} | 8 | | 10 | | 12 | | ns |
| t _{CK (2)} | 13 | | 15 | | 15 | | ns |
| t _{CK (1)} | 30 | | 30 | | 30 | | ns |

| SYMBOL* | -8A | | -10 | | -12 | | UNITS |
|------------------|-----|-----|-----|-----|-----|-----|-------|
| | MIN | MAX | MIN | MAX | MIN | MAX | |
| t _{CKH} | 1 | | 1 | | 1 | | ns |
| t _{CKS} | 2 | | 3 | | 3 | | ns |
| t _{CMH} | 1 | | 1 | | 1 | | ns |
| t _{CMS} | 2 | | 3 | | 3 | | ns |
| t _{RC} | 80 | | 90 | | 105 | | ns |
| t _{RP} | 30 | | 30 | | 36 | | ns |

*CAS latency indicated in parentheses.

NOTE: 1. DQM represents DQML and DQMH. DQML controls the lower byte, and DQMH controls the upper byte.

SELF REFRESH MODE



TIMING PARAMETERS

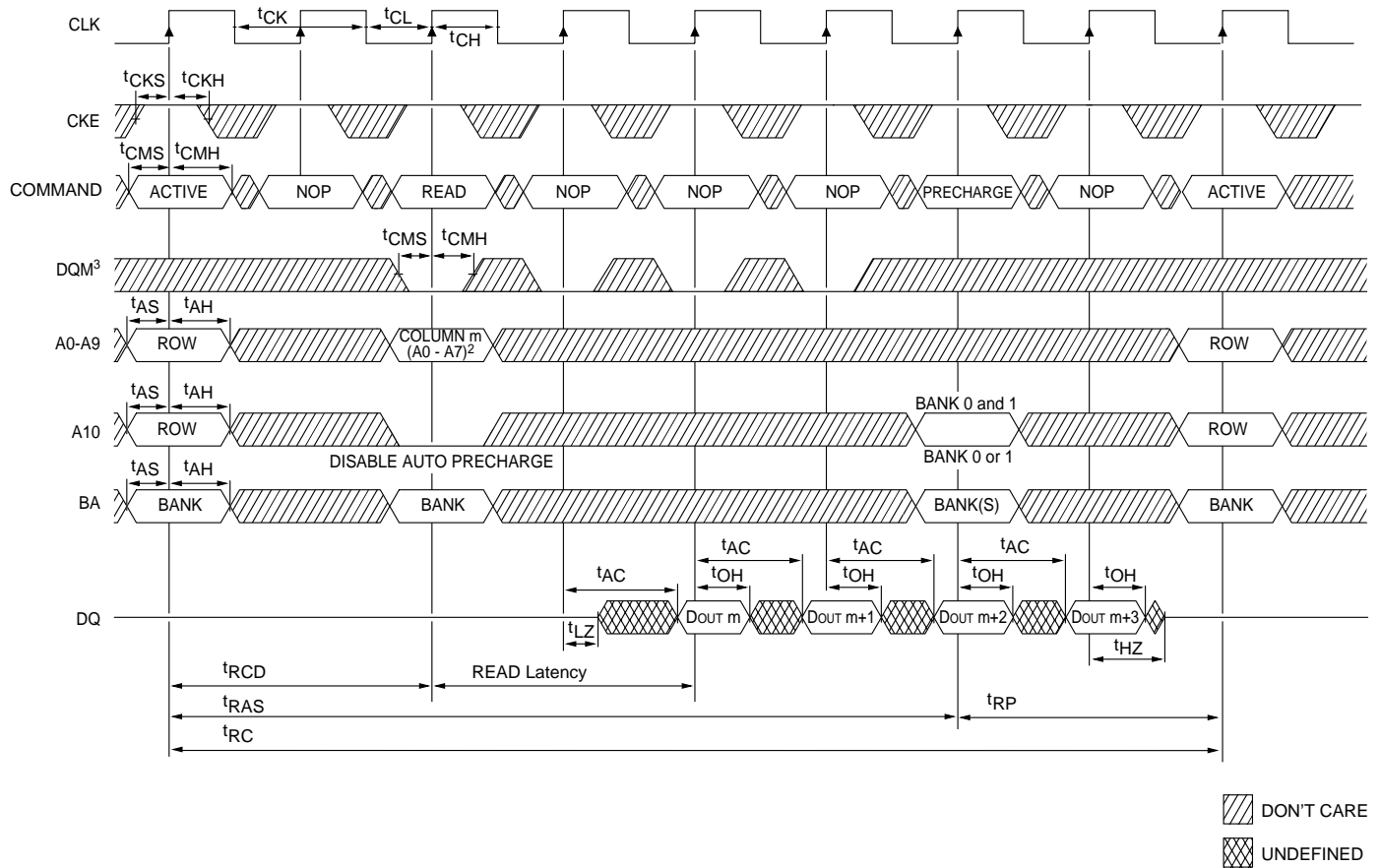
| SYMBOL* | -8A | | -10 | | -12 | | UNITS |
|--------------|-----|-----|-----|-----|-----|-----|-------|
| | MIN | MAX | MIN | MAX | MIN | MAX | |
| t_{AH} | 1 | | 1 | | 1 | | ns |
| t_{AS} | 2 | | 3 | | 3 | | ns |
| t_{CH} | 3 | | 3.5 | | 3.5 | | ns |
| t_{CL} | 3 | | 3.5 | | 3.5 | | ns |
| $t_{CK} (3)$ | 8 | | 10 | | 12 | | ns |
| $t_{CK} (2)$ | 13 | | 15 | | 15 | | ns |
| $t_{CK} (1)$ | 30 | | 30 | | 30 | | ns |

| SYMBOL* | -8A | | -10 | | -12 | | UNITS |
|-----------|-----|------|-----|------|-----|------|-------|
| | MIN | MAX | MIN | MAX | MIN | MAX | |
| t_{CKH} | 1 | | 1 | | 1 | | ns |
| t_{CKS} | 2 | | 3 | | 3 | | ns |
| t_{CMH} | 1 | | 1 | | 1 | | ns |
| t_{CMS} | 2 | | 3 | | 3 | | ns |
| t_{RAS} | 50 | 120K | 60 | 120K | 72 | 120K | ns |
| t_{RP} | 30 | | 30 | | 36 | | ns |
| t_{XSR} | 80 | | 96 | | 105 | | ns |

*CAS latency indicated in parentheses.

NOTE: 1. DQM represents DQML and DQMH. DQML controls the lower byte, and DQMH controls the upper byte.

READ – WITHOUT AUTO PRECHARGE ¹



TIMING PARAMETERS

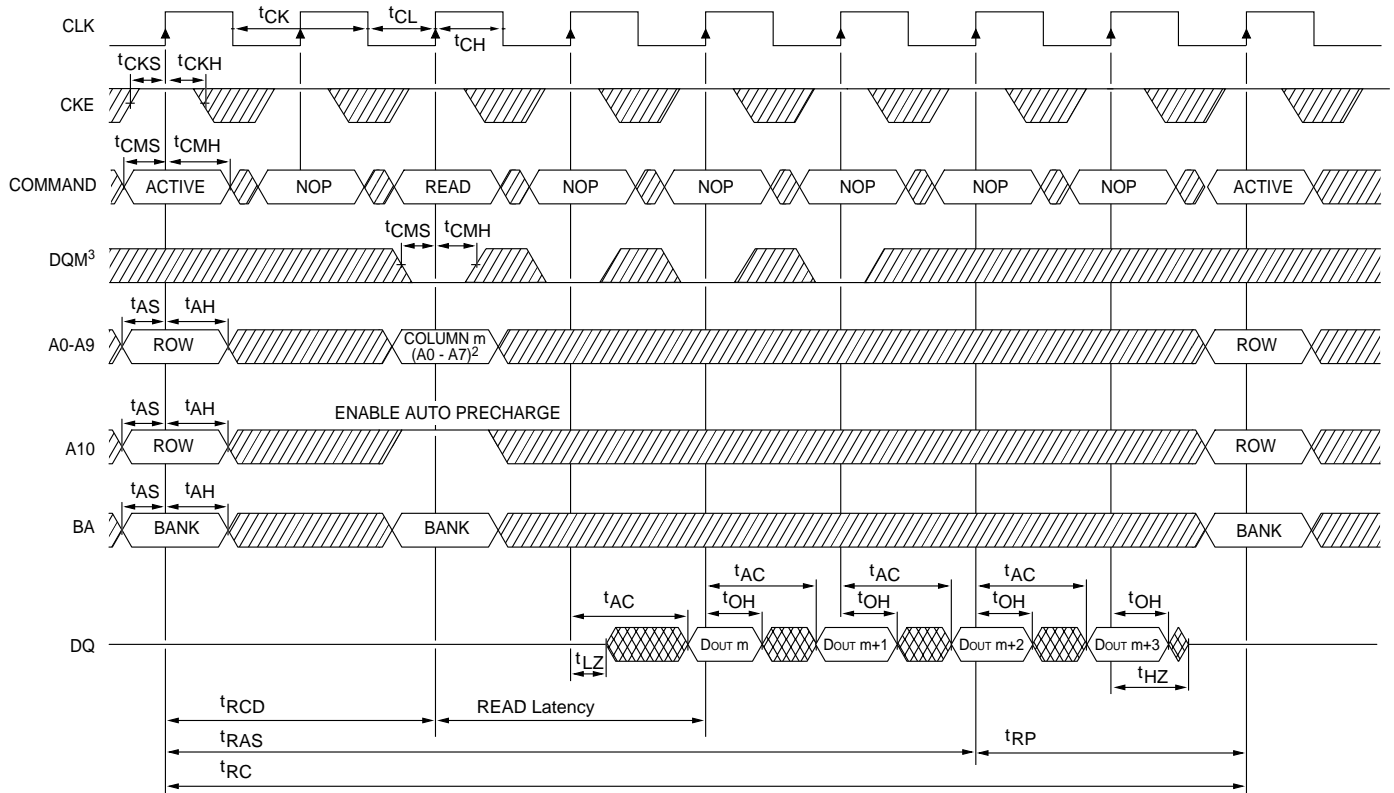
| SYMBOL* | -8A | | -10 | | -12 | | UNITS |
|---------------------|-----|-----|-----|-----|-----|-----|-------|
| | MIN | MAX | MIN | MAX | MIN | MAX | |
| t ^{AC} (3) | | 6 | | 7.5 | | 9 | ns |
| t ^{AC} (2) | | 9 | | 9 | | 9 | ns |
| t ^{AC} (1) | | 27 | | 27 | | 27 | ns |
| t ^{AH} | 1 | | 1 | | 1 | | ns |
| t ^{AS} | 2 | | 3 | | 3 | | ns |
| t ^{CH} | 3 | | 3.5 | | 3.5 | | ns |
| t ^{CL} | 3 | | 3.5 | | 3.5 | | ns |
| t ^{CK} (3) | 8 | | 10 | | 12 | | ns |
| t ^{CK} (2) | 13 | | 15 | | 15 | | ns |
| t ^{CK} (1) | 30 | | 30 | | 30 | | ns |
| t ^{CKH} | 1 | | 1 | | 1 | | ns |
| t ^{CKS} | 2 | | 3 | | 3 | | ns |

| SYMBOL* | -8A | | -10 | | -12 | | UNITS |
|---------------------|-----|------|-----|------|-----|------|-------|
| | MIN | MAX | MIN | MAX | MIN | MAX | |
| t ^{CMH} | 1 | | 1 | | 1 | | ns |
| t ^{CMS} | 2 | | 3 | | 3 | | ns |
| t ^{HZ} (3) | | 6 | | 8 | | 9 | ns |
| t ^{HZ} (2) | | 7 | | 10 | | 10 | ns |
| t ^{HZ} (1) | | 15 | | 15 | | 15 | ns |
| t ^{LZ} | 1 | | 2 | | 2 | | ns |
| t ^{OH} | 3 | | 3 | | 3 | | ns |
| t ^{RAS} | 50 | 120K | 60 | 120K | 72 | 120K | ns |
| t ^{RC} | 80 | | 90 | | 105 | | ns |
| t ^{RCD} | 30 | | 30 | | 30 | | ns |
| t ^{RP} | 30 | | 30 | | 36 | | ns |

*CAS latency indicated in parentheses.

- NOTE:**
1. For this example, the burst length = 4, the READ latency = 2 and the READ burst is followed by a "manual" PRECHARGE.
 2. A8 and A9 = "Don't Care."
 3. DQM represents DQML and DQMH. DQML controls the lower byte, and DQMH controls the upper byte.

READ – WITH AUTO PRECHARGE 1



▨ DON'T CARE
▩ UNDEFINED

TIMING PARAMETERS

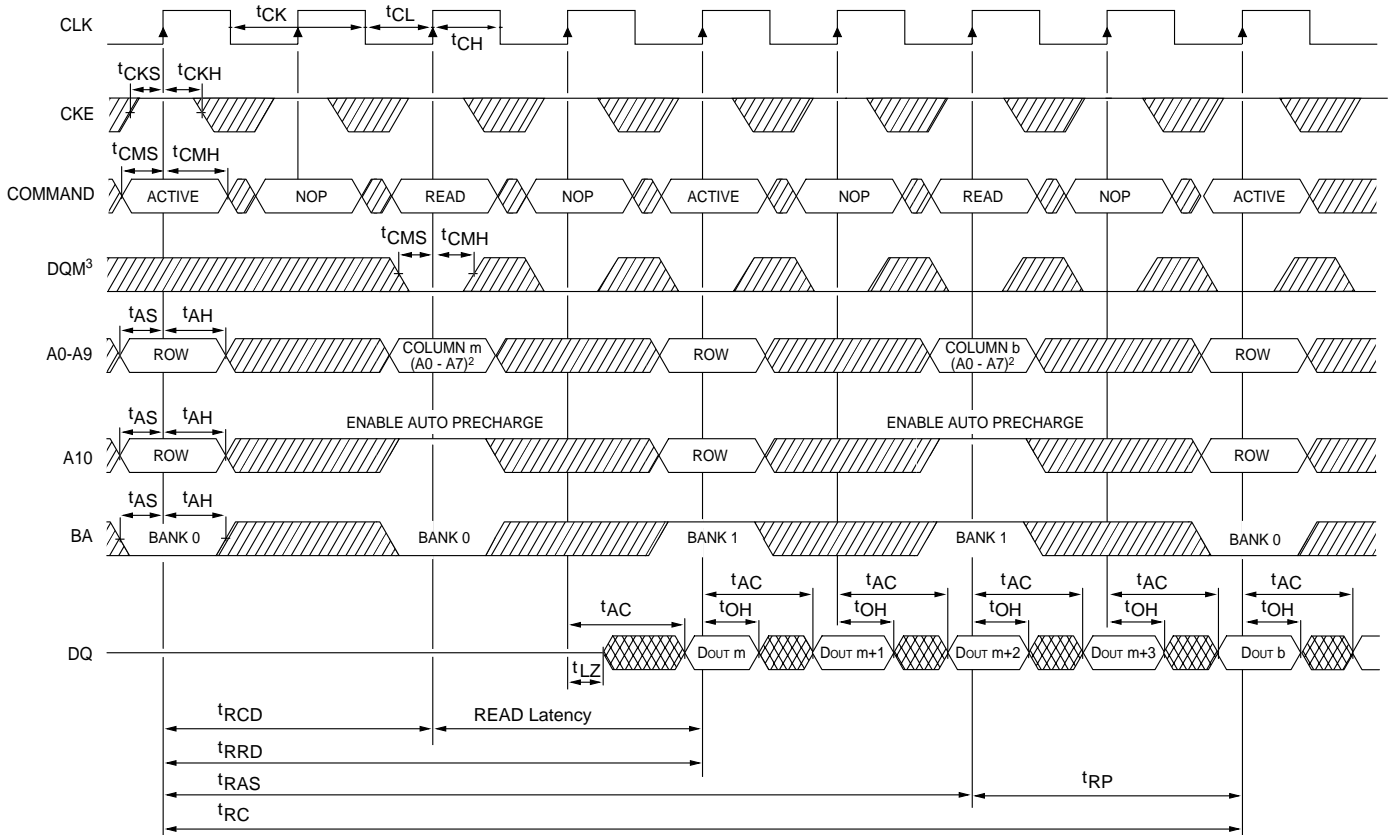
| SYMBOL* | -8A | | -10 | | -12 | | UNITS |
|-----------------------|-----|-----|-----|-----|-----|-----|-------|
| | MIN | MAX | MIN | MAX | MIN | MAX | |
| t ¹ AC(3) | | 6 | | 7.5 | | 9 | ns |
| t ¹ AC(2) | | 9 | | 9 | | 9 | ns |
| t ¹ AC(1) | | 27 | | 27 | | 27 | ns |
| t ¹ AH | 1 | | 1 | | 1 | | ns |
| t ¹ AS | 2 | | 3 | | 3 | | ns |
| t ¹ CH | 3 | | 3.5 | | 3.5 | | ns |
| t ¹ CL | 3 | | 3.5 | | 3.5 | | ns |
| t ¹ CK (3) | 8 | | 10 | | 12 | | ns |
| t ¹ CK (2) | 13 | | 15 | | 15 | | ns |
| t ¹ CK (1) | 30 | | 30 | | 30 | | ns |
| t ¹ CKH | 1 | | 1 | | 1 | | ns |
| t ¹ CKS | 2 | | 3 | | 3 | | ns |

| SYMBOL* | -8A | | -10 | | -12 | | UNITS |
|-----------------------|-----|------|-----|------|-----|------|-------|
| | MIN | MAX | MIN | MAX | MIN | MAX | |
| t ¹ CMH | 1 | | 1 | | 1 | | ns |
| t ¹ CMS | 2 | | 3 | | 3 | | ns |
| t ¹ HZ (3) | | 6 | | 8 | | 9 | ns |
| t ¹ HZ (2) | | 7 | | 10 | | 10 | ns |
| t ¹ HZ (1) | | 15 | | 15 | | 15 | ns |
| t ¹ LZ | 1 | | 2 | | 2 | | ns |
| t ¹ OH | 3 | | 3 | | 3 | | ns |
| t ¹ RAS | 50 | 120K | 60 | 120K | 72 | 120K | ns |
| t ¹ RC | 80 | | 90 | | 105 | | ns |
| t ¹ RCD | 30 | | 30 | | 30 | | ns |
| t ¹ RP | 30 | | 30 | | 36 | | ns |

*CAS latency indicated in parentheses.

- NOTE:**
1. For this example, the burst length = 4, and the READ latency = 2.
 2. A8 and A9 = "Don't Care."
 3. DQM represents DQML and DQMH. DQML controls the lower byte, and DQMH controls the upper byte.

ALTERNATING BANK READ ACCESSES 1



▨ DON'T CARE
▩ UNDEFINED

TIMING PARAMETERS

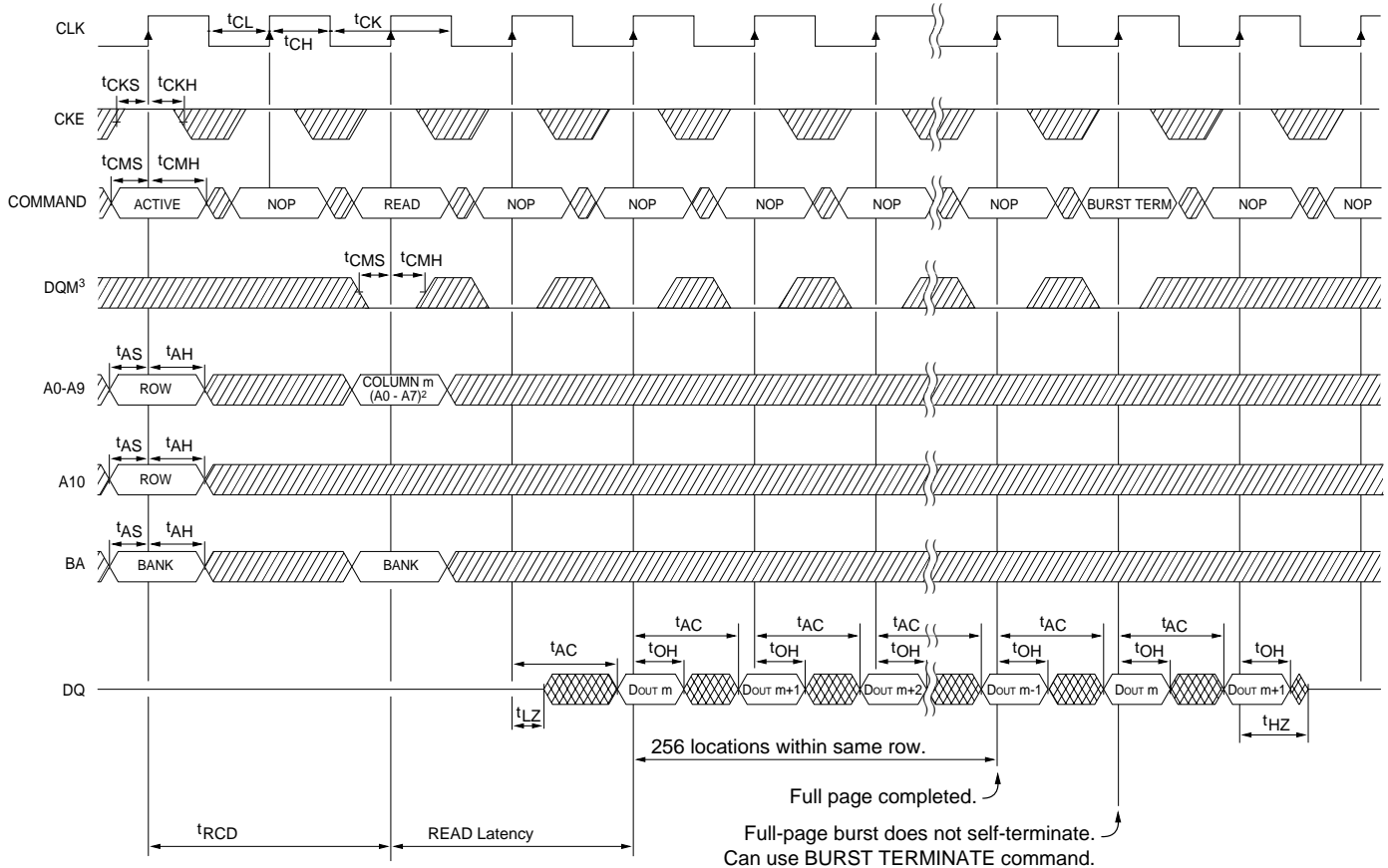
| SYMBOL* | -8A | | -10 | | -12 | | UNITS |
|-----------------------|-----|-----|-----|-----|-----|-----|-------|
| | MIN | MAX | MIN | MAX | MIN | MAX | |
| t ¹ AC(3) | | 6 | | 7.5 | | 9 | ns |
| t ¹ AC(2) | | 9 | | 9 | | 9 | ns |
| t ¹ AC(1) | | 27 | | 27 | | 27 | ns |
| t ¹ AH | 1 | | 1 | | 1 | | ns |
| t ¹ AS | 2 | | 3 | | 3 | | ns |
| t ¹ CH | 3 | | 3.5 | | 3.5 | | ns |
| t ¹ CL | 3 | | 3.5 | | 3.5 | | ns |
| t ¹ CK (3) | 8 | | 10 | | 12 | | ns |
| t ¹ CK (2) | 13 | | 15 | | 15 | | ns |
| t ¹ CK (1) | 30 | | 30 | | 30 | | ns |
| t ¹ CKH | 1 | | 1 | | 1 | | ns |

| SYMBOL* | -8A | | -10 | | -12 | | UNITS |
|--------------------|-----|------|-----|------|-----|------|-------|
| | MIN | MAX | MIN | MAX | MIN | MAX | |
| t ¹ CKS | 2 | | 3 | | 3 | | ns |
| t ¹ CMH | 1 | | 1 | | 1 | | ns |
| t ¹ CMS | 2 | | 3 | | 3 | | ns |
| t ¹ LZ | 1 | | 2 | | 2 | | ns |
| t ¹ OH | 3 | | 3 | | 3 | | ns |
| t ¹ RAS | 50 | 120K | 60 | 120K | 72 | 120K | ns |
| t ¹ RC | 80 | | 90 | | 105 | | ns |
| t ¹ RCD | 30 | | 30 | | 30 | | ns |
| t ¹ RP | 30 | | 30 | | 36 | | ns |
| t ¹ RRD | 20 | | 20 | | 20 | | ns |

*CAS latency indicated in parentheses.

- NOTE:**
1. For this example, the burst length = 4, and the READ latency = 2.
 2. A8 and A9 = "Don't Care."
 3. DQM represents DQML and DQMH. DQML controls the lower byte, and DQMH controls the upper byte.

READ – FULL-PAGE BURST ¹



DONT CARE
 UNDEFINED

TIMING PARAMETERS

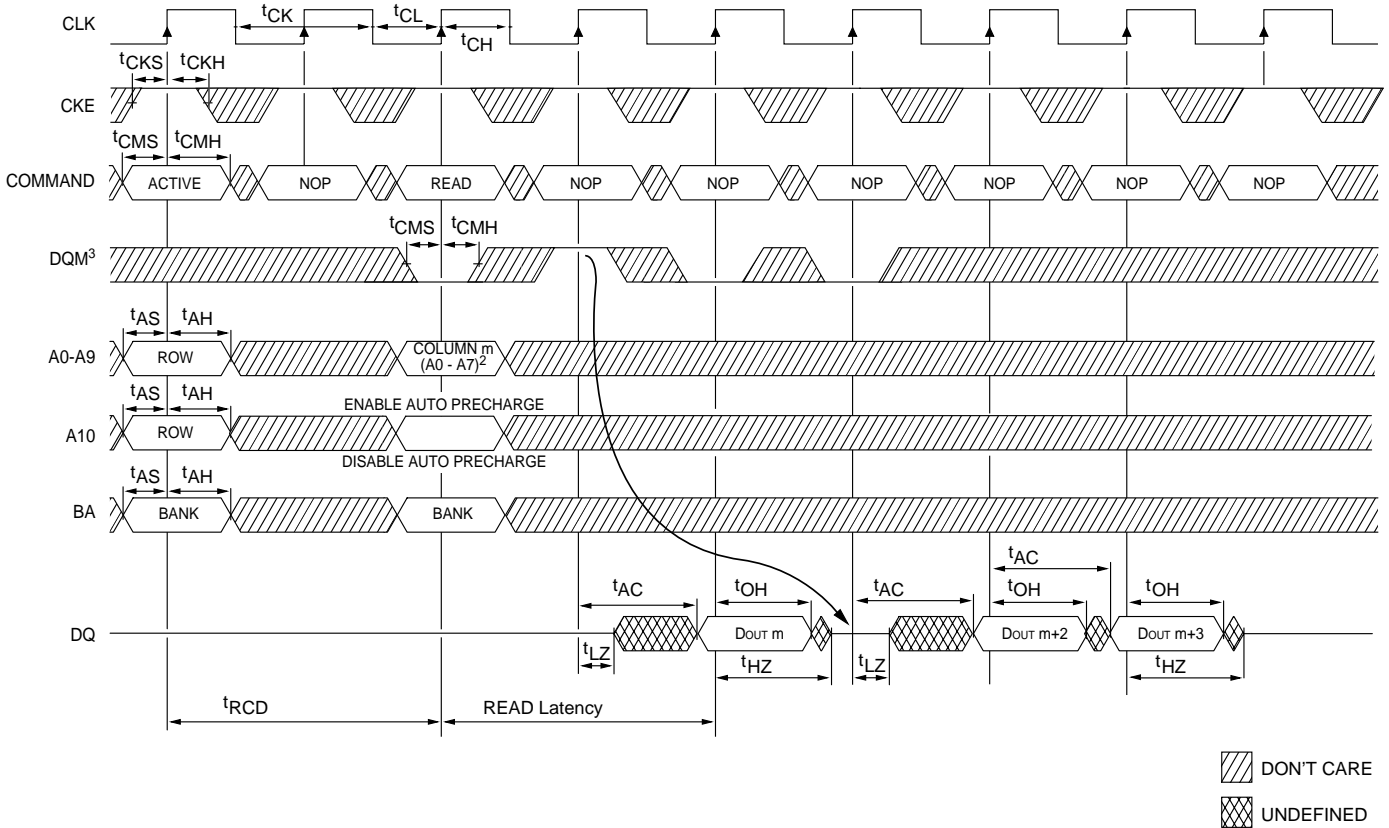
| SYMBOL* | -8A | | -10 | | -12 | | UNITS |
|---------------------|-----|-----|-----|-----|-----|-----|-------|
| | MIN | MAX | MIN | MAX | MIN | MAX | |
| ^t AC(3) | | 6 | | 7.5 | | 9 | ns |
| ^t AC(2) | | 9 | | 9 | | 9 | ns |
| ^t AC(1) | | 27 | | 27 | | 27 | ns |
| ^t AH | 1 | | 1 | | 1 | | ns |
| ^t AS | 2 | | 3 | | 3 | | ns |
| ^t CH | 3 | | 3.5 | | 3.5 | | ns |
| ^t CL | 3 | | 3.5 | | 3.5 | | ns |
| ^t CK (3) | 8 | | 10 | | 12 | | ns |
| ^t CK (2) | 13 | | 15 | | 15 | | ns |
| ^t CK (1) | 30 | | 30 | | 30 | | ns |

| SYMBOL* | -8A | | -10 | | -12 | | UNITS |
|---------------------|-----|-----|-----|-----|-----|-----|-------|
| | MIN | MAX | MIN | MAX | MIN | MAX | |
| ^t CKH | 1 | | 1 | | 1 | | ns |
| ^t CKS | 2 | | 3 | | 3 | | ns |
| ^t CMH | 1 | | 1 | | 1 | | ns |
| ^t CMS | 2 | | 3 | | 3 | | ns |
| ^t HZ (3) | | 6 | | 8 | | 9 | ns |
| ^t HZ (2) | | 7 | | 10 | | 10 | ns |
| ^t HZ (1) | | 15 | | 15 | | 15 | ns |
| ^t LZ | 1 | | 2 | | 2 | | ns |
| ^t OH | 3 | | 3 | | 3 | | ns |
| ^t RCD | 30 | | 30 | | 30 | | ns |

*CAS latency indicated in parentheses.

- NOTE:**
1. For this example, the READ latency = 2.
 2. A8 and A9 = "Don't Care."
 3. DQM represents DQML and DQMH. DQML controls the lower byte, and DQMH controls the upper byte.

READ – DQM OPERATION 1



TIMING PARAMETERS

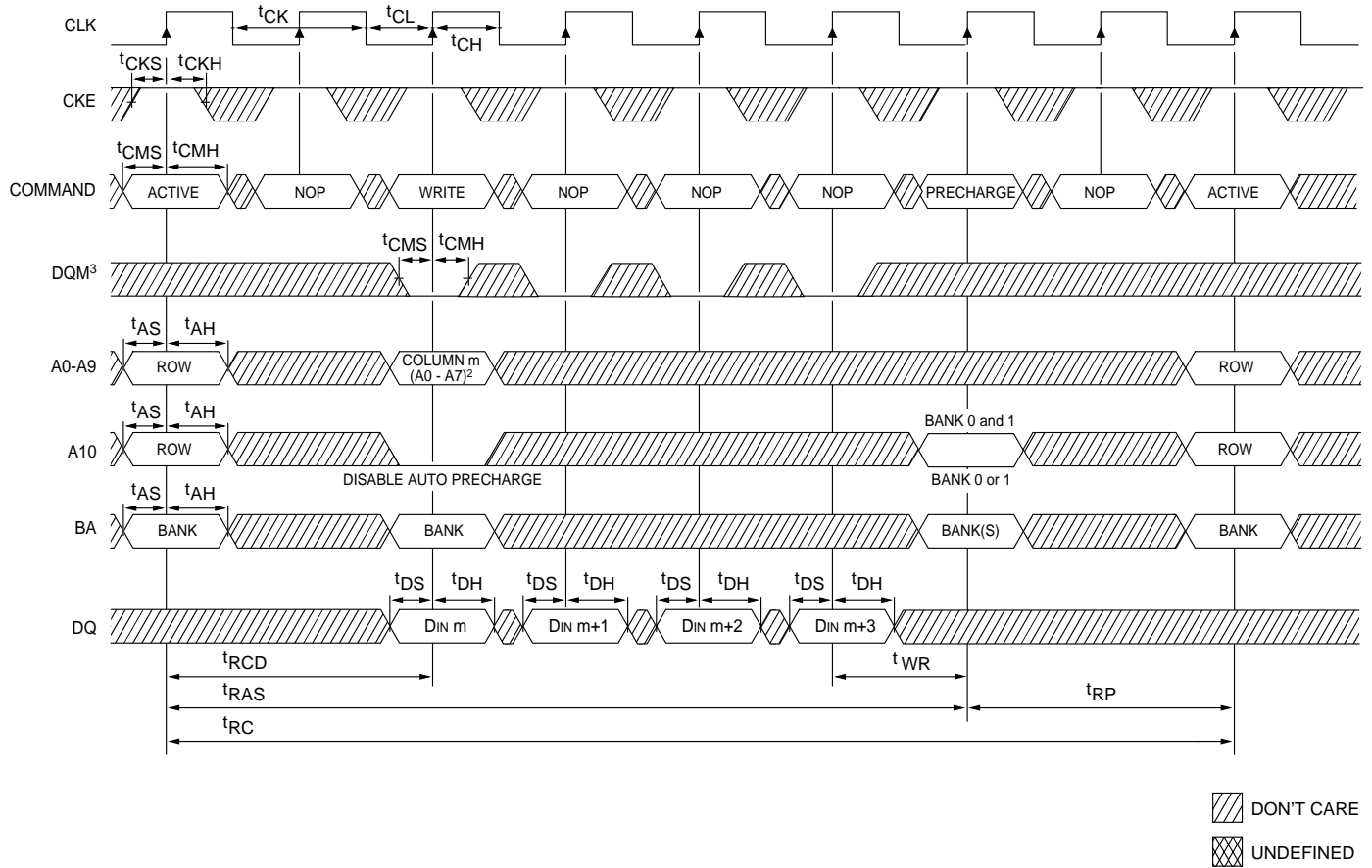
| SYMBOL* | -8A | | -10 | | -12 | | UNITS |
|-------------|-----|-----|-----|-----|-----|-----|-------|
| | MIN | MAX | MIN | MAX | MIN | MAX | |
| $t_{AC}(3)$ | | 6 | | 7.5 | | 9 | ns |
| $t_{AC}(2)$ | | 9 | | 9 | | 9 | ns |
| $t_{AC}(1)$ | | 27 | | 27 | | 27 | ns |
| t_{AH} | 1 | | 1 | | 1 | | ns |
| t_{AS} | 2 | | 3 | | 3 | | ns |
| t_{CH} | 3 | | 3.5 | | 3.5 | | ns |
| t_{CL} | 3 | | 3.5 | | 3.5 | | ns |
| $t_{CK}(3)$ | 8 | | 10 | | 12 | | ns |
| $t_{CK}(2)$ | 13 | | 15 | | 15 | | ns |
| $t_{CK}(1)$ | 30 | | 30 | | 30 | | ns |

| SYMBOL* | -8A | | -10 | | -12 | | UNITS |
|-------------|-----|-----|-----|-----|-----|-----|-------|
| | MIN | MAX | MIN | MAX | MIN | MAX | |
| t_{CKH} | 1 | | 1 | | 1 | | ns |
| t_{CKS} | 2 | | 3 | | 3 | | ns |
| t_{CMH} | 1 | | 1 | | 1 | | ns |
| t_{CMS} | 2 | | 3 | | 3 | | ns |
| $t_{HZ}(3)$ | | 6 | | 8 | | 9 | ns |
| $t_{HZ}(2)$ | | 7 | | 10 | | 10 | ns |
| $t_{HZ}(1)$ | | 15 | | 15 | | 15 | ns |
| t_{LZ} | 1 | | 2 | | 2 | | ns |
| t_{OH} | 3 | | 3 | | 3 | | ns |
| t_{RCD} | 30 | | 30 | | 30 | | ns |

*CAS latency indicated in parentheses.

- NOTE:**
1. For this example, the burst length = 4, and the READ latency = 2.
 2. A8 and A9 = "Don't Care."
 3. DQM represents DQML and DQMH. DQML controls the lower byte, and DQMH controls the upper byte.

WRITE – WITHOUT AUTO PRECHARGE 1



TIMING PARAMETERS

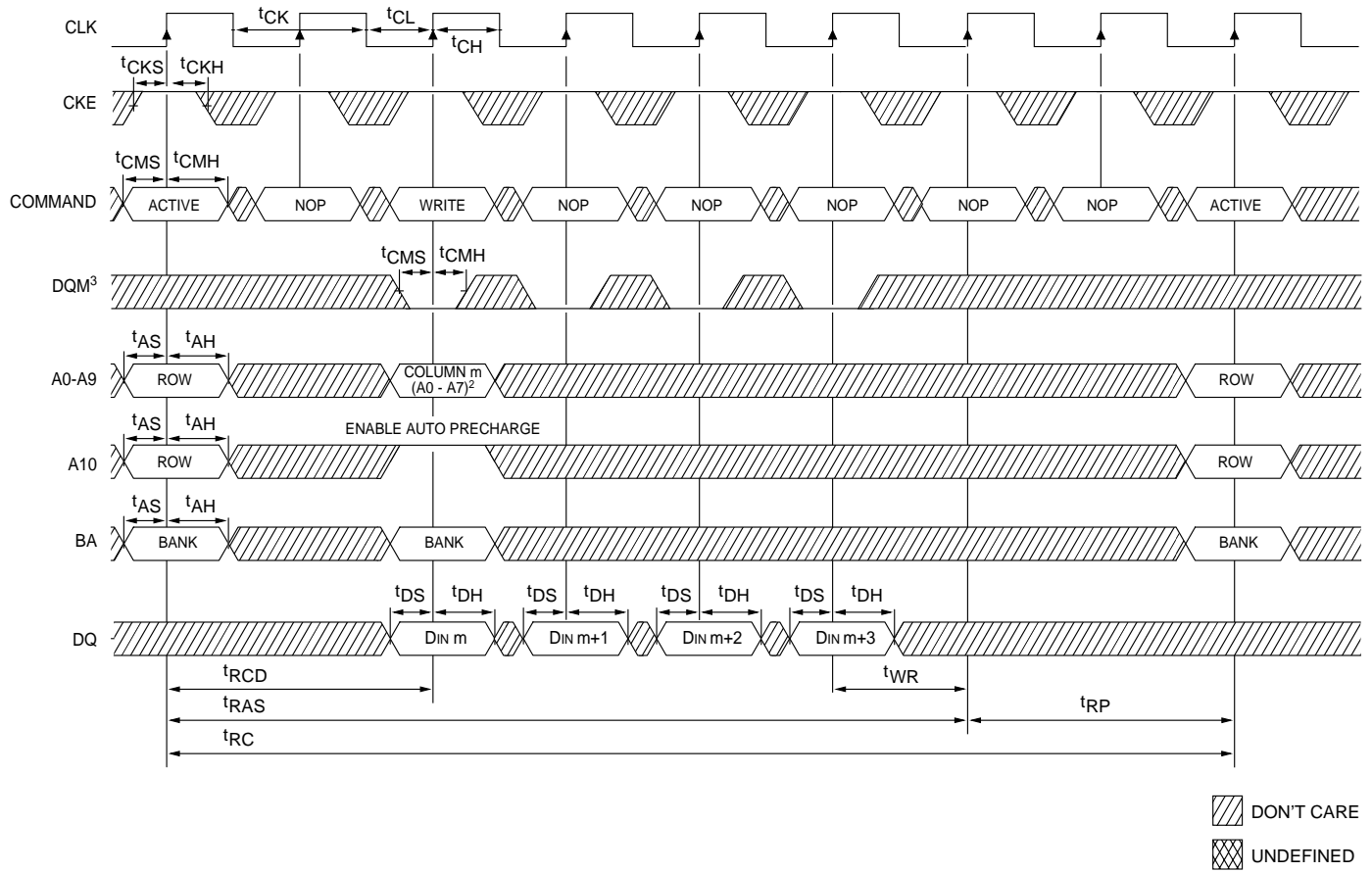
| SYMBOL* | -8A | | -10 | | -12 | | UNITS |
|---------------------|-----|-----|-----|-----|-----|-----|-------|
| | MIN | MAX | MIN | MAX | MIN | MAX | |
| t _{AH} | 1 | | 1 | | 1 | | ns |
| t _{AS} | 2 | | 3 | | 3 | | ns |
| t _{CH} | 3 | | 3.5 | | 3.5 | | ns |
| t _{CL} | 3 | | 3.5 | | 3.5 | | ns |
| t _{CK (3)} | 8 | | 10 | | 12 | | ns |
| t _{CK (2)} | 13 | | 15 | | 15 | | ns |
| t _{CK (1)} | 30 | | 30 | | 30 | | ns |
| t _{CKH} | 1 | | 1 | | 1 | | ns |
| t _{CKS} | 2 | | 3 | | 3 | | ns |

| SYMBOL* | -8A | | -10 | | -12 | | UNITS |
|------------------|-----|------|-----|------|-----|------|-----------------|
| | MIN | MAX | MIN | MAX | MIN | MAX | |
| t _{CMH} | 1 | | 1 | | 1 | | ns |
| t _{CMS} | 2 | | 3 | | 3 | | ns |
| t _{DH} | 1 | | 1 | | 1 | | ns |
| t _{DS} | 2 | | 3 | | 3 | | ns |
| t _{RAS} | 50 | 120K | 60 | 120K | 72 | 120K | ns |
| t _{RC} | 80 | | 90 | | 105 | | ns |
| t _{RCD} | 30 | | 30 | | 30 | | ns |
| t _{RP} | 30 | | 30 | | 36 | | ns |
| t _{WR} | 1 | | 1 | | 1 | | t _{CK} |

*CAS latency indicated in parentheses.

- NOTE:**
1. For this example, the burst length = 4, and the WRITE burst is followed by "manual" PRECHARGE.
 2. A8 and A9 = "Don't Care."
 3. DQM represents DQML and DQMH. DQML controls the lower byte, and DQMH controls the upper byte.

WRITE – WITH AUTO PRECHARGE 1



TIMING PARAMETERS

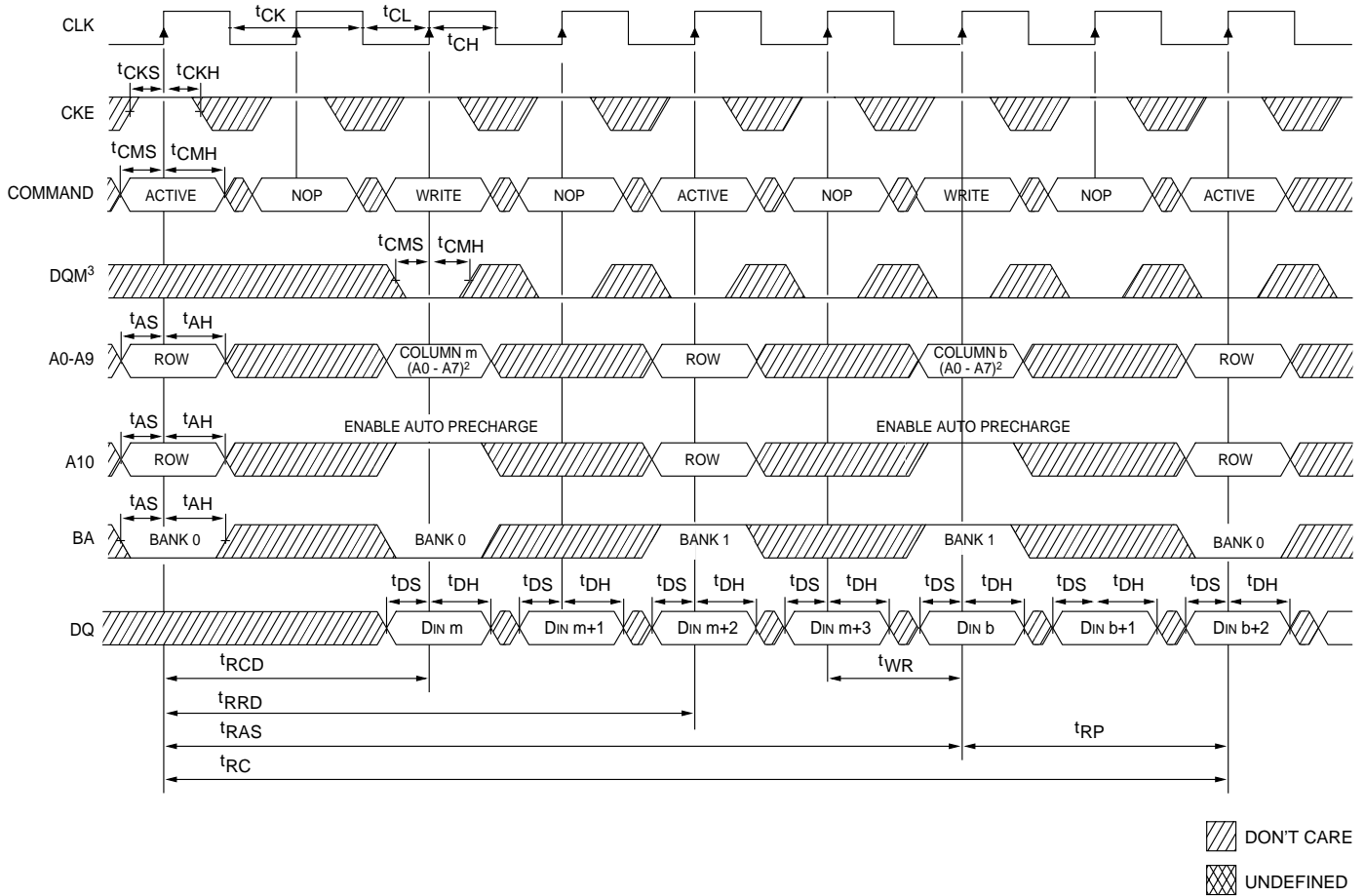
| SYMBOL* | -8A | | -10 | | -12 | | UNITS |
|---------------------|-----|-----|-----|-----|-----|-----|-------|
| | MIN | MAX | MIN | MAX | MIN | MAX | |
| t _{AH} | 1 | | 1 | | 1 | | ns |
| t _{AS} | 2 | | 3 | | 3 | | ns |
| t _{CH} | 3 | | 3.5 | | 3.5 | | ns |
| t _{CL} | 3 | | 3.5 | | 3.5 | | ns |
| t _{CK} (3) | 8 | | 10 | | 12 | | ns |
| t _{CK} (2) | 13 | | 15 | | 15 | | ns |
| t _{CK} (1) | 30 | | 30 | | 30 | | ns |
| t _{CKH} | 1 | | 1 | | 1 | | ns |
| t _{CKS} | 2 | | 3 | | 3 | | ns |

| SYMBOL* | -8A | | -10 | | -12 | | UNITS |
|------------------|-----|------|-----|------|-----|------|-----------------|
| | MIN | MAX | MIN | MAX | MIN | MAX | |
| t _{CMH} | 1 | | 1 | | 1 | | ns |
| t _{CMS} | 2 | | 3 | | 3 | | ns |
| t _{DH} | 1 | | 1 | | 1 | | ns |
| t _{DS} | 2 | | 3 | | 3 | | ns |
| t _{RAS} | 50 | 120K | 60 | 120K | 72 | 120K | ns |
| t _{RC} | 80 | | 90 | | 105 | | ns |
| t _{RCD} | 30 | | 30 | | 30 | | ns |
| t _{RP} | 30 | | 30 | | 36 | | ns |
| t _{WR} | 1 | | 1 | | 1 | | t _{CK} |

*CAS latency indicated in parentheses.

- NOTE:**
1. For this example, the burst length = 4.
 2. A8 and A9 = "Don't Care."
 3. DQM represents DQML and DQMH. DQML controls the lower byte, and DQMH controls the upper byte.

ALTERNATING BANK WRITE ACCESSES 1



TIMING PARAMETERS

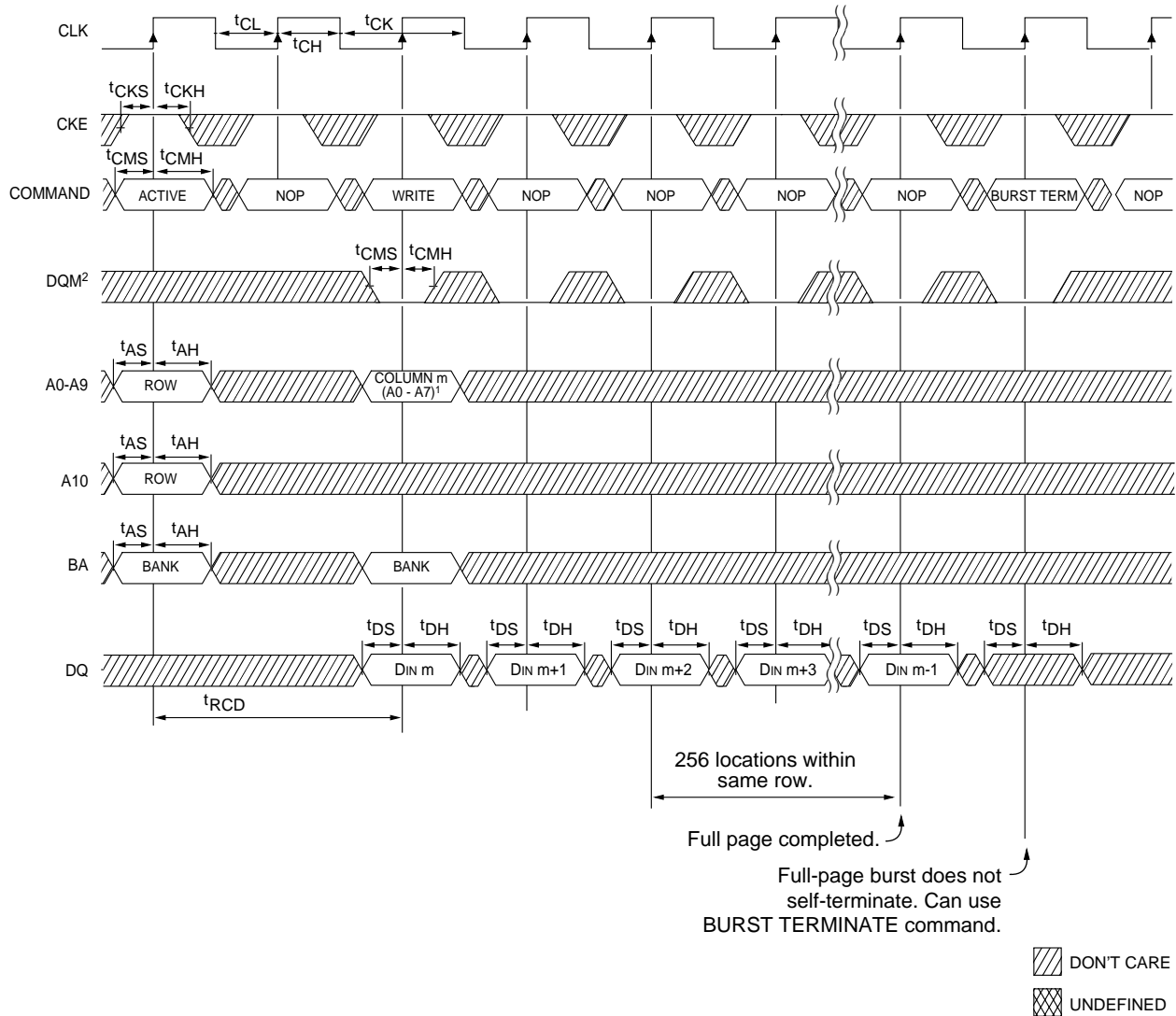
| SYMBOL* | -8A | | -10 | | -12 | | UNITS |
|---------------------|-----|-----|-----|-----|-----|-----|-------|
| | MIN | MAX | MIN | MAX | MIN | MAX | |
| t _{AH} | 1 | | 1 | | 1 | | ns |
| t _{AS} | 2 | | 3 | | 3 | | ns |
| t _{CH} | 3 | | 3.5 | | 3.5 | | ns |
| t _{CL} | 3 | | 3.5 | | 3.5 | | ns |
| t _{CK} (3) | 8 | | 10 | | 12 | | ns |
| t _{CK} (2) | 13 | | 15 | | 15 | | ns |
| t _{CK} (1) | 30 | | 30 | | 30 | | ns |
| t _{CKH} | 1 | | 1 | | 1 | | ns |
| t _{CKS} | 2 | | 3 | | 3 | | ns |
| t _{CMH} | 1 | | 1 | | 1 | | ns |

| SYMBOL* | -8A | | -10 | | -12 | | UNITS |
|------------------|-----|------|-----|------|-----|------|-----------------|
| | MIN | MAX | MIN | MAX | MIN | MAX | |
| t _{CMS} | 2 | | 3 | | 3 | | ns |
| t _{DH} | 1 | | 1 | | 1 | | ns |
| t _{DS} | 2 | | 3 | | 3 | | ns |
| t _{RAS} | 50 | 120K | 60 | 120K | 72 | 120K | ns |
| t _{RC} | 80 | | 90 | | 105 | | ns |
| t _{RCD} | 30 | | 30 | | 30 | | ns |
| t _{RP} | 30 | | 30 | | 36 | | ns |
| t _{RRD} | 20 | | 20 | | 20 | | ns |
| t _{WR} | 1 | | 1 | | 1 | | t _{CK} |

*CAS latency indicated in parentheses.

- NOTE:**
1. For this example, the burst length = 4.
 2. A8 and A9 = "Don't Care."
 3. DQM represents DQML and DQMH. DQML controls the lower byte, and DQMH controls the upper byte.

WRITE – FULL-PAGE BURST



TIMING PARAMETERS

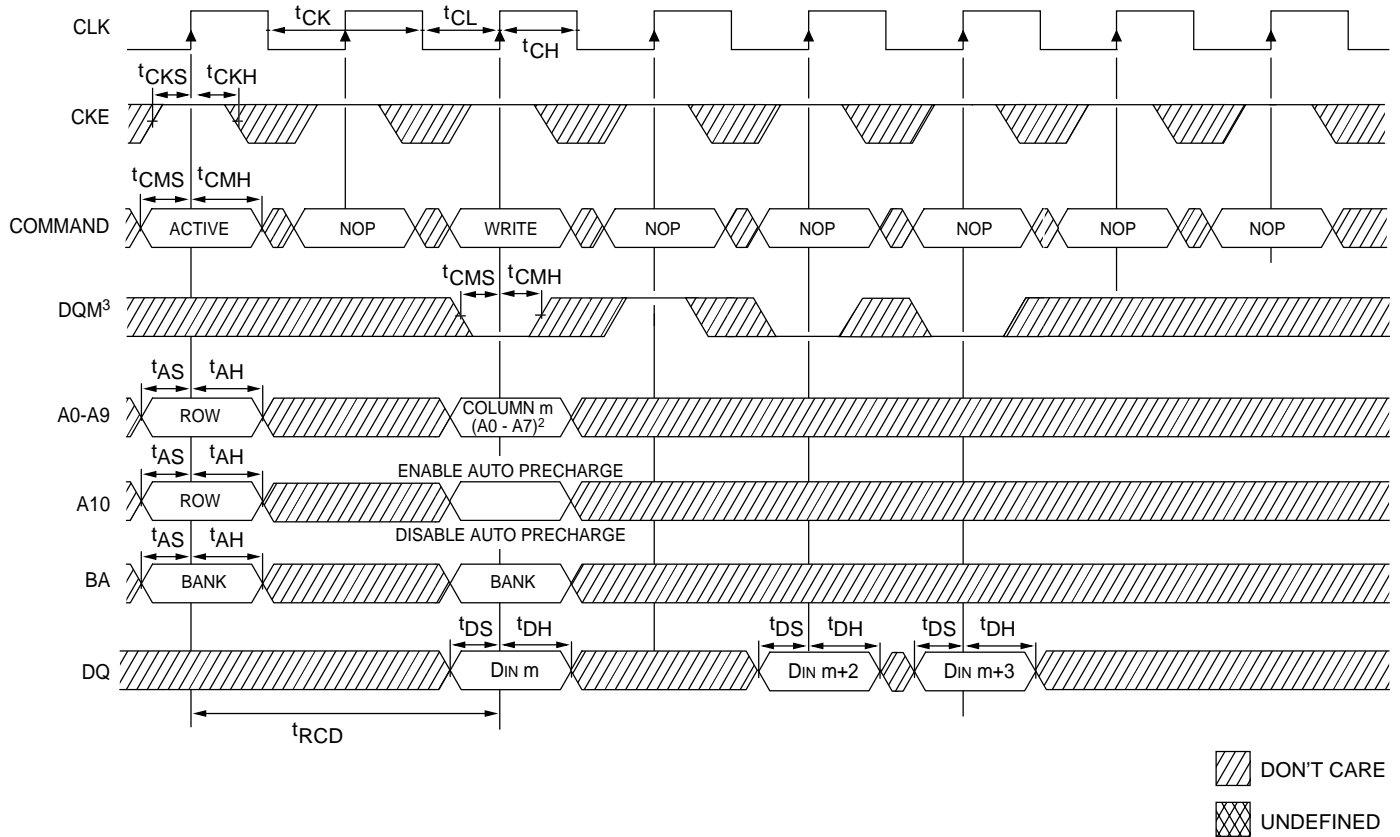
| SYMBOL* | -8A | | -10 | | -12 | | UNITS |
|--------------|-----|-----|-----|-----|-----|-----|-------|
| | MIN | MAX | MIN | MAX | MIN | MAX | |
| t_{AH} | 1 | | 1 | | 1 | | ns |
| t_{AS} | 2 | | 3 | | 3 | | ns |
| t_{CH} | 3 | | 3.5 | | 3.5 | | ns |
| t_{CL} | 3 | | 3.5 | | 3.5 | | ns |
| t_{CK} (3) | 8 | | 10 | | 12 | | ns |
| t_{CK} (2) | 13 | | 15 | | 15 | | ns |
| t_{CK} (1) | 30 | | 30 | | 30 | | ns |

| SYMBOL* | -8A | | -10 | | -12 | | UNITS |
|-----------|-----|-----|-----|-----|-----|-----|-------|
| | MIN | MAX | MIN | MAX | MIN | MAX | |
| t_{CKH} | 1 | | 1 | | 1 | | ns |
| t_{CKS} | 2 | | 3 | | 3 | | ns |
| t_{CMH} | 1 | | 1 | | 1 | | ns |
| t_{CMS} | 2 | | 3 | | 3 | | ns |
| t_{DH} | 1 | | 1 | | 1 | | ns |
| t_{DS} | 2 | | 3 | | 3 | | ns |
| t_{RCD} | 30 | | 30 | | 30 | | ns |

*CAS latency indicated in parentheses.

NOTE: 1. A8 and A9 = "Don't Care."
2. DQM represents DQML and DQMH. DQML controls the lower byte, and DQMH controls the upper byte.

WRITE – DQM OPERATION ¹



TIMING PARAMETERS

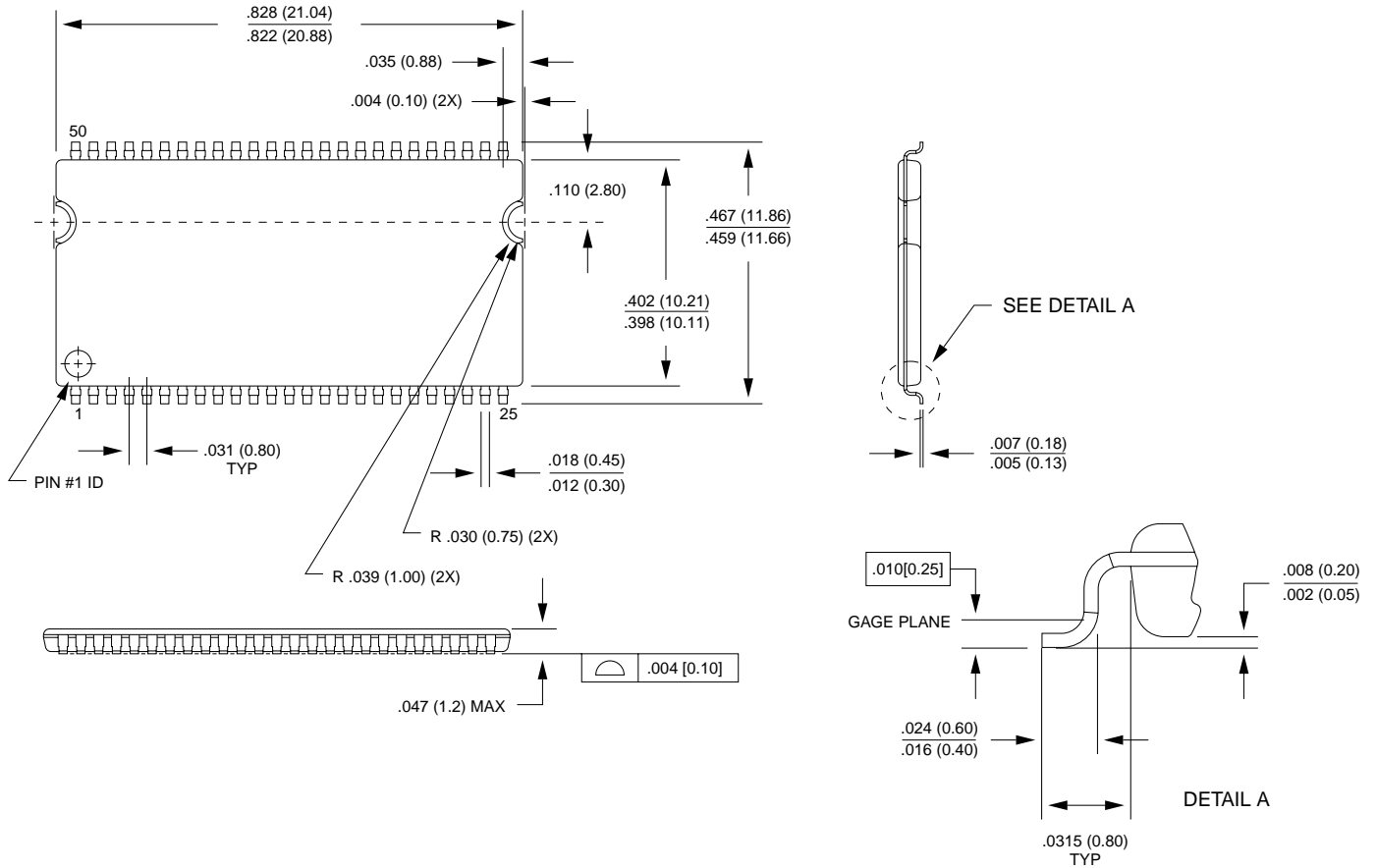
| SYMBOL* | -8A | | -10 | | -12 | | UNITS |
|---------------------|-----|-----|-----|-----|-----|-----|-------|
| | MIN | MAX | MIN | MAX | MIN | MAX | |
| t _{AH} | 1 | | 1 | | 1 | | ns |
| t _{AS} | 2 | | 3 | | 3 | | ns |
| t _{CH} | 3 | | 3.5 | | 3.5 | | ns |
| t _{CL} | 3 | | 3.5 | | 3.5 | | ns |
| t _{CK} (3) | 8 | | 10 | | 12 | | ns |
| t _{CK} (2) | 13 | | 15 | | 15 | | ns |
| t _{CK} (1) | 30 | | 30 | | 30 | | ns |

| SYMBOL* | -8A | | -10 | | -12 | | UNITS |
|------------------|-----|-----|-----|-----|-----|-----|-------|
| | MIN | MAX | MIN | MAX | MIN | MAX | |
| t _{CKH} | 1 | | 1 | | 1 | | ns |
| t _{CKS} | 2 | | 3 | | 3 | | ns |
| t _{CMH} | 1 | | 1 | | 1 | | ns |
| t _{CMS} | 2 | | 3 | | 3 | | ns |
| t _{DH} | 1 | | 1 | | 1 | | ns |
| t _{DS} | 2 | | 3 | | 3 | | ns |
| t _{RCD} | 30 | | 30 | | 30 | | ns |

*CAS latency indicated in parentheses.

- NOTE:**
1. For this example, the burst length = 4.
 2. A8 and A9 = "Don't Care."
 3. DQM represents DQML and DQMH. DQML controls the lower byte, and DQMH controls the upper byte.

50-PIN PLASTIC TSOP (400 mil)



- NOTE:**
1. All dimensions in inches (millimeters) $\frac{\text{MAX}}{\text{MIN}}$ or typical where noted.
 2. Package width and length do not include mold protrusion; allowable mold protrusion is .01" per side.