

PRELIMINARY

August 1996

COP87L84RG

8-Bit One-Time Programmable (OTP) Microcontroller with 32 Kbytes of Program Memory

General Description

The COP87L84RG is a member of the COP8TM OTP microcontroller family. It is pin and software compatible to the mask ROM COP888EG product family.

(Continued)

Key Features

- Full duplex UART
- Three 16-bit timers, each with two 16-bit registers supporting:
 - Processor independent PWM mode
 - External event counter mode
 - Input capture mode
- 32 kbytes on-board EPROM with security feature

Note: Mask ROMed devices with equivalent on-chip features and program memory sizes of 4k and 8k are available (see Table I).

■ 512 bytes on-board RAM

Additional Peripheral Features

- Idle timer
- Multi-Input Wake Up (MIWU) with optional interrupts (8)
- One analog comparator
- WATCHDOG™ and clock monitor logic
- MICROWIRE/PLUS™ serial I/O

I/O Features

- Memory mapped I/O
- Software selectable I/O options (TRI-STATE® output, push-pull output, weak pull-up input, high impedance input

- Schmitt trigger inputs on ports G and L
- Packages:
 - 28 DIP or SO each with 24 I/O pins

CPU/Instruction Set Features

- 1 µs instruction cycle time
- Fourteen multi-source vectored interrupts servicing
 - External interrupt
 - Idle timer T0
 - Two timers (each with 2 interrupts)
 - MICROWIRE/PLUS
 - Multi-Input Wake Up
 - Software trap
 - UART (2)
 - Default VIS (default interrupt)
- Versatile instruction set with true bit manipulation
- 8-bit Stack Pointer SP (stack in RAM)
- Two 8-bit register indirect data memory pointers (B and X)

Fully Static CMOS

- Two power saving modes: HALT and IDLE
- Single supply operation: 2.7V-5.5V
- Temperature range: -40°C to +85°C

Development Support

- Emulation device for the COP884EG
- Real time emulation and full program debug offered by MetaLink Development Systems

Block Diagram

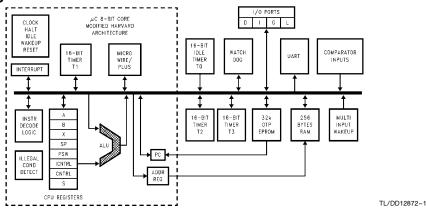


FIGURE 1. Block Diagram

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TL/DD12872

RRD-B30M96/Printed in U. S.

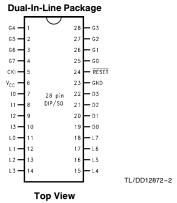
General Description (Continued)

The device is a fully static part, fabricated using double-metal silicon gate microCMOS technology. Features include an 8-bit memory mapped architecture, MICROWIRE/PLUS™ serial I/O, three 16-bit timer/counters supporting three modes (Processor Independent PWM generation, External Event counter, and Input Capture mode capabilities), full duplex UART and two comparators. Each I/O pin has software selectable configurations. The devices operates over a voltage range of 2.7V to 5.5V. High throughput is achieved with an efficient, regular instruction set operating at a maximum rate of 1 μs per instruction.

TABLE I. COP888CG/EG Family Members

Device	ROM/EPROM (Bytes)	RAM (Bytes)	Key Common Features
COP888CG	4k ROM	192	3 Timers
COP888EG	8k ROM	256	UART
COP87L84EG	8k OTP EPROM	256	1 Comparator
COP87L84RG	32k OTP EPROM	256	

Note: -X Crystal Oscillator -E Halt Enable



Order Number COP87L84RGM-XE or COP87L84RGN-XE See NS Package Number M28B or N28B

FIGURE 2. Connection Diagrams

Connection Diagrams (Continued)

Pinouts for 28-Pin Package

Port	Туре	Alt. Fun	Alt. Fun	28-Pin DIP/SO
LO	1/0	MIWU		11
L1	1/0	MIWU	CKX	12
L2	1/0	MIWU	TDX	13
L3	1/0	MIWU	RDX	14
L4	1/0	MIWU	T2A	15
L5	1/0	MIWU	T2B	16
L6	1/0	MIWU	T3A	17
L7	1/0	MIWU	ТЗВ	18
G0	1/0	INT		25
G1	WDOUT			26
G2	1/0	T1B		27
G3	1/0	T1A		28
G4	1/0	SO		1
G5	1/0	SK		2
G6	1	SI		3
G7	I/CKO	HALT Restart		4
D0	0			19
D1	0			20
D2	0			21
D3	0			22
10	ı			7
11	1	COMP1IN-		8
12	1	COMP1IN+		9
13	1	COMP1OUT		10
V _{CC}				6
GND				23
CKI				5
RESET				24

Absolute Maximum Ratings (Note)

If Military/Aerospace specified devices are required, please contact the National Semiconductor Sales Office/Distributors for availability and specifications.

Supply Voltage (V_{CC}) 7V Voltage at Any Pin -0.3V to V_{CC} +0.3V Total Current into V_{CC} Pin (Source) 100 mA

Total Current out of GND Pin (Sink)
Storage Temperature Range

110 mA -65°C to +140°C

Note: Absolute maximum ratings indicate limits beyond which damage to the device may occur. DC and AC electrical specifications are not ensured when operating the device at absolute maximum ratings.

DC Electrical Characteristics $-40^{\circ}C \le T_{A} \le +85^{\circ}C$ unless otherwise specified

Parameter	Conditions	Min	Тур	Max	Units
Operating Voltage		2.7		5.5	V
Power Supply Ripple (Note 1)	Peak-to-Peak			0.1 V _{CC}	V
Supply Current (Note 2) CKI = 10 MHz CKI = 4 MHz	$V_{CC} = 5.5V, t_{c} = 1 \mu s$ $V_{CC} = 4.0V, t_{c} = 2.5 \mu s$			16.5 6.5	mA mA
HALT Current (Note 3)	$V_{CC} = 5.5V$, CKI = 0 MHz $V_{CC} = 4.0V$, CKI = 0 MHz			. 12 8	μA μA
IDLE Current CKI = 10 MHz CKI = 1 MHz	$V_{CC} = 5.5V, t_{C} = 1 \mu s$ $V_{CC} = 4.0V, t_{C} = 10 \mu s$			3.5 0.7	mA mA
Input Levels RESET Logic High Logic Low CKI (External and Crystal Osc. Modes) Logic High Logic Low All Other Inputs		0.8 V _{CC}		0.2 V _{CC}	V V V
Logic High Logic Low		0.7 V _{CC}		0.2 V _{CC}	V V
Hi-Z Input Leakage	V _{CC} = 5.5V	-2		+2	μΑ
Input Pullup Current	V _{CC} = 5.5V	40		250	μΑ
G and L Port Input Hysteresis			0.05 V _{CC}	0.35 V _{CC}	V
Output Current Levels D Outputs Source Sink (Note 4) All Others Source (Weak Pull-Up Mode) Source (Push-Pull Mode) Sink (Push-Pull Mode)	$\begin{aligned} & \text{V}_{\text{CC}} = 4.5\text{V}, \text{V}_{\text{OH}} = 3.3\text{V} \\ & \text{V}_{\text{CC}} = 4.5\text{V}, \text{V}_{\text{OL}} = 1\text{V} \end{aligned}$ $\begin{aligned} & \text{V}_{\text{CC}} = 4.5\text{V}, \text{V}_{\text{OH}} = 2.7\text{V} \\ & \text{V}_{\text{CC}} = 4.5\text{V}, \text{V}_{\text{OH}} = 3.3\text{V} \\ & \text{V}_{\text{CC}} = 4.5\text{V}, \text{V}_{\text{OL}} = 0.4\text{V} \end{aligned}$	0.4 10 10 0.4 1.6		100	mA mA μA mA mA
TRI-STATE Leakage	$V_{CC} = 5.5V$	-2		+2	μΑ
Allowable Sink/Source Current per Pin D Outputs (Sink) All others				15 3	mA mA
Maximum Input Current without Latchup (Note 5)	T _A = 25°C			±100	mA
RAM Retention Voltage, V _r	500 ns Rise and Fall Time (Min)	2			V
Input Capacitance				7	pF
Load Capacitance on D2				1000	pF

Note 1: Rate of voltage change must be less then 0.5 V/ms.

Note 5: Pins G6 and $\overline{\text{RESET}}$ are designed with a high voltage input network for factory testing. These pins allow input voltages greater than V_{CC} and the pins will have sink current to V_{CC} when biased at voltages greater than V_{CC} (the pins do not have source current when biased at a voltage below V_{CC}). The effective resistance to V_{CC} is 750 Ω (typical). These two pins will not latch up. The voltage at the pins must be limited to less than 14V.

Note 2: Supply current is measured after running 2000 cycles with a square wave CKI input, CKO open, inputs at rails and outputs open.

Note 3: The HALT mode will stop CKI from oscillating in the RC and the Crystal configurations by bringing CKI high. Test conditions: All inputs tied to V_{CC}, L and G ports in the TRI-STATE mode and tied to ground, all outputs low and tied to ground. The clock monitor is disabled.

Note 4: The user must guarantee that D2 pin does not source more than 10 mA during RESET. If D2 sources more than 10 mA during reset, the device will go into programming mode.

Parameter	Conditions	Min	Тур	Max	Units
Instruction Cycle Time (t _c)					
Crystal, Resonator,		1		DC	μs
R/C Oscillator		3		DC	μS
Inputs					
t _{SETUP}		200			ns
thold		60			ns
Output Propagation Delay	$R_L = 2.2k, C_L = 100 pF$				
t _{PD1} , t _{PD0}					
SO, SK				0.7	μs
All Others				1	μs
MICROWIRE™ Setup Time (t _{UWS})		20			ns
MICROWIRE Hold Time (t _{UWH})		56			ns
MICROWIRE Output Propagation Delay (t _{UPD})				220	ns
Input Pulse Width					
Interrupt Input High Time		1			t _c
Interrupt Input Low Time		1			t _c
Timer Input High Time		1			t _c
Timer Input Low Time		1			t _c
Reset Pulse Width		1			цS

Comparator AC and DC Characteristics $V_{CC} = 5V, T_A = 25^{\circ}C$

Parameter	Conditions	Min	Тур	Max	Units
Input Offset Voltage	$0.4V \leq V_{IN} \leq V_{CC} - 1.5V$		±10	±25	mV
Input Common Mode Voltage Range		0.4		V _{CC} - 1.5	V
Low Level Output Current	$V_{OL} = 0.4V$	1.6			mA
High Level Output Current	V _{OH} = 4.6V	1.6			mA
Comparator DC Supply Current (When Enabled)				250	μΑ
Response Time	TBD mV Step, TBD mV Overdrive, 100 pF Load		1		μs

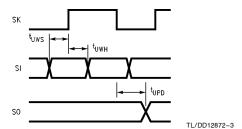


FIGURE 3. MICROWIRE/PLUS Timing

Pin Descriptions

 $V_{\mbox{\footnotesize CC}}$ and GND are the power supply pins.

CKI is the clock input. This can come from an R/C generated oscillator, or a crystal oscillator (in conjunction with CKO). See Oscillator Description section.

RESET is the master reset input. See Reset Description section

The device contains two bidirectional 8-bit I/O ports (G and L), where each individual bit may be independently configured as an input (Schmitt trigger inputs on ports L and G), output or TRI-STATE under program control. Three data memory address locations are allocated for each of these I/O ports. Each I/O port has two associated 8-bit memory mapped registers, the CONFIGURATION register and the output DATA register. A memory mapped address is also reserved for the input pins of each I/O port. (See the memory map for the various addresses associated with the I/O ports.) Figure 4 shows the I/O port configurations. The DATA and CONFIGURATION registers allow for each port bit to be individually configured under software control as shown below:

CONFIGURATION Register	DATA Register	Port Set-Up
0	0	Hi-Z Input (TRI-STATE Output)
0	1	Input with Weak Pull-Up
1	0	Push-Pull Zero Output
1	1	Push-Pull One Output

PORT L is an 8-bit I/O port. All L-pins have Schmitt triggers on the inputs.

Port L supports Multi-Input Wake Up (MIWU) on all eight pins. L1 is used for the UART external clock. L2 and L3 are used for the UART transmit and receive. L4 and L5 are used for the timer input functions T2A and T2B. L6 and L7 are used for the timer input functions T3A and T3B.

Port L has the following alternate features:

- LO MIWU
- L1 MIWU or CKX
- L2 MIWU or TDX
- L3 MIWU or RDX
- L4 MIWU or T2A
- L5 MIWU or T2B
- L6 MIWU or T3A
- L7 MIWU or T3B

Port G is an 8-bit port with 5 I/O pins (G0, G2-G5), an input pin (G6), and two dedicated output pins (G1 and G7). Pins G0 and G2-G6 all have Schmitt Triggers on their inputs. Pin G1 serves as the dedicated WDOUT WATCHDOG output, while pin G7 is either input or output depending on the oscillator mask option selected. With the crystal oscillator option selected, G7 serves as the dedicated output pin for the CKO clock output. With the single-pin R/C oscillator mask option selected, G7 serves as a general purpose input pin but is also used to bring the device out of HALT mode with a low to high transition on G7. There are two registers associated with the G Port, a data register and a configuration register. Therefore, each of the 5 I/O bits (G0, G2-G5) can be individually configured under software control.

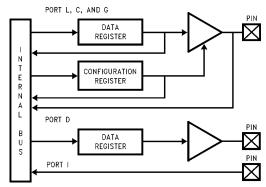


FIGURE 4. I/O Port Configurations

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Pin Descriptions (Continued)

Since G6 is an input only pin and G7 is the dedicated CKO clock output pin (crystal clock option) or general purpose input (R/C clock option), the associated bits in the data and configuration registers for G6 and G7 are used for special purpose functions as outlined below. Reading the G6 and G7 data bits will return zeros.

Note that the chip will be placed in the HALT mode by writing a "1" to bit 7 of the Port G Data Register. Similarly the chip will be placed in the IDLE mode by writing a "1" to bit 6 of the Port G Data Register.

Writing a "1" to bit 6 of the Port G Configuration Register enables the MICROWIRE/PLUS to operate with the alternate phase of the SK clock. The G7 configuration bit, if set high, enables the clock start up delay after HALT when the R/C clock configuration is used.

	Config Reg.	Data Reg.
G7	CLKDLY	HALT
G6	Alternate SK	IDLE

Port G has the following alternate features:

- G0 INTR (External Interrupt Input)
- G2 T1B (Timer T1 Capture Input)
- G3 T1A (Timer T1 I/O)
- G4 SO (MICROWIRE Serial Data Output)
- G5 SK (MICROWIRE Serial Clock)
- G6 SI (MICROWIRE Serial Data Input)

Port G has the following dedicated functions:

- G1 WDOUT WATCHDOG and/or Clock Monitor dedicated output
- G7 CKO Oscillator dedicated output or general purpose input

Port I is a four-bit Hi-Z input port.

Port I1-I3 are used for Comparator 1. Port I4-I6 are used for Comparator 2.

The Port I has the following alternate features.

- I1 COMP1-IN (Comparator 1 Negative Input)
- I2 COMP1+IN (Comparator 1 Positive Input)
- I3 COMP1OUT (Comparator 1 Output)

Port D is a recreated 4-bit output port that is preset high when RESET goes low. D port recreation is one clock cycle behind normal port timing. The user can tie two or more D port outputs (except D2) together in order to get a higher drive.

Functional Description

The architecture of the device is modified Harvard architecture. With the Harvard architecture, the control store program memory (ROM) is separated from the data store memory (RAM). Both ROM and RAM have their own separate addressing space with separate address buses. The architecture, though based on Harvard architecture, permits transfer of data from ROM to RAM.

CPU REGISTERS

The CPU can do an 8-bit addition, subtraction, logical or shift operation in one instruction ($t_{\rm c}$) cycle time.

There are six CPU registers:

A is the 8-bit Accumulator Register

PC is the 15-bit Program Counter Register

PU is the upper 7 bits of the program counter (PC)

PL is the lower 8 bits of the program counter (PC)

B is an 8-bit RAM address pointer, which can be optionally

post auto incremented or decremented.

X is an 8-bit alternate RAM address pointer, which can be optionally post auto incremented or decremented.

SP is the 8-bit stack pointer, which points to the subroutine/interrupt stack (in RAM). The SP is initialized to RAM address 06F with reset.

S is the 8-bit Data Segment Address Register used to extend the lower half of the address range (00 to 7F) into 256 data segments of 128 bytes each.

All the CPU registers are memory mapped with the exception of the Accumulator (A) and the Program Counter (PC).

PROGRAM MEMORY

The program memory consists of 32 kbytes of OTP EPROM. These bytes may hold program instructions or constant data (data tables for the LAID instruction, jump vectors for the JID instruction, and interrupt vectors for the VIS instruction). The program memory is addressed by the 15-bit program counter (PC). All interrupts in the devices vector to program memory location 0FF Hex.

The device can be configured to inhibit external reads of the program memory. This is done by programming the Security Byte.

Note: Mask ROMed devices with equivalent on-chip features and program memory sizes of 4k and 8k are available.

SECURITY FEATURE

The program memory array has an associate Security Byte that is located outside of the program address range. This byte can be addressed only from programming mode by a programmer tool.

Security is an optional feature and can only be asserted after the memory array has been programmed and verified. A secured part will read all 00(hex) by a programmer. The part will fail Blank Check and will fail Verify operations. A Read operaiton will fill the programmer's memory with 00(hex). The Security Byte itself is always readable with value of 00(hex) if unsecure and FF(hex) if secure.

DATA MEMORY

The data memory address space includes the on-chip RAM and data registers, the I/O registers (Configuration, Data and Pin), the control registers, the MICROWIRE/PLUS SIO shift register, and the various registers, and counters associated with the timers (with the exception of the IDLE timer). Data memory is addressed directly by the instruction or indirectly by the B, X, SP pointers and S register.

Functional Description (Continued)

The data memory consists of 256 bytes of RAM. Sixteen bytes of RAM are mapped as "registers" at addresses 0F0 to 0FF Hex. These registers can be loaded immediately, and also decremented and tested with the DRSZ (decrement register and skip if zero) instruction. The memory pointer registers X, SP, B and S are memory mapped into this space at address locations 0FC to 0FF Hex respectively, with the other registers being available for general usage.

The instruction set permits any bit in memory to be set, reset or tested. All I/O and registers (except A and PC) are memory mapped; therefore, I/O bits and register bits can be directly and individually set, reset and tested. The accumulator (A) bits can also be directly and individually tested.

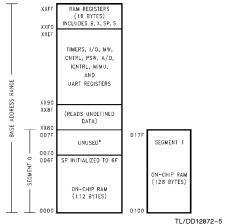
Note: RAM contents are undefined upon power-up.

Data Memory Segment RAM Extension

Data memory address 0FF is used as a memory mapped location for the Data Segment Address Register (S).

The data store memory is either addressed directly by a single byte address within the instruction, or indirectly relative to the reference of the B. X. or SP pointers (each contains a single-byte address). This single-byte address allows an addressing range of 256 locations from 00 to FF hex. The upper bit of this single-byte address divides the data store memory into two separate sections as outlined previously. With the exception of the RAM register memory from address locations 00F0 to 00FF, all RAM memory is memory mapped with the upper bit of the single-byte address being equal to zero. This allows the upper bit of the single-byte address to determine whether or not the base address range (from 0000 to 00FF) is extended. If this upper bit equals one (representing address range 0080 to 00FF), then address extension does not take place. Alternatively, if this upper bit equals zero, then the data segment extension register S is used to extend the base address range (from 0000 to 007F) from XX00 to XX7F, where XX represents the 8 bits from the S register. Thus the 128-byte data segment extensions are located from addresses 0100 to 017F for data segment 1, 0200 to 027F for data segment 2, etc., up to FF00 to FF7F for data segment 255. The base address range from 0000 to 007F represents data segment 0.

Figure 5 illustrates how the S register data memory extension is used in extending the lower half of the base address range (00 to 7F hex) into 256 data segments of 128 bytes each, with a total addressing range of 32 kbytes from XX00 to XX7F. This organization allows a total of 256 data segments of 128 bytes each with an additional upper base segment of 128 bytes. Furthermore, all addressing modes are available for all data segments. The S register must be changed under program control to move from one data segment (128 bytes) to another. However, the upper base segment (containing the 16 memory registers, I/O registers, control registers, etc.) is always available regardless of the contents of the S register, since the upper base segment (address range 0080 to 00FF) is independent of data segment extension.



*Reads as all ones.

FIGURE 5. RAM Organization

The instructions that utilize the stack pointer (SP) always reference the stack as part of the base segment (Segment 0), regardless of the contents of the S register. The S register is not changed by these instructions. Consequently, the stack (used with subroutine linkage and interrupts) is always located in the base segment. The stack pointer will be intitialized to point at data memory location 006F as a result of reset

The 128 bytes of RAM contained in the base segment are split between the lower and upper base segments. The first 116 bytes of RAM are resident from address 0000 to 006F in the lower base segment, while the remaining 16 bytes of RAM represent the 16 data memory registers located at addresses 00F0 to 00FF of the upper base segment. No RAM is located at the upper sixteen addresses (0070 to 007F) of the lower base segment.

Additional RAM beyond these initial 128 bytes, however, will always be memory mapped in groups of 128 bytes (or less) at the data segment address extensions (XX00 to XX7F) of the lower base segment. The additional 128 bytes of RAM are memory mapped at address locations 0100 to 017F hex.

Reset

The RESET input when pulled low initializes the microcontroller. Initialization will occur whenever the RESET input is pulled low. Upon initialization, the data and configuration registers for ports L, G and C are cleared, resulting in these

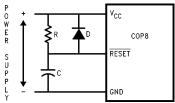
Reset (Continued)

Ports being initialized to the TRI-STATE mode. Pin G1 of the G Port is an exception (as noted below) since pin G1 is dedicated as the WATCHDOG and/or Clock Monitor error output pin. Port D is set high. The PC, PSW, ICNTRL, CNTRL, T2CNTRL and T3CNTRL control registers are cleared. The UART registers PSR, ENU (except that TBMT bit is set), ENUR and ENUI are cleared. The Comparator Select Register is cleared. The S register is initialized to zero. The Multi-Input Wake Up registers WKEN, WKEDG and WKPND are cleared. The stack pointer, SP, is initialized to 6F Hex.

The device comes out of reset with both the WATCHDOG logic and the Clock Monitor detector armed, with the WATCHDOG service window bits set and the Clock Monitor bit set. The WATCHDOG and Clock Monitor circuits are inhibited during reset. The WATCHDOG service window bits being initialized high default to the maximum WATCHDOG service window of 64k $t_{\rm c}$ clock cycles. The Clock Monitor bit being initialized high will cause a Clock Monitor error following reset if the clock has not reached the minimum specified frequency at the termination of reset. A Clock Monitor error will cause an active low error output on pin G1. This error output will continue until 16 $t_{\rm c}$ –32 $t_{\rm c}$ clock cycles following the clock frequency reaching the minimum specified value, at which time the G1 output will enter the TRI-STATE mode.

The external RC network shown in Figure 6 should be used to ensure that the $\overline{\text{RESET}}$ pin is held low until the power supply to the chip stabilizes.

Note: Continual state of reset will cause the device to draw excessive current.



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 $RC > 5 \times Power Supply Rise Time$

FIGURE 6. Recommended Reset Circuit

Oscillator Circuits

The chip can be driven by a clock input on the CKI input pin which can be between DC and 10 MHz. The CKO output clock is on pin G7 (crystal configuration). The CKI input frequency is divided down by 10 to produce the instruction cycle clock $(1/t_c)$.

Figure 7 shows the Crystal and R/C diagrams.

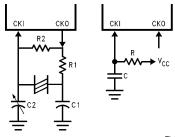


FIGURE 7. Crystal and R/C Oscillator Diagrams

CRYSTAL OSCILLATOR

CKI and CKO can be connected to make a closed loop crystal (or resonator) controlled oscillator.

Table I shows the component values required for various standard crystal values.

TABLE I. Crystal Oscillator Configuration, $T_A = 25^{\circ}C$

R1 (kΩ)	R2 (MΩ)	C1 (pF)	C2 (pF)	CKI Freq (MHz)	Conditions
0	1	30	30-36	10	$V_{CC} = 5V$
0	1	30	30-36	4	$V_{CC} = 5V$
0	1	200	100-150	0.455	$V_{CC} = 5V$

R/C OSCILLATOR

By selecting CKI as a single pin oscillator input, a single pin R/C oscillator circuit can be connected to it. CKO is available as a general purpose input, and/or HALT restart pin.

Table II shows the variation in the oscillator frequencies as functions of the component (R and C) values.

TABLE II. R/C Oscillator Configuration, $T_A=25^{\circ}C$

R (kΩ)	C (pF)	CKI Freq (MHz)	Instr. Cycle (μs)	Conditions
3.3	82	2.2-2.7	3.7-4.6	$V_{CC} = 5V$
5.6	100	1.1-1.3	7.4-9.0	$V_{CC} = 5V$
6.8	100	0.9-1.1	8.8-10.8	$V_{CC} = 5V$

Note: $3k \le R \le 200k$ $50 \text{ pF} \le C \le 200 \text{ pF}$

Control Registers

CNTRL Register (Address X'00EE)

The Timer1 (T1) and MICROWIRE/PLUS control register contains the following bits:

SL1 & SL0 Select the MICROWIRE/PLUS clock divide by (00 = 2, 01 = 4, 1x = 8)

IEDG External interrupt edge polarity select

(0 = Rising edge, 1 = Falling edge)

MSEL Selects G5 and G4 as MICROWIRE/PLUS

signals

SK and SO respectively

T1C0 Timer T1 Start/Stop control in timer

modes 1 and 2

Timer T1 Underflow Interrupt Pending Flag in

timer mode 3

T1C1 Timer T1 mode control bit
T1C2 Timer T1 mode control bit

T1C2 Timer T1 mode control bit
T1C3 Timer T1 mode control bit

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Control Registers (Continued)

PSW Register (Address X'00EF)

The PSW register contains the following select bits:

GIE Global interrupt enable (enables interrupts)

EXEN Enable external interrupt

BUSY MICROWIRE/PLUS busy shifting flag

EXPND External interrupt pending

T1ENA Timer T1 Interrupt Enable for Timer Underflow

or T1A Input capture edge

T1PNDA Timer T1 Interrupt Pending Flag (Autoreload RA

in mode 1, T1 Underflow in Mode 2, T1A cap-

ture edge in mode 3)

C Carry Flag
HC Half Carry Flag

HC C T1PNDA T1ENA EXPND BUSY EXEN GIE

The Half-Carry bit is also affected by all the instructions that affect the Carry flag. The SC (Set Carry) and RC (Reset Carry) instructions will respectively set or clear both the carry flags. In addition to the SC and RC instructions, ADC, SUBC, RRC and RLC instructions affect the carry and Half Carry flags.

ICNTRL Register (Address X'00E8)

The ICNTRL register contains the following bits:

T1ENB Timer T1 Interrupt Enable for T1B Input capture edge

T1PNDB Timer T1 Interrupt Pending Flag for T1B cap-

ture edge

WEN Enable MICROWIRE/PLUS interrupt

WPND MICROWIRE/PLUS interrupt pending
T0EN Timer T0 Interrupt Enable (Bit 12 toggle)

T0PND Timer T0 Interrupt pending

LPEN L Port Interrupt Enable (Multi-Input Wake Up/

Interrupt)

Bit 7 could be used as a flag

Unused	LPEN	T0PND	T0EN	WPND	WEN	T1PNDB	T1ENB
Bit 7							Bit 0

T2CNTRL Register (Address X'00C6)

The T2CNTRL register contains the following bits:

T2ENB Timer T2 Interrupt Enable for T2B Input capture edge

T2PNDB Timer T2 Interrupt Pending Flag for T2B capture edge

T2ENA Timer T2 Interrupt Enable for Timer Underflow or T2A Input capture edge

T2PNDA Timer T2 Interrupt Pending Flag (Autoreload RA in mode 1, T2 Underflow in mode 2, T2A capture edge in mode 3)

T2C0 Timer T2 Start/Stop control in timer modes 1 and 2 Timer T2 Underflow Interrupt Pending Flag in timer mode 3 T2C1 Timer T2 mode control bit
T2C2 Timer T2 mode control bit
T2C3 Timer T2 mode control bit

T2C3	T2C2	T2C1	T2C0	T2PNDA	T2ENA	T2PNDB	T2ENB	
Bit 7							Bit 0	

T3CNTRL Register (Address X'00B6)

The T3CNTRL register contains the following bits:

T3ENB Timer T3 Interrupt Enable for T3B

T3PNDB Timer T3 Interrupt Pending Flag for T3B pin

(T3B capture edge)

T3ENA Timer T3 Interrupt Enable for Timer Underflow

or T3A pin

T3PNDA Timer T3 Interrupt Pending Flag (Autoload RA in mode 1, T3 Underflow in mode 2, T3a cap-

ture edge in mode 3)

T3C0 Timer T3 Start/Stop control in timer modes 1

and 2

Timer T3 Underflow Interrupt Pending Flag in

Rit 0

timer mode 3

T3C1 Timer T3 mode control bit
T3C2 Timer T3 mode control bit
T3C3 Timer T3 mode control bit

T3C3 T3C2 T3C1 T3C0 T3PNDA T3ENA T3PNDB T3ENB

Timers

Bit 7

The device contains a very versatile set of timers (T0, T1, T2, T3). All timers and associated autoreload/capture registers power up containing random data.

TIMER TO (IDLE TIMER)

The devices support applications that require maintaining real time and low power with the IDLE mode. This IDLE mode support is furnished by the IDLE timer T0, which is a 16-bit timer. The Timer T0 runs continuously at the fixed rate of the instruction cycle clock, $t_{\rm c}$. The user cannot read or write to the IDLE Timer T0, which is a count down timer. The Timer T0 supports the following functions:

Exit out of the Idle Mode (See Idle Mode description) WATCHDOG logic (See WATCHDOG description) Start up delay out of the HALT mode

The IDLE Timer T0 can generate an interrupt when the thirteenth bit toggles. This toggle is latched into the T0PND pending flag, and will occur every 4 ms at the maximum clock frequency ($t_{\rm C}=1~\mu{\rm s}$). A control flag T0EN allows the interrupt from the thirteenth bit of Timer T0 to be enabled or disabled. Setting T0EN will enable the interrupt, while resetting it will disable the interrupt.

Timers (Continued)

TIMER T1, TIMER T2 AND TIMER T3

The devices have a set of three powerful timer/counter blocks, T1, T2 and T3. The associated features and functioning of a timer block are described by referring to the timer block Tx. Since the three timer blocks, T1, T2 and T3 are identical, all comments are equally applicable to any of the three timer blocks.

Each timer block consists of a 16-bit timer, Tx, and two supporting 16-bit autoreload/capture registers, RxA and RxB. Each timer block has two pins associated with it, TxA and TxB. The pin TxA supports I/O required by the timer block, while the pin TxB is an input to the timer block. The powerful and flexible timer block allows the device to easily perform all timer functions with minimal software overhead. The timer block has three operating modes: Processor Independent PWM mode, External Event Counter mode, and Input Capture mode.

The control bits TxC3, TxC2, and TxC1 allow selection of the different modes of operation.

Mode 1. Processor Independent PWM Mode

As the name suggests, this mode allows the device to generate a PWM signal with very minimal user intervention. The user only has to define the parameters of the PWM signal (ON time and OFF time). Once begun, the timer block will continuously generate the PWM signal completely independent of the microcontroller. The user software services the timer block only when the PWM parameters require updating.

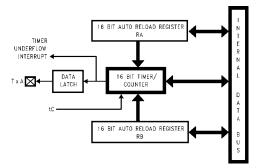
In this mode the timer Tx counts down at a fixed rate of t_c . Upon every underflow the timer is alternately reloaded with the contents of supporting registers, RxA and RxB. The very first underflow of the timer causes the timer to reload from the register RxA. Subsequent underflows cause the timer to be reloaded from the registers alternately beginning with the register RxB.

The Tx Timer control bits, TxC3, TxC2 and TxC1 set up the timer for PWM mode operation.

Figure ϑ shows a block diagram of the timer in PWM mode. The underflows can be programmed to toggle the TxA output pin. The underflows can also be programmed to generate interrupts.

Underflows from the timer are alternately latched into two pending flags, TxPNDA and TxPNDB. The user must reset these pending flags under software control. Two control enable flags, TxENA and TxENB, allow the interrupts from the timer underflow to be enabled or disabled. Setting the timer enable flag TxENA will cause an interrupt when a timer underflow causes the RxA register to be reloaded into the timer. Setting the timer enable flag TxENB will cause an interrupt when a timer underflow causes the RxB register to be reloaded into the timer. Resetting the timer enable flags will disable the associated interrupts.

Either or both of the timer underflow interrupts may be enabled. This gives the user the flexibility of interrupting once per PWM period on either the rising or falling edge of the PWM output. Alternatively, the user may choose to interrupt on both edges of the PWM output.



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FIGURE 8. Timer in PWM Mode

Mode 2. External Event Counter Mode

This mode is quite similar to the processor independent PWM mode described above. The main difference is that the timer, Tx, is clocked by the input signal from the TxA pin. The Tx timer control bits, TxC3, TxC2 and TxC1 allow the timer to be clocked either on a positive or negative edge from the TxA pin. Underflows from the timer are latched into the TxPNDA pending flag. Setting the TxENA control flag will cause an interrupt when the timer underflows.

In this mode the input pin TxB can be used as an independent positive edge sensitive interrupt input if the TxENB control flag is set. The occurrence of a positive edge on the TxB input pin is latched into the TxPNDB flag.

Figure 9 shows a block diagram of the timer in External Event Counter mode.

Note: The PWM output is not available in this mode since the TxA pin is being used as the counter input clock.

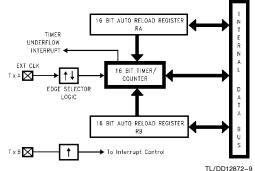


FIGURE 9. Timer in External Event Counter Mode

Mode 3. Input Capture Mode

The device can precisely measure external frequencies or time external events by placing the timer block, Tx, in the input capture mode.

In this mode, the timer Tx is constantly running at the fixed t_{c} rate. The two registers, RxA and RxB, act as capture registers. Each register acts in conjunction with a pin. The register RxA acts in conjunction with the TxA pin and the register RxB acts in conjunction with the TxB pin.

Timers (Continued)

The timer value gets copied over into the register when a trigger event occurs on its corresponding pin. Control bits, TxC3, TxC2 and TxC1, allow the trigger events to be specified either as a positive or a negative edge. The trigger condition for each input pin can be specified independently.

The trigger conditions can also be programmed to generate interrupts. The occurrence of the specified trigger condition on the TxA and TxB pins will be respectively latched into the pending flags, TxPNDA and TxPNDB. The control flag TxENA allows the interrupt on TxA to be either enabled or disabled. Setting the TxENA flag enables interrupts to be generated when the selected trigger condition occurs on the TxA pin. Similarly, the flag TxENB controls the interrupts from the TxB pin.

Underflows from the timer can also be programmed to generate interrupts. Underflows are latched into the timer TxC0 pending flag (the TxC0 control bit serves as the timer underflow interrupt pending flag in the Input Capture mode). Consequently, the TxC0 control bit should be reset when entering the Input Capture mode. The timer underflow interrupt is enabled with the TxENA control flag. When a TxA interrupt occurs in the Input Capture mode, the user must check both the TxPNDA and TxC0 pending flags in order to determine whether a TxA input capture or a timer underflow (or both) caused the interrupt.

Figure 10 shows a block diagram of the timer in Input Capture mode.

TIMER CONTROL FLAGS

The timers T1, T2 and T3 have indentical control structures. The control bits and their functions are summarized below.

TxC0 Timer Start/Stop control in Modes 1 and 2 (Processor Independent PWM and External Event Counter), where 1 = Start, 0 = Stop Timer Underflow Interrupt Pending Flag in Mode 3 (Input Capture)

TxC3 Timer mode control TxC2 Timer mode control TxC1 Timer mode control

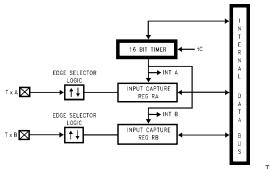


FIGURE 10. Timer in Input Capture Mode

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Timers (Continued)

The timer mode control bits (TxC3, TxC2 and TxC1) are detailed below:

TxC3	TxC2	TxC1	Timer Mode	Interrupt A Source	Interrupt B Source	Timer Counts On
0	0	0	MODE 2 (External Event Counter)	`		TxA Pos. Edge
0	0	1	MODE 2 (External Event Counter)	Timer Underflow	Pos. TxB Edge	TxA Neg. Edge
1	0	1	MODE 1 (PWM) TxA Toggle	Autoreload RA	Autoreload RB	t _c
1	0	0	MODE 1 (PWM) No TxA Toggle	Autoreload RA	Autoreload RB	t _c
0	1	0	MODE 3 (Capture) Captures: TxA Pos. Edge TxB Pos. Edge	Pos. TxA Edge or Timer Underflow	Pos. TxB Edge	t _c
1	1	0	MODE 3 (Capture) Captures: TxA Pos. Edge TxB Neg. Edge	MODE 3 (Capture) Captures: Edge or TxA Pos. Edge Timer Pos. TxA Neg. TxB Edge Edge		t _c
0	1	1	MODE 3 (Capture) Captures: TxA Neg. Edge TxB Pos. Edge	B (Capture) Neg. TxA Pos. TxB s: Edge or g. Edge Timer		t _c
1	1	1	MODE 3 (Capture) Captures: TxA Neg. Edge TxB Neg. Edge	Neg. TxA Edge or Timer Underflow	Neg. TxB Edge	t _c

Power Save Modes

The devices offer the user two power save modes of operation: HALT and IDLE. In the HALT mode, all microcontroller activities are stopped. In the IDLE mode, the on-board oscillator circuitry the WATCHDOG logic, the Clock Monitor and timer T0 are active but all other microcontroller activities are stopped. In either mode, all on-board RAM, registers, I/O states, and timers (with the exception of T0) are unaltered.

HALT MODE

The devices can be placed in the HALT mode by writing a "1" to the HALT flag (G7 data bit). All microcontroller activities, including the clock and timers, are stopped. The WATCHDOG logic on the device is disabled during the HALT mode. However, the clock monitor circuitry if enabled remains active and will cause the WATCHDOG output pin (WDOUT) to go low. If the HALT mode is used and the user does not want to activate the WDOUT pin, the Clock Monitor should be disabled after the device comes out of reset (resetting the Clock Monitor control bit with the first write to the WDSVR register). In the HALT mode, the power requirements of the device are minimal and the applied voltage (VCC) may be decreased to V_{Γ} ($\mathrm{V}_{\Gamma}=2.0\mathrm{V}$) without altering the state of the machine.

The devices support three different ways of exiting the HALT mode. The first method of exiting the HALT mode is with the Multi-Input Wake Up feature on the L port. The second method is with a low to high transition on the CKO (G7) pin. This method precludes the use of the crystal clock

configuration (since CKO becomes a dedicated output), and so may be used with an RC clock configuration. The third method of exiting the HALT mode is by pulling the $\overline{\text{RESET}}$ pin low.

Since a crystal or ceramic resonator may be selected as the oscillator, the Wake Up signal is not allowed to start the chip running immediately since crystal oscillators and ceramic resonators have a delayed start up time to reach full amplitude and frequency stability. The IDLE timer is used to generate a fixed delay to ensure that the oscillator has indeed stabilized before allowing instruction execution. In this case, upon detecting a valid Wake Up signal, only the oscillator circuitry is enabled. The IDLE timer is loaded with a value of 256 and is clocked with the t_c instruction cycle clock. The t_c clock is derived by dividing the oscillator clock down by a factor of 10. The Schmitt trigger following the CKI inverter on the chip ensures that the IDLE timer is clocked only when the oscillator has a sufficiently large amplitude to meet the Schmitt trigger specifications. This Schmitt trigger is not part of the oscillator closed loop. The startup timeout from the IDLE timer enables the clock signals to be routed to the rest of the chip.

If an RC clock option is being used, the fixed delay is introduced optionally. A control bit, CLKDLY, mapped as configuration bit G7, controls whether the delay is to be introduced or not. The delay is included if CLKDLY is set, and excluded if CLKDLY is reset. The CLKDLY bit is cleared on reset.

Power Save Modes (Continued)

The WATCHDOG detector circuit is inhibited during the HALT mode. However, the clock monitor circuit if enabled remains active during HALT mode in order to ensure a clock monitor error if the device inadvertently enters the HALT mode as a result of a runaway program or power glitch.

IDLE MODE

The device is placed in the IDLE mode by writing a "1" to the IDLE flag (G6 data bit). In this mode, all activities, except the associated on-board oscillator circuitry, the WATCHDOG logic, the clock monitor and the IDLE Timer T0, are stopped. The power supply requirements of the micro-controller in this mode of operation are typically around 30% of normal power requirement of the microcontroller.

As with the HALT mode, the device can be returned to normal operation with a reset, or with a Multi-Input Wake Up from the L Port. Alternately, the microcontroller resumes normal operation from the IDLE mode when the thirteenth bit (representing 4.096 ms at internal clock frequency of 1 MHz, $t_{\rm C}=1~\mu{\rm s})$ of the IDLE Timer toggles.

This toggle condition of the thirteenth bit of the IDLE Timer T0 is latched into the T0PND pending flag.

The user has the option of being interrupted with a transition on the thirteenth bit of the IDLE Timer T0. The interrupt can be enabled or disabled via the T0EN control bit. Setting the T0EN flag enables the interrupt and vice versa.

The user can enter the IDLE mode with the Timer T0 interrupt enabled. In this case, when the T0PND bit gets set, the device will first execute the Timer T0 interrupt service routine and then return to the instruction following the "Enter Idle Mode" instruction.

Alternatively, the user can enter the IDLE mode with the IDLE Timer T0 interrupt disabled. In this case, the device will resume normal operation with the instruction immediately following the "Enter IDLE Mode" instruction.

Note: It is necessary to program two NOP instructions following both the set HALT mode and set IDLE mode instructions. These NOP instructions are necessary to allow clock resynchronization following the HALT or IDLE modes.

Due to the on-board 8k EPROM with port recreation logic, the HALT/IDLE current is much higher compared to the equivalent masked port.

Multi-Input Wake Up

The Multi-Input Wake Up feature is ued to return (Wake Up) the device from either the HALT or IDLE modes. Alternately Multi-Input Wake Up/Interrupt feature may also be used to generate up to 8 edge selectable external interrupts.

Figure 11 shows the Multi-Input Wake Up logic. The Multi-Input Wake Up feature utilizes the L Port. The user selects which particular L port bit (or combination of L Port bits) will cause the device to exit the HALT or IDLE modes. The selection is done through the Reg: WKEN. The Reg: WKEN

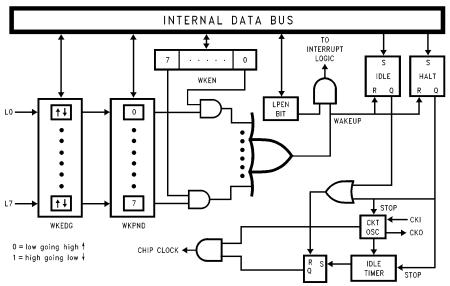


FIGURE 11. Multi-Input Wake Up Logic

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Multi-Input Wake Up (Continued)

is an 8-bit read/write register, which contains a control bit for every L port bit. Setting a particular WKEN bit enables a Wake Up from the associated L port pin.

The user can select whether the trigger condition on the selected L Port pin is going to be either a positive edge (low to high transition) or a negative edge (high to low transition). This selection is made via the Reg: WKEDG, which is an 8-bit control register with a bit assigned to each L Port pin. Setting the control bit will select the trigger condition to be a negative edge on that particular L Port pin. Resetting the bit selects the trigger condition to be a positive edge. Changing an edge select entails several steps in order to avoid a pseudo Wake Up condition as a result of the edge change. First, the associated WKEN bit should be reset, followed by the edge select change in WKEDG. Next, the associated WKPND bit should be cleared, followed by the associated WKEN bit being re-enabled.

An example may serve to clarify this procedure. Suppose we wish to change the edge select from positive (low going high) to negative (high going low) for L Port bit 5, where bit 5 has previously been enabled for an input interrupt. The program would be as follows:

RBIT 5, WKEN
SBIT 5, WKEDG
RBIT 5, WKPND
SBIT 5. WKEN

If the L port bits have been used as outputs and then changed to inputs with Multi-Input Wake Up/Interrupt, a safety procedure should also be followed to avoid inherited pseudo wakeup conditions. After the selected L port bits have been changed from output to input but before the associated WKEN bits are enabled, the associated edge select bits in WKEDG should be set or reset for the desired edge selects, followed by the associated WKPND bits being cleared

This same procedure should be used following reset, since the L port inputs are left floating as a result of reset.

The occurrence of the selected trigger condition for Multi-Input Wake Up is latched into a pending register called WKPND. The respective bits of the WKPND register will be set on the occurrence of the selected trigger edge on the corresponding Port L pin. The user has the responsibility of clearing these pending flags. Since WKPND is a pending register for the occurrence of selected Wake Up conditions, the device will not enter the HALT mode if any Wake Up bit is both enabled and pending. Consequently, the user has the responsibility of clearing the pending flags before attempting to enter the HALT mode.

WKEN, WKPND and WKEDG are all read/write registers, and are cleared at reset.

PORT L INTERRUPTS

Port L provides the user with an additional eight fully selectable, edge sensitive interrupts which are all vectored into the same service subroutine.

The interrupt from Port L shares logic with the wake up circuitry. The register WKEN allows interrupts from Port L to be individually enabled or disabled. The register WKEDG specifies the trigger condition to be either a positive or a negative edge. Finally, the register WKPND latches in the pending trigger conditions.

The GIE (Global Interrupt Enable) bit enables the interrupt function.

A control flag, LPEN, functions as a global interrupt enable for Port L interrupts. Setting the LPEN flag will enable interrupts and vice versa. A separate global pending flag is not needed since the register WKPND is adequate.

Since Port L is also used for waking the device out of the HALT or IDLE modes, the user can elect to exit the HALT or IDLE modes either with or without the interrupt enabled. If he elects to disable the interrupt, then the device will restart execution from the instruction immediately following the instruction that placed the microcontroller in the HALT or IDLE modes. In the other case, the device will first execute the interrupt service routine and then revert to normal operation.

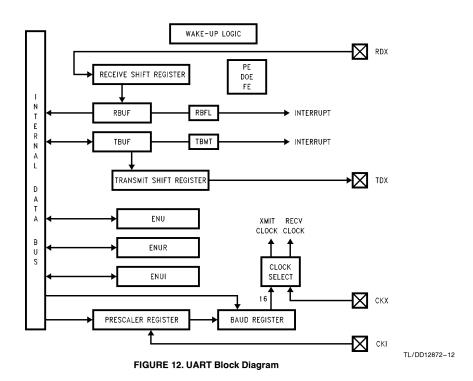
The Wake Up signal will not start the chip running immediately since crystal oscillators or ceramic resonators have a finite start up time. The IDLE Timer (T0) generates a fixed delay to ensure that the oscillator has indeed stabilized before allowing the device to execute instructions. In this case, upon detecting a valid Wake Up signal, only the oscillator circuitry and the IDLE Timer T0 are enabled. The IDLE Timer is loaded with a value of 256 and is clocked from the tc instruction cycle clock. The t_{c} clock is derived by dividing down the oscillator clock by a factor of 10. A Schmitt trigger following the CKI on-chip inverter ensures that the IDLE timer is clocked only when the oscillator has a sufficiently large amplitude to meet the Schmitt trigger specifications. This Schmitt trigger is not part of the oscillator closed loop. The startup timeout from the IDLE timer enables the clock signals to be routed to the rest of the chip.

If the RC clock option is used, the fixed delay is under software control. A control flag, CLKDLY, in the G7 configuration bit allows the clock start up delay to be optionally inserted. Setting CLKDLY flag high will cause clock start up delay to be inserted and resetting it will exclude the clock start up delay. The CLKDLY flag is cleared during reset, so the clock start up delay is not present following reset with the RC clock options.

UART

The device contains a full-duplex software programmable UART. The UART (Figure 12) consists of a transmit shift register, a receiver shift register and seven addressable registers, as follows: a transmit buffer register (TBUF), a receiver buffer register (RBUF), a UART control and status register (ENU), a UART receive control and status register (ENUR), a UART interrupt and clock source register (ENUI), a prescaler select register (PSR) and baud (BAUD) register. The ENU register contains flags for transmit and receive functions; this register also determines the length of the data frame (7, 8 or 9 bits), the value of the ninth bit in transmission, and parity selection bits. The ENUR register flags framing, data overrun and parity errors while the UART is receiving.

Other functions of the ENUR register include saving the ninth bit received in the data frame, enabling or disabling the UART's attention mode of operation and providing additional receiver/transmitter status information via RCVG and XMTG bits. The determination of an internal or external clock source is done by the ENUI register, as well as selecting the number of stop bits and enabling or disabling transmit and receive interrupts. A control flag in this register can also select the UART mode of operation: asynchronous or synchronous.



UART (Continued)

UART CONTROL AND STATUS REGISTERS

The operation of the UART is programmed through three registers: ENU, ENUR and ENUI. The function of the individual bits in these registers is as follows:

ENU-UART Control and Status Register (Address at 0BA)

PEN	PSEL1	XBIT9/	CHL1	CHL0	ERR	RBFL	ТВМТ
		PSEL0					
0RW	0RW	0RW	0RW	0RW	0R	0R	1R

Bit 7 Bit 0

ENUR-UART Receive Control and Status Register (Address at 0BB)

DOE	FE	PE	SPARE	RBIT9	ATTN	XMTG	RCVG
0RD	0RD	0RD	0RW*	0R	0RW	0R	0R

Bit7 Bit0

ENUI-UART Interrupt and Clock Source Register (Address at 0BC)

STP2	STP78	ETDX	SSEL	XRCLK	XTCLK	ERI	ETI
0RW	0RW	0RW	0RW	0RW	0RW	0RW	0RW

Bit7
*Bit is not used.

Bit is cleared on reset

- o Bit is cleared on reset.
- 1 Bit is set to one on reset.
- R Bit is read-only; it cannot be written by software.

RW Bit is read/write.

D Bit is cleared on read; when read by software as a one, it is cleared automatically. Writing to the bit does not affect its state.

DESCRIPTION OF UART REGISTER BITS

ENU—UART CONTROL AND STATUS REGISTER

TBMT: This bit is set when the UART transfers a byte of data from the TBUF register into the TSFT register for transmission. It is automatically reset when software writes into the TBUF register.

RBFL: This bit is set when the UART has received a complete character and has copied it into the RBUF register. It is automatically reset when software reads the character from RBUF.

ERR: This bit is a global UART error flag which gets set if any or a combination of the errors (DOE, FE, PE) occur.

CHL1, CHL0: These bits select the character frame format. Parity is not included and is generated/verified by hardware.

CHL1 = 0, CHL0 = 0 The frame contains eight data bits. CHL1 = 0, CHL0 = 1 The frame contains seven data

bits.

back to receiver input. Nine bit framing format is used.

ams the ninth bit for transmission

XBIT9/PSEL0: Programs the ninth bit for transmission when the UART is operating with nine data bits per frame. For seven or eight data bits per frame, this bit in conjunction with PSEL1 selects parity.

PSEL1, PSEL0: Parity select bits.

PSEL1 = 0, PSEL0 = 0 Odd Parity (if Parity enabled)
PSEL1 = 0, PSEL0 = 1 Odd Parity (if Parity enabled)

PSEL1 = 1, PSEL0 = 0 Mark(1) (if Parity enabled) PSEL1 = 1, PSEL1 = 1 Space(0) (if Parity enabled)

PEN: This bit enables/disables Parity (7- and 8-bit modes only).

PEN = 0 Parity disabled.

PEN = 1 Parity enabled.

ENUR—UART RECEIVE CONTROL AND STATUS REGISTER

RCVG: This bit is set high whenever a framing error occurs and goes low when RDX goes high.

XMTG: This bit is set to indicate that the UART is transmitting. It gets reset at the end of the last frame (end of last Stop bit).

ATTN: ATTENTION Mode is enabled while this bit is set. This bit is cleared automatically on receiving a character with data bit nine set.

RBIT9: Contains the ninth data bit received when the UART is operating with nine data bits per frame.

SPARE: Reserved for future use.

PE: Flags a Parity Error.

Bit0

PE=0 Indicates no Parity Error has been detected since the last time the ENUR register was read.

PE = 1 Indicates the occurrence of a Parity Error.

FE: Flags a Framing Error.

FE=0 Indicates no Framing Error has been detected since the last time the ENUR register was read.

FE = 1 Indicates the occurrence of a Framing Error.

DOE: Flags a Data Overrun Error.

DOE = 0 Indicates no Data Overrun Error has been detected since the last time the ENUR register was read.

DOE = 1 Indicates the occurrence of a Data Overrun Error.

ENUI—UART INTERRUPT AND CLOCK SOURCE REGISTER

ETI: This bit enables/disables interrupt from the transmitter section.

 $\mathsf{ETI} = \mathsf{0}$ Interrupt from the transmitter is disabled.

ETI = 1 Interrupt from the transmitter is enabled.

ERI: This bit enables/disables interrupt from the receiver section.

ERI = 0 Interrupt from the receiver is disabled.

ERI = 1 Interrupt from the receiver is enabled.

XTCLK: This bit selects the clock source for the transmitter section.

 $\mathsf{XTCLK} = \mathsf{0}$ The clock source is selected through the PSR and BAUD registers.

XTCLK = 1 Signal on CKX (L1) pin is used as the clock.

XRCLK: This bit selects the clock source for the receiver section.

 $\mathsf{XRCLK} = 0$ The clock source is selected through the PSR and BAUD registers.

XRCLK = 1 Signal on CKX (L1) pin is used as the clock.

SSEL: UART mode select.

SSEL = 0 Asynchronous Mode.

SSEL = 1 Synchronous Mode.

UART (Continued)

ETDX: TDX (UART Transmit Pin) is the alternate function assigned to Port L pin L2; it is selected by setting ETDX bit. To simulate line break generation, software should reset ETDX bit and output logic zero to TDX pin through Port L data and configuration registers.

STP78: This bit is set to program the last Stop bit to be 7/8th of a bit in length.

STP2: This bit programs the number of Stop bits to be transmitted.

STP2 = 0 One Stop bit transmitted. STP2 = 1 Two Stop bits transmitted.

Associated I/O Pins

Data is transmitted on the TDX pin and received on the RDX pin. TDX is the alternate function assigned to Port L pin L2; it is selected by setting ETDX (in the ENUI register) to one. RDX is an inherent function of Port L pin L3, requiring no setup.

The baud rate clock for the UART can be generated onchip, or can be taken from an external source. Port L pin L1 (CKX) is the external clock I/O pin. The CKX pin can be either an input or an output, as determined by Port L Configuration and Data registers (Bit 1). As an input, it accepts a clock signal which may be selected to drive the transmitter and/or receiver. As an output, it presents the internal Baud Rate Generator output.

UART Operation

The UART has two modes of operation: asynchronous mode and synchronous mode.

ASYNCHRONOUS MODE

This mode is selected by resetting the SSEL (in the ENUI register) bit to zero. The input frequency to the UART is 16 times the baud rate.

The TSFT and TBUF registers double-buffer data for transmission. While TSFT is shifting out the current character on the TDX pin, the TBUF register may be loaded by software with the next byte to be transmitted. When TSFT finishes transmitting the current character the contents of TBUF are transferred to the TSFT register and the Transmit Buffer Empty Flag (TBMT in the ENU register) is set. The TBMT flag is automatically reset by the UART when software loads a new character into the TBUF register. There is also the XMTG bit which is set to indicate that the UART is transmitting. This bit gets reset at the end of the last frame (end of last Stop bit). TBUF is a read/write register.

The RSFT and RBUF registers double-buffer data being received. The UART receiver continually monitors the signal on the RDX pin for a low level to detect the beginning of a Start bit. Upon sensing this low level, it waits for half a bit time and samples again. If the RDX pin is still low, the receiver considers this to be a valid Start bit, and the remaining bits in the character frame are each sampled a single time, at the mid-bit position. Serial data input on the RDX pin is shifted into the RSFT register. Upon receiving the complete character, the contents of the RSFT register are copied into the RBUF register and the Received Buffer Full Flag (RBFL) is set. RBFL is automatically reset when software reads the character from the RBUF register. RBUF is a read only register. There is also the RCVG bit which is set high

when a framing error occurs and goes low once RDX goes high. TBMT, XMTG, RBFL and RCVG are read only bits.

SYNCHRONOUS MODE

In this mode data is transferred synchronously with the clock. Data is transmitted on the rising edge and received on the falling edge of the synchronous clock.

This mode is selected by setting SSEL bit in the ENUI register. The input frequency to the UART is the same as the baud rate.

When an external clock input is selected at the CKX pin, data transmit and receive are performed synchronously with this clock through TDX/RDX pins.

If data transmit and receive are selected with the CKX pin as clock output, the device generates the synchronous clock output at the CKX pin. The internal baud rate generator is used to produce the synchronous clock. Data transmit and receive are performed synchronously with this clock.

FRAMING FORMATS

The UART supports several serial framing formats (Figure 13). The format is selected using control bits in the ENU, ENUR and ENUI registers.

The first format (1, 1a, 1b, 1c) for data transmission (CHL0 = 1, CHL1 = 0) consists of Start bit, seven Data bits (excluding parity) and 7/8, one or two Stop bits. In applications using parity, the parity bit is generated and verified by hardware

The second format (CHL0 = 0, CHL1 = 0) consists of one Start bit, eight Data bits (excluding parity) and 7/8, one or two Stop bits. Parity bit is generated and verified by hardware

The third format for transmission (CHL0 = 0, CHL1 = 1) consists of one Start bit, nine Data bits and 7/8, one or two Stop bits. This format also supports the UART "ATTENTION" feature. When operating in this format, all eight bits of TBUF and RBUF are used for data. The ninth data bit is transmitted and received using two bits in the ENU and ENUR registers, called XBIT9 and RBIT9. RBIT9 is a read only bit. Parity is not generated or verified in this mode.

For any of the above framing formats, the last Stop bit can be programmed to be 7/8th of a bit in length. If two Stop bits are selected and the 7/8th bit is set (selected), the second Stop bit will be 7/8th of a bit in length.

The parity is enabled/disabled by PEN bit located in the ENU register. Parity is selected for 7- and 8-bit modes only. If parity is enabled (PEN = 1), the parity selection is then performed by PSEL0 and PSEL1 bits located in the ENU register.

Note that the XBIT9/PSEL0 bit located in the ENU register serves two mutually exclusive functions. This bit programs the ninth bit for transmission when the UART is operating with nine data bits per frame. There is no parity selection in this framing format. For other framing formats XBIT9 is not needed and the bit is PSEL0 used in conjunction with PSEL1 to select parity.

The frame formats for the receiver differ from the transmitter in the number of Stop bits required. The receiver only requires one Stop bit in a frame, regardless of the setting of the Stop bit selection bits in the control register. Note that an implicit assumption is made for full duplex UART operation that the framing formats are the same for the transmitter and receiver.

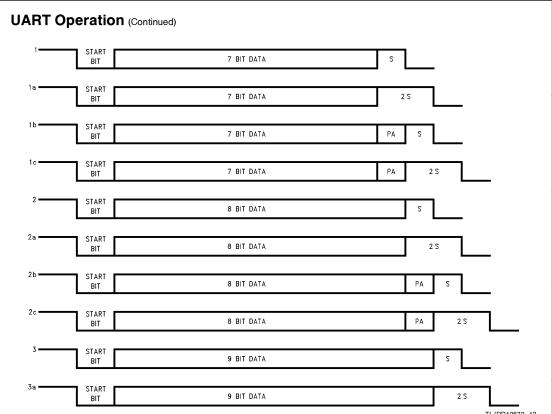


FIGURE 13. Framing Formats

TL/DD12872-13

UART INTERRUPTS

The UART is capable of generating interrupts. Interrupts are generated on Receive Buffer Full and Transmit Buffer Empty. Both interrupts have individual interrupt vectors. Two bytes of program memory space are reserved for each interrupt vector. The two vectors are located at addresses 0xEC to 0xEF Hex in the program memory space. The interrupts can be individually enabled or disabled using Enable Transmit Interrupt (ETI) and Enable Receive Interrupt (ERI) bits in the ENUI register.

The interrupt from the transmitter is set pending, and remains pending, as long as both the TBMT and ETI bits are set. To remove this interrupt, software must either clear the ETI bit or write to the TBUF register (thus clearing the TBMT bit).

The interrupt from the receiver is set pending, and remains pending, as long as both the RBFL and ERI bits are set. To remove this interrupt, software must either clear the ERI bit or read from the RBUF register (thus clearing the RBFL bit).

Baud Clock Generation

The clock inputs to the transmitter and receiver sections of the UART can be individually selected to come either from an external source at the CKX pin (port L, pin L1) or from a

source selected in the PSR and BAUD registers. Internally, the basic baud clock is created from the oscillator frequency through a two-stage divider chain consisting of a 1-16 (increments of 0.5) prescaler and an 11-bit binary counter. (Figure 14) The divide factors are specified through two read/write registers shown in Figure 15. Note that the 11-bit Baud Rate Divisor spills over into the Prescaler Select Register (PSR). PSR is cleared upon reset.

As shown in Table III, a Prescaler Factor of 0 corresponds to NO CLOCK. NO CLOCK condition is the UART power down mode where the UART clock is turned off for power saving purpose. The user must also turn the UART clock off when a different baud rate is chosen.

The correspondences between the 5-bit Prescaler Select and Prescaler factors are shown in Table III. There are many ways to calculate the two divisor factors, but one particularly effective method would be to achieve a 1.8432 MHz frequency coming out of the first stage. The 1.8432 MHz prescaler output is then used to drive the software programmable baud rate counter to create a x16 clock for the following baud rates: 110, 134.5, 150, 300, 600, 1200, 1800, 2400, 3600, 4800, 7200, 9600, 19200 and 38400 (Table IV). Other baud rates may be created by using appropriate divisors. The x16 clock is then divided by 16 to provide the rate for the serial shift registers of the transmitter and receiver.

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Baud Clock Generation (Continued)

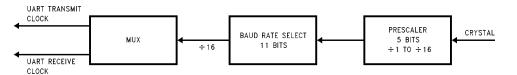


FIGURE 14. UART BAUD Clock Generation

TL/DD12872-14

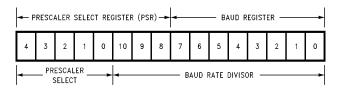


FIGURE 15. UART BAUD Clock Divisor Registers

TABLE III. Prescaler Factors

Prescaler Select	Prescaler Factor	Prescaler Select	Prescaler Factor
00000	NO CLOCK	10000	8.5
00001	1	10001	9
00010	1.5	10010	9.5
00011	2	10011	10
00100	2.5	10100	10.5
00101	3	10101	11
00110	3.5	10110	11.5
00111	4	10111	12
01000	4.5	11000	12.5
01001	5	11001	13
01010	5.5	11010	13.5
01011	6	11011	14
01100	6.5	11100	14.5
01101	7	11101	15
01110	7.5	11110	15.5
01111	8	11111	16

TABLE IV. Baud Rate Divisors (1.8432 MHz Prescaler Output)

TL/DD12872-15

-	
Baud Rate	Baud Rate Divisor – 1 (N-1)
110 (110.03)	1046
134.5 (134.58)	855
150	767
300	383
600	191
1200	95
1800	63
2400	47
3600	31
4800	23
7200	15
9600	11
19200	5
38400	2

Note: The entries in Table IV assume a prescaler output of 1.8432 MHz. In the asynchronous mode the baud rate could be as high as 625k.

As an example, considering the Asynchronous Mode and a CKI clock of 4.608 MHz, the prescaler factor selected is:

The 2.5 entry is available in Table III. The 1.8432 MHz prescaler output is then used with proper Baud Rate Divisor (Table II) to obtain different baud rates. For a baud rate of 19200 e.g., the entry in Table IV is 5.

N-1=5 (N -1 is the value from Table IV)

N = 6 (N is the Baud Rate Divisor)

Baud Rate = $1.8432 \, \text{MHz}/(16 \times 6) = 19200$

The divide by 16 is performed because in the asynchronous mode, the input frequency to the UART is 16 times the baud rate. The equation to calculate baud rates is given below.

The actual Baud Rate may be found from:

 $BR = Fc/(16 \times N \times P)$

Baud Clock Generation (Continued)

Where:

BR is the Baud Rate

Fc is the CKI frequency

N is the Baud Rate Divisor (Table IV).

P is the Prescaler Divide Factor selected by the value in the Prescaler Select Register (Table III)

Note: In the Synchronous Mode, the divisor 16 is replaced by two. **Example:**

Asynchronous Mode:

Crystal Frequency = 5 MHz

Desired baud rate = 9600

Using the above equation N \times P can be calculated first.

$$N \times P = (5 \times 10^6)/(16 \times 9600) = 32.552$$

Now 32.552 is divided by each Prescaler Factor (Table III) to obtain a value closest to an integer. This factor happens to be 6.5 (P=6.5).

$$N = 32.552/6.5 = 5.008 (N = 5)$$

The programmed value (from Table IV) should be 4 (N - 1). Using the above values calculated for N and P:

BR =
$$(5 \times 10^6)/(16 \times 5 \times 6.5) = 9615.384$$

% error = $(9615.385 - 9600)/9600 = 0.16$

Effect of HALT/IDLE

The UART logic is reinitialized when either the HALT or IDLE modes are entered. This reinitialization sets the TBMT flag and resets all read only bits in the UART control and status registers. Read/Write bits remain unchanged. The Transmit Buffer (TBUF) is not affected, but the Transmit Shift register (TSFT) bits are set to one. The receiver registers RBUF and RSFT are not affected.

The device will exit from the HALT/IDLE modes when the Start bit of a character is detected at the RDX (L3) pin. This feature is obtained by using the Multi-Input Wake Up scheme provided on the device.

Before entering the HALT or IDLE modes the user program must select the Wake Up source to be on the RDX pin. This selection is done by setting bit 3 of WKEN (Wake Up Enable) register. The Wake Up trigger condition is then selected to be high to low transition. This is done via the WKEDG register (Bit 3 is one.)

If the device is halted and crystal oscillator is used, the Wake Up signal will not start the chip running immediately because of the finite start up time requirement of the crystal oscillator. The idle timer (To) generates a fixed (256 $\rm t_c)$ delay to ensure that the oscillator has indeed stabilized before allowing the device to execute code. The user has to consider this delay when data transfer is expected immediately after exiting the HALT mode.

Diagnostic

Bits CHARL0 and CHARL1 in the ENU register provide a loopback feature for diagnostic testing of the UART. When these bits are set to one, the following occur: The receiver input pin (RDX) is internally connected to the transmitter output pin (TDX); the output of the Transmitter Shift Register is "looped back" into the Receive Shift Register input. In this mode, data that is transmitted is immediately received. This feature allows the processor to verify the transmit and receive data paths of the UART.

Note that the framing format for this mode is the nine bit format; one Start bit, nine data bits, and 7/8, one or two Stop bits. Parity is not generated or verified in this mode.

Attention Mode

The UART Receiver section supports an alternate mode of operation, referred to as ATTENTION Mode. This mode of operation is selected by the ATTN bit in the ENUR register. The data format for transmission must also be selected as having nine Data bits and either 7/8, one or two Stop bits.

The ATTENTION mode of operation is intended for use in networking the device with other processors. Typically in such environments the messages consists of device addresses, indicating which of several destinations should receive them, and the actual data. This Mode supports a scheme in which addresses are flagged by having the ninth bit of the data field set to a 1. If the ninth bit is reset to a zero the byte is a Data byte.

While in ATTENTION mode, the UART monitors the communication flow, but ignores all characters until an address character is received. Upon receiving an address character, the UART signals that the character is ready by setting the RBFL flag, which in turn interrupts the processor if UART Receiver interrupts are enabled. The ATTN bit is also cleared automatically at this point, so that data characters as well as address characters are recognized. Software examines the contents of the RBUF and responds by deciding either to accept the subsequent data stream (by leaving the ATTN bit reset) or to wait until the next address character is seen (by setting the ATTN bit again).

Operation of the UART Transmitter is not affected by selection of this Mode. The value of the ninth bit to be transmitted is programmed by setting XBIT9 appropriately. The value of the ninth bit received is obtained by reading RBIT9. Since this bit is located in ENUR register where the error flags reside, a bit operation on it will reset the error flags.

Comparators

The device contains one differential comparator, with a pair of inputs (positive and negative) and an output. Ports I1-I3 is used for the comparator. The following is the Port I assignment:

- 11 Comparator negative input
- 12 Comparator positive input
- 13 Comparator output

A Comparator Select Register (CMPSL) is used to enable the comparator, read the output of the comparator internally, and enable the output of the comparator to the pins. Two control bits (enable and output enable) and one result bit are associated with the comparator. The comparator result bits (CMP1RD) is read only bits which will read as zero if the associated comparator is not enabled. The Comparator Select Register is cleared with reset, resulting in the comparator being disabled. The comparator should also be disabled before entering either the HALT or IDLE modes in order to save power. The configuration of the CMPSL register is as follows:

Comparators (Continued)

CMPSL REGISTER (ADDRESS X'00B7)

The CMPSL register contains the following bits:

CMP1EN Enable comparator

CMP1RD Comparator result (this is a read only bit, which

will read as 0 if the comparator is not enabled)

CMP10E Selects pin I3 as comparator output provided

that CMPIEN is set to enable the comparator

Unused	Unused	Unused	Unused	CMP10E	CMP1RD	CMP1EN	Unused
D:: -							D'1 0

Note that the two unused bits of CMPSL may be used as software flags.

Comparator output has the same spec as Ports L and G except that the rise and fall times are symmetrical.

Interrupts

The devices support a vectored interrupt scheme. It supports a total of fourteen interrupt sources. The following table lists all the possible device interrupt sources, their arbitration ranking and the memory locations reserved for the interrupt vector for each source.

Two bytes of program memory space are reserved for each interrupt source. All interrupt sources except the software interrupt are maskable. Each of the maskable interrupts have an Enable bit and a Pending bit. A maskable interrupt is active if its associated enable and pending bits are set. If GIE = 1 and an interrupt is active, then the processor will be interrupted as soon as it is ready to start executing an instruction except if the above conditions happen during the Software Trap service routine. This exception is described in the Software Trap sub-section.

The interruption process is accomplished with the INTR instruction (opcode 00), which is jammed inside the Instruction Register and replaces the opcode about to be executed. The following steps are performed for every interrupt:

- 1. The GIE (Global Interrupt Enable) bit is reset.
- 2. The address of the instruction about to be executed is pushed into the stack.
- 3. The PC (Program Counter) branches to address 00FF. This procedure takes 7 t_{C} cycles to execute.

Arbitration Ranking	Source	Description	Vector Address Hi-Low Byte
(1) Highest	Software	INTR Instruction	0yFE-0yFF
(2)	Reserved	for Future Use	0yFC-0yFD
(3)	External	Pin G0 Edge	0yFA-0yFB
(4)	Timer T0	Underflow	0yF8-0yF9
(5)	Timer T1	T1A/Underflow	0yF6-0yF7
(6)	Timer T1	T1B	0yF4-0yF5
(7)	MICROWIRE/PLUS	BUSY Goes Low	0yF2-0yF3
(8)	Reserved	for Future Use	0yF0-0yF1
(9)	UART	Receive	0yEE-0yEF
(10)	UART	Transmit	0yEC-0yED
(11)	Timer T2	T2A/Underflow	0yEA-0yEB
(12)	Timer T2	T2B	0yE8-0yE9
(13)	Timer T3	T3A/Underflow	0yE6-0yE7
(14)	Timer T3	ТЗВ	0yE4-0yE5
(15)	Port L/Wake Up	Port L Edge	0yE2-0yE3
(16) Lowest	Default	VIS Instr. Execution without Any Interrupts	0yE0-0yE1

y is VIS page, $y \neq 0$.

Interrupts (Continued)

At this time, since GIE=0, other maskable interrupts are disabled. The user is now free to do whatever context switching is required by saving the context of the machine in the stack with PUSH instructions. The user would then program a VIS (Vector Interrupt Select) instruction in order to branch to the interrupt service routine of the highest priority interrupt enabled and pending at the time of the VIS. Note that this is not necessarily the interrupt that caused the branch to address location 00FF Hex prior to the context switching.

Thus, if an interrupt with a higher rank than the one which caused the interruption becomes active before the decision of which interrupt to service is made by the VIS, then the interrupt with the higher rank will override any lower ones and will be acknowledged. The lower priority interrupt(s) are still pending, however, and will cause another interrupt immediately following the completion of the interrupt service routine associated with the higher priority interrupt just serviced. This lower priority interrupt will occur immediately following the RETI (Return from Interrupt) instruction at the end of the interrupt service routine just completed.

Inside the interrupt service routine, the associated pending bit has to be cleared by software. The RETI (Return from Interrupt) instruction at the end of the interrupt service routine will set the GIE (Global Interrupt Enable) bit, allowing the processor to be interrupted again if another interrupt is active and pending.

The VIS instruction looks at all the active interrupts at the time it is executed and performs an indirect jump to the beginning of the service routine of the one with the highest rank

The addresses of the different interrupt service routines, called vectors, are chosen by the user and stored in ROM in a table starting at 01E0 (assuming that VIS is located between 00FF and 01DF). The vectors are 15-bit wide and therefore occupy 2 ROM locations.

VIS and the vector table must be located in the same 256-byte block (0y00 to 0yFF) except if VIS is located at the last address of a block. In this case, the table must be in the next block. The vector table cannot be inserted in the first 256-byte block (y \neq 0).

The vector of the maskable interrupt with the lowest rank is located at 0yE0 (Hi-Order byte) and 0yE1 (Lo-Order byte) and so forth in increasing rank number. The vector of the maskable interrupt with the highest rank is located at 0yFA (Hi-Order byte) and 0yFB (Lo-Order byte).

The Software Trap has the highest rank and its vector is located at 0yFE and 0yFF.

If, by accident, a VIS gets executed and no interrupt is active, then the PC (Program Counter) will branch to a vector located at 0yE0-0yE1.

WARNING

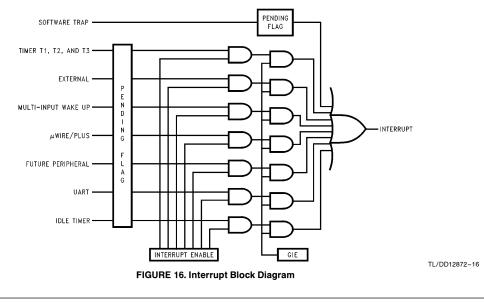
A Default VIS interrupt handler routine must be present. As a minimum, this handler should confirm that the GIE bit is cleared (this indicates that the interrupt sequence has been taken), take care of any required housekeeping, restore context and return. Some sort of Warm Restart procedure should be implemented. These events can occur without any error on the part of the system designer or programmer.

Note: There is always the possibility of an interrupt occurring during an instruction which is attempting to reset the GIE bit or any other interrupt enable bit. If this occurs when a single cycle instruction is being used to reset the interrupt enable bit, the interrupt enable bit will be reset but an interrupt may still occur. This is because interrupt processing is started at the same time as the interrupt bit is being reset. To avoid this scenario, the user should always use a two, three, or four cycle instruction to reset interrupt enable bits.

Figure 16 shows the Interrupt block diagram.

SOFTWARE TRAP

The Software Trap (ST) is a special kind of non-maskable interrupt which occurs when the INTR instruction (used to acknowledge interrupts) is fetched from ROM and placed inside the instruction register. This may happen when the PC is pointing beyond the available ROM address space or when the stack is over-popped.



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Interrupts (Continued)

When an ST occurs, the user can re-initialize the stack pointer and do a recovery procedure (similar to reset, but not necessarily containing all of the same initialization procedures) before restarting.

The occurrence of an ST is latched into the ST pending bit. The GIE bit is not affected and the ST pending bit (not accessible by the user) is used to inhibit other interrupts and to direct the program to the ST service routine with the VIS instruction. The RPND instruction is used to clear the software interrupt pending bit. This pending bit is also cleared on reset.

The ST has the highest rank among all interrupts.

Nothing (except another ST) can interrupt an ST being serviced.

WATCHDOG

The devices contain a WATCHDOG and clock monitor. The WATCHDOG is designed to detect the user program getting stuck in infinite loops resulting in loss of program control or "runaway" programs. The Clock Monitor is used to detect the absence of a clock or a very slow clock below a specified rate on the CKI pin.

The WATCHDOG consists of two independent logic blocks: WD UPPER and WD LOWER. WD UPPER establishes the upper limit on the service window and WD LOWER defines the lower limit of the service window.

Servicing the WATCHDOG consists of writing a specific value to a WATCHDOG Service Register named WDSVR which is memory mapped in the RAM. This value is composed of three fields, consisting of a 2-bit Window Select, a 5-bit Key Data field, and the 1-bit Clock Monitor Select field. Table V shows the WDSVR register.

The lower limit of the service window is fixed at 2048 instruction cycles. Bits 7 and 6 of the WDSVR register allow the user to pick an upper limit of the service window.

Table VI shows the four possible combinations of lower and upper limits for the WATCHDOG service window. This flexibility in choosing the WATCHDOG service window prevents any undue burden on the user software.

Bits 5, 4, 3, 2 and 1 of the WDSVR register represent the 5-bit Key Data field. The key data is fixed at 01100. Bit 0 of the WDSVR Register is the Clock Monitor Select bit.

TABLE V. WATCHDOG Service Register (WDSVR)

	dow ect	Key Data					Clock Monitor
Χ	Χ	0	1	1	0	0	Υ
7	6	5	4	3	2	1	0

TABLE VI. WATCHDOG Service Window Select

WDSVR Bit 7	WDSVR Bit 6	Service Window (Lower-Upper Limits)
0	0	2k-8k t _c Cycles
0	1	2k-16k t _c Cycles
1	0	2k-32k t _c Cycles
1	1	2k-64k t _c Cycles

Clock Monitor

The Clock Monitor aboard the device can be selected or deselected under program control. The Clock Monitor is guaranteed not to reject the clock if the instruction cycle clock ($1/t_c$) is greater or equal to 10 kHz. This equates to a clock input rate on CKI of greater or equal to 100 kHz.

WATCHDOG Operation

The WATCHDOG and Clock Monitor are disabled during reset. The device comes out of reset with the WATCHDOG armed, the WATCHDOG Window Select bits (bits 6, 7 of the WDSVR Register) set, and the Clock Monitor bit (bit 0 of the WDSVR Register) enabled. Thus, a Clock Monitor error will occur after coming out of reset, if the instruction cycle clock frequency has not reached a minimum specified value, including the case where the oscillator fails to start.

The WDSVR register can be written to only once after reset and the key data (bits 5 through 1 of the WDSVR Register) must match to be a valid write. This write to the WDSVR register involves two irrevocable choices: (i) the selection of the WATCHDOG service window (ii) enabling or disabling of the Clock Monitor. Hence, the first write to WDSVR Register involves selecting or deselecting the Clock Monitor, select the WATCHDOG service window and match the WATCHDOG key data. Subsequent writes to the WDSVR register will compare the value being written by the user to the WATCHDOG service window value and the key data (bits 7 through 1) in the WDSVR Register. Table VII shows the sequence of events that can occur.

The user must service the WATCHDOG at least once before the upper limit of the service window expires. The WATCHDOG may not be serviced more than once in every lower limit of the service window. The user may service the WATCHDOG as many times as wished in the time period between the lower and upper limits of the service window. The first write to the WDSVR Register is also counted as a WATCHDOG service.

The WATCHDOG has an output pin associated with it. This is the WDOUT pin, on pin 1 of the port G. WDOUT is active low. The WDOUT pin is in the high impedance state in the inactive state. Upon triggering the WATCHDOG, the logic will pull the WDOUT (G1) pin low for an additional $16\ t_{\rm C}-32\ t_{\rm C}$ cycles after the signal level on WDOUT pin goes below the lower Schmitt trigger threshold. After this delay, the device will stop forcing the WDOUT output low.

TABLE VII. WATCHDOG Service Actions

Key Data	Window Data	Clock Monitor	Action
Match	Match	Match	Valid Service: Restart Service Window
Don't Care	Mismatch	Don't Care	Error: Generate WATCHDOG Output
Mismatch	Don't Care	Don't Care	Error: Generate WATCHDOG Output
Don't Care	Don't Care	Mismatch	Error: Generate WATCHDOG Output

WATCHDOG Operation (Continued)

The WATCHDOG service window will restart when the WDOUT pin goes high. It is recommended that the user tie the WDOUT pin back to V_{CC} through a resistor in order to pull WDOUT high.

A WATCHDOG service while the WDOUT signal is active will be ignored. The state of the WDOUT pin is not guaranteed on reset, but if it powers up low then the WATCHDOG will time out and WDOUT will enter high impedance state.

The Clock Monitor forces the G1 pin low upon detecting a clock frequency error. The Clock Monitor error will continue until the clock frequency has reached the minimum specified value, after which the G1 output will enter the high impedance TRI-STATE mode following 16 $\rm t_c{-}32~t_c$ clock cycles. The Clock Monitor generates a continual Clock Monitor error if the oscillator fails to start, or fails to reach the minimum specified frequency. The specification for the Clock Monitor is as follows:

 $1/t_{\rm C} >$ 10 kHz—No clock rejection.

 $1/t_{\rm C} < 10$ Hz—Guaranteed clock rejection.

WATCHDOG AND CLOCK MONITOR SUMMARY

The following salient points regarding the WATCHDOG and CLOCK MONITOR should be noted:

- Both the WATCHDOG and CLOCK MONITOR detector circuits are inhibited during RESET.
- Following RESET, the WATCHDOG and CLOCK MONI-TOR are both enabled, with the WATCHDOG having the maximum service window selected.
- The WATCHDOG service window and CLOCK MONI-TOR enable/disable option can only be changed once, during the initial WATCHDOG service following RESET.
- The initial WATCHDOG service must match the key data value in the WATCHDOG Service register WDSVR in order to avoid a WATCHDOG error.
- Subsequent WATCHDOG services must match all three data fields in WDSVR in order to avoid WATCHDOG errors
- The correct key data value cannot be read from the WATCHDOG Service register WDSVR. Any attempt to read this key data value of 01100 from WDSVR will read as key data value of all 0's.

- The WATCHDOG detector circuit is inhibited during both the HALT and IDLE modes.
- The CLOCK MONITOR detector circuit is active during both the HALT and IDLE modes. Consequently, the COP888 inadvertently entering the HALT mode will be detected as a CLOCK MONITOR error (provided that the CLOCK MONITOR enable option has been selected by the program).
- With the single-pin R/C oscillator mask option selected and the CLKDLY bit reset, the WATCHDOG service window will resume following HALT mode from where it left off before entering the HALT mode.
- With the crystal oscillator mask option selected, or with the single-pin R/C oscillator mask option selected and the CLKDLY bit set, the WATCHDOG service window will be set to its selected value from WDSVR following HALT. Consequently, the WATCHDOG should not be serviced for at least 2048 instruction cycles following HALT, but must be serviced within the selected window to avoid a WATCHDOG error.
- The IDLE timer T0 is not initialized with RESET.
- The user can sync in to the IDLE counter cycle with an IDLE counter (T0) interrupt or by monitoring the T0PND flag. The T0PND flag is set whenever the thirteenth bit of the IDLE counter toggles (every 4096 instruction cycles).
 The user is responsible for resetting the T0PND flag.
- A hardware WATCHDOG service occurs just as the device exits the IDLE mode. Consequently, the WATCHDOG should not be serviced for at least 2048 instruction cycles following IDLE, but must be serviced within the selected window to avoid a WATCHDOG error.
- Following RESET, the initial WATCHDOG service (where the service window and the CLOCK MONITOR enable/ disable must be selected) may be programmed anywhere within the maximum service window (65,536 instruction cycles) initialized by RESET. Note that this initial WATCHDOG service may be programmed within the initial 2048 instruction cycles without causing a WATCHDOG error.

25

Detection of Illegal Conditions

The device can detect various illegal conditions resulting from coding errors, transient noise, power supply voltage drops, runaway programs, etc.

Reading of undefined ROM gets zeros. The opcode for software interrupt is zero. If the program fetches instructions from undefined ROM, this will force a software interrupt, thus signaling that an illegal condition has occurred.

The subroutine stack grows down for each call (jump to subroutine), interrupt, or PUSH, and grows up for each return or POP. The stack pointer is initialized to RAM location 06F Hex during reset. Consequently, if there are more returns than calls, the stack pointer will point to addresses 070 and 071 Hex (which are undefined RAM). Undefined RAM from addresses 070 to 07F (Segment 0), 140 to 17F (Segment 1), and all other segments (i.e., Segments 3 ... etc.) is read as all 1's, which in turn will cause the program to return to address 7FFF Hex. This is an undefined ROM location and the instruction fetched (all 0's) from this location will generate a software interrupt signaling an illegal condition.

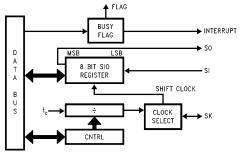
Thus, the chip can detect the following illegal conditions:

- 1. Executing from undefined ROM
- Over "POP"ing the stack by having more returns than calls.

When the software interrupt occurs, the user can re-initialize the stack pointer and do a recovery procedure before restarting (this recovery program is probably similar to that following reset, but might not contain the same program initialization procedures). The recovery program should reset the software interrupt pending bit using the RPND instruction.

MICROWIRE/PLUS

MICROWIRE/PLUS is a serial synchronous communications interface. The MICROWIRE/PLUS capability enables the device to interface with any of National Semiconductor's MICROWIRE peripherals (i.e. A/D converters, display drivers, E2PROMs etc.) and with other microcontrollers which sperial shift register (SIO) with serial data input (SI), serial data output (SO) and serial shift clock (SK). Figure 17 shows a block diagram of the MICROWIRE/PLUS logic.



TL/DD12872-17

FIGURE 17. MICROWIRE/PLUS Block Diagram

The shift clock can be selected from either an internal source or an external source. Operating the MICROWIRE/PLUS arrangement with the internal clock source is called the Master mode of operation. Similarly, operating the MICROWIRE/PLUS arrangement with an external shift clock is called the Slave mode of operation.

The CNTRL register is used to configure and control the MICROWIRE/PLUS mode. To use the MICROWIRE/PLUS, the MSEL bit in the CNTRL register is set to one. In the master mode, the SK clock rate is selected by the two bits, SL0 and SL1, in the CNTRL register. Table VIII details the different clock rates that may be selected.

TABLE VIII. MICROWIRE/PLUS Master Mode Clock Select

SL1	SL0	SK
0	0	$2 \times t_{c}$
0	1	$4 imes t_{c}$
1	х	$8 \times t_{c}$

Where t_{C} is the instruction cycle clock

MICROWIRE/PLUS (Continued)

MICROWIRE/PLUS OPERATION

Setting the BUSY bit in the PSW register causes the MICROWIRE/PLUS to start shifting the data. It gets reset when eight data bits have been shifted. The user may reset the BUSY bit by software to allow less than 8 bits to shift. If enabled, an interrupt is generated when eight data bits have been shifted. The device may enter the MICROWIRE/PLUS mode either as a Master or as a Slave. Figure 18 shows how two devices, microcontrollers and several peripherals may be interconnected using the MICROWIRE/PLUS arrangements.

Warning

The SIO register should only be loaded when the SK clock is low. Loading the SIO register while the SK clock is high will result in undefined data in the SIO register. SK clock is normally low when not shifting.

Setting the BUSY flag when the input SK clock is high in the MICROWIRE/PLUS slave mode may cause the current SK clock for the SIO shift register to be narrow. For safety, the BUSY flag should only be set when the input SK clock is low.

MICROWIRE/PLUS Master Mode Operation

In the MICROWIRE/PLUS Master mode of operation the shift clock (SK) is generated internally by the device. The MICROWIRE Master always initiates all data exchanges. The MSEL bit in the CNTRL register must be set to enable the SO and SK functions onto the G Port. The SO and SK pins must also be selected as outputs by setting appropriate bits in the Port G configuration register. Table IX summarizes the bit settings required for Master mode of operation.

MICROWIRE/PLUS Slave Mode Operation

In the MICROWIRE/PLUS Slave mode of operation the SK clock is generated by an external source. Setting the MSEL bit in the CNTRL register enables the SO and SK functions onto the G Port. The SK pin must be selected as an input and the SO pin is selected as an output pin by setting and resetting the appropriate bit in the Port G configuration register. Table IX summarizes the settings required to enter the Slave mode of operation.

The user must set the BUSY flag immediately upon entering the Slave mode. This will ensure that all data bits sent by the Master will be shifted properly. After eight clock pulses the BUSY flag will be cleared and the sequence may be repeated.

Alternate SK Phase Operation

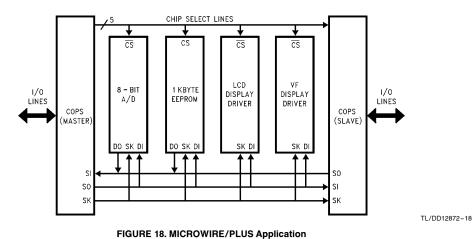
The device allows either the normal SK clock or an alternate phase SK clock to shift data in and out of the SIO register. In both the modes the SK is normally low. In the normal mode data is shifted in on the rising edge of the SK clock and the data is shifted out on the falling edge of the SK clock. The SIO register is shifted on each falling edge of the SK clock. In the alternate SK phase operation, data is shifted in on the falling edge of the SK clock and shifted out on the rising edge of the SK clock.

A control flag, SKSEL, allows either the normal SK clock or the alternate SK clock to be selected. Resetting SKSEL causes the MICROWIRE/PLUS logic to be clocked from the normal SK signal. Setting the SKSEL flag selects the alternate SK clock. The SKSEL is mapped into the G6 configuration bit. The SKSEL flag will power up in the reset condition, selecting the normal SK signal.

TABLE IX. MICROWIRE/PLUS Mode Selection

G4 (SO) Config. Bit	G5 (SK) Config. Bit	G4 Fun.	G5 Fun.	Operation
1	1	so		MICROWIRE/PLUS Master
0	1	TRI- STATE		MICROWIRE/PLUS Master
1	0	so		MICROWIRE/PLUS Slave
0	0	TRI- STATE		MICROWIRE/PLUS Slave

Note: This table assumes that the control flag MSEL is set.



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Memory Map

All RAM, ports and registers (except A and PC) are mapped into data memory address space.

Address S/ADD REG	Contents						
0000 to 006F	On-Chip RAM bytes (112 bytes)						
0070 to 007F	Unused RAM Address Space (Reads As All Ones)						
xx80 to xxAF	Unused RAM Address Space (Reads Undefined Data)						
xxB0	Timer T3 Lower Byte						
xxB1	Timer T3 Upper Byte						
xxB2	Timer T3 Autoload Register T3RA Lower Byte						
xxB3	Timer T3 Autoload Register T3RA Upper Byte						
xxB4	Timer T3 Autoload Register T3RB Lower Byte						
xxB5	Timer T3 Autoload Register T3RB Upper Byte						
xxB6	Timer T3 Control Register						
xxB7	Comparator Select Register (CMPSL)						
xxB8	UART Transmit Buffer (TBUF)						
xxB9	UART Receive Buffer (RBUF) UART Control and Status Register						
xxBA	(ENU)						
xxBB	UART Receive Control and Status Register (ENUR)						
xxBC	UART Interrupt and Clock Source Register (ENUI)						
xxBD	UART Baud Register (BAUD)						
xxBE	UART Prescale Select Register (PSR)						
xxBF	Reserved for UART						
xxC0	Timer T2 Lower Byte						
xxC1	Timer T2 Upper Byte						
xxC2	Timer T2 Autoload Register T2RA Lower Byte						
xxC3	Timer T2 Autoload Register T2RA Upper Byte						
xxC4	Timer T2 Autoload Register T2RB Lower Byte						
xxC5	Timer T2 Autoload Register T2RB Upper Byte						
xxC6	Timer T2 Control Register						
xxC7	WATCHDOG Service Register (Reg:WDSVR)						
xxC8	MIWU Edge Select Register (Reg:WKEDG)						
xxC9	MIWU Enable Register (Reg:WKEN)						
xxCA	MIWU Pending Register (Reg:WKPND)						
xxCB	Reserved						
xxCC	Reserved						
xxCD to xxCF	Reserved						

Address	Contents							
S/ADD REG								
xxD0	Port L Data Register							
xxD1	Port L Configuration Register							
xxD2	Port L Input Pins (Read Only)							
xxD3	Reserved for Port L							
xxD4	Port G Data Register							
xxD5	Port G Configuration Register							
xxD6	Port G Input Pins (Read Only)							
xxD7	Port I Input Pins (Read Only)							
xxD8	Reserved							
xxD9	Reserved							
xxDA	Reserved							
xxDB	Reserved							
xxDC	Port D							
xxDD to DF	Reserved for Port D							
xxE0 to xxE5	Reserved for EE Control Registers							
xxE6	Timer T1 Autoload Register T1RB							
	Lower Byte							
xxE7	Timer T1 Autoload Register T1RB							
	Upper Byte							
xxE8	ICNTRL Register							
xxE9	MICROWIRE/PLUS Shift Register							
xxEA	Timer T1 Lower Byte							
xxEB	Timer T1 Upper Byte							
xxEC	Timer T1 Autoload Register T1RA							
	Lower Byte							
xxED	Timer T1 Autoload Register T1RA							
	Upper Byte							
xxEE	CNTRL Control Register							
xxEF	PSW Register							
xxF0 to FB	On-Chip RAM Mapped as Registers							
xxFC	X Register							
xxFD	SP Register							
xxFE	B Register							
xxFF	S Register							
0100-017F	On-Chip 128 RAM Bytes							

Note: Reading memory locations 0070H–007FH (Segment 0) will return all ones. Reading unused memory locations 0080H–00AFH (Segment 0) will return undefined data. Reading memory locations from other Segments (i.e., Segment 2, Segment 3, ... etc.) will return all ones.

Addressing Modes

There are ten addressing modes, six for operand addressing and four for transfer of control.

OPERAND ADDRESSING MODES

Register Indirect

This is the "normal" addressing mode. The operand is the data memory addressed by the B pointer or X pointer.

Register Indirect (with auto post increment or decrement of pointer)

This addressing mode is used with the LD and X instructions. The operand is the data memory addressed by the B pointer or X pointer. This is a register indirect mode that automatically post increments or decrements the B or X register after executing the instruction.

Direct

The instruction contains an 8-bit address field that directly points to the data memory for the operand.

Immediate

The instruction contains an 8-bit immediate field as the operand.

Short Immediate

This addressing mode is used with the Load B Immediate instruction. The instruction contains a 4-bit immediate field as the operand.

Indirect

This addressing mode is used with the LAID instruction. The contents of the accumulator are used as a partial address (lower 8 bits of PC) for accessing a data operand from the program memory.

TRANSFER OF CONTROL ADDRESSING MODES

Relative

This mode is used for the JP instruction, with the instruction field being added to the program counter to get the new program location. JP has a range from -31 to +32 to allow a 1-byte relative jump (JP + 1 is implemented by a NOP instruction). There are no "pages" when using JP, since all 15 bits of PC are used.

Absolute

This mode is used with the JMP and JSR instructions, with the instruction field of 12 bits replacing the lower 12 bits of the program counter (PC). This allows jumping to any location in the current 4k program memory segment.

Absolute Long

This mode is used with the JMPL and JSRL instructions, with the instruction field of 15 bits replacing the entire 15 bits of the program counter (PC). This allows jumping to any location up to 32k in the program memory space.

Indirect

This mode is used with the JID instruction. The contents of the accumulator are used as a partial address (lower 8 bits of PC) for accessing a location in the program memory. The contents of this program memory location serve as a partial address (lower 8 bits of PC) for the jump to the next instruction

Note: The VIS is a special case of the Indirect Transfer of Control addressing mode, where the double byte vector associated with the interrupt is transferred from adjacent addresses in the program memory into the program counter (PC) in order to jump to the associated interrupt service routine.

Instruction Set

Register and Symbol Definition

	Registers						
Α	8-Bit Accumulator Register						
В	8-Bit Address Register						
X	8-Bit Address Register						
SP	8-Bit Stack Pointer Register						
PC	15-Bit Program Counter Register						
PU	Upper 7 Bits of PC						
PL	Lower 8 Bits of PC						
С	1 Bit of PSW Register for Carry						
HC	1 Bit of PSW Register for Half Carry						
GIE	1 Bit of PSW Register for Global						
	Interrupt Enable						
VU	Interrupt Vector Upper Byte						
VL	Interrupt Vector Lower Byte						

	Symbols							
[B]	Memory Indirectly Addressed by B Register							
[X]	Memory Indirectly Addressed by X Register							
MD	Direct Addressed Memory							
Mem	Direct Addressed Memory or [B]							
Meml	Direct Addressed Memory or [B] or Immediate Data							
lmm	8-Bit Immediate Data							
Reg	Register Memory: Addresses F0 to FF (Includes B, X and SP)							
Bit	Bit Number (0 to 7)							
←	Loaded with							
\longleftrightarrow	Exchanged with							

Instruction Set (Continued)

INSTRUCTION SET

ADD	A,Meml	ADD	A ← A + Meml
ADC	A,Meml	ADD with Carry	$A \leftarrow A + Meml + C, C \leftarrow Carry,$
		•	HC ← Half Carry
SUBC	A,Meml	Subtract with Carry	$A \leftarrow A - Meml + C, C \leftarrow Carry,$
		-	HC ← Half Carry
AND	A,Meml	Logical AND	A ← A and Meml
ANDSZ	A,Imm	Logical AND Immed., Skip if Zero	Skip next if (A and Imm) = 0
OR	A,Meml	Logical OR	A ← A or Meml
XOR	A,Meml	Logical EXclusive OR	A ← A xor Meml
IFEQ	MD,Imm	IF EQual	Compare MD and Imm, Do next if MD = Imm
IFEQ	A,Meml	IF EQual	Compare A and Meml, Do next if $A = Meml$
IFNE	A,Meml	IF Not Equal	Compare A and Meml, Do next if $A \neq Meml$
IFGT	A,Meml	IF Greater Than	Compare A and Meml, Do next if A > Meml
IFBNE	#	If B Not Equal	Do next if lower 4 bits of B \neq Imm
DRSZ	Reg	Decrement Reg., Skip if Zero	$Reg \leftarrow Reg - 1$, Skip if $Reg = 0$
SBIT	#,Mem	Set BIT	1 to bit, Mem (bit = 0 to 7 immediate)
RBIT	#,Mem	Reset BIT	0 to bit, Mem
IFBIT	#,Mem	IF BIT	If bit in A or Mem is true do next instruction
RPND		Reset PeNDing Flag	Reset Software Interrupt Pending Flag
X	A,Mem	EXchange A with Memory	A ←→ Mem
X	A,[X]	EXchange A with Memory [X]	$A \longleftrightarrow [X]$
LD	A,Meml	LoaD A with Memory	A ← Meml
LD	A,[X]	LoaD A with Memory [X]	$A \leftarrow [X]$
LD	B,Imm	LoaD B with Immed.	B ← Imm
LD	Mem,Imm	LoaD Memory Immed.	Mem ← Imm
LD	Reg,Imm	LoaD Register Memory Immed.	Reg ← Imm
X	A, [B ±]	EXchange A with Memory [B]	$A \longleftrightarrow [B], (B \longleftarrow B \pm 1)$
X	A, $[X \pm]$	EXchange A with Memory [X]	$A \longleftrightarrow [X], (X \leftarrow \pm 1)$
LD	A, $[B\pm]$	LoaD A with Memory [B]	$A \leftarrow [B], (B \leftarrow B \pm 1)$
LD	A, $[X \pm]$	LoaD A with Memory [X]	$A \leftarrow [X], (X \leftarrow X \pm 1)$
LD	$[B\pm]$,Imm	LoaD Memory [B] Immed.	$[B] \leftarrow Imm, (B \leftarrow B \pm 1)$
CLR	Α	CLeaR A	A ← 0
INC	Α	INCrement A	A ← A + 1
DEC	Α	DECrementA	A ← A − 1
LAID		Load A InDirect from ROM	A ← ROM (PU,A)
DCCD			
DCOR	Α	Decimal CORrect A	A ← BCD correction of A (follows ADC, SUBC)
RRC	A A	Decimal CORrect A Rotate A Right thru C	A \leftarrow BCD correction of A (follows ADC, SUBC) C \rightarrow A7 \rightarrow \rightarrow A0 \rightarrow C
		Rotate A Right thru C Rotate A Left thru C	$C \rightarrow A7 \rightarrow \dots \rightarrow A0 \rightarrow C$ $C \leftarrow A7 \leftarrow \dots \leftarrow A0 \leftarrow C$
RRC RLC SWAP	Α	Rotate A Right thru C Rotate A Left thru C SWAP nibbles of A	$C \rightarrow A7 \rightarrow \dots \rightarrow A0 \rightarrow C$ $C \leftarrow A7 \leftarrow \dots \leftarrow A0 \leftarrow C$ $A7 \dots A4 \longleftrightarrow A3 \dots A0$
RRC RLC SWAP SC	A A	Rotate A Right thru C Rotate A Left thru C SWAP nibbles of A Set C	$C \rightarrow A7 \rightarrow \dots \rightarrow A0 \rightarrow C$ $C \leftarrow A7 \leftarrow \dots \leftarrow A0 \leftarrow C$ $A7 \dots A4 \longleftrightarrow A3 \dots A0$ $C \leftarrow 1, HC \leftarrow 1$
RRC RLC SWAP SC RC	A A	Rotate A Right thru C Rotate A Left thru C SWAP nibbles of A Set C Reset C	$C \rightarrow A7 \rightarrow \dots \rightarrow A0 \rightarrow C$ $C \leftarrow A7 \leftarrow \dots \leftarrow A0 \leftarrow C$ $A7 \dots A4 \longleftrightarrow A3 \dots A0$ $C \leftarrow 1, HC \leftarrow 1$ $C \leftarrow 0, HC \leftarrow 0$
RRC RLC SWAP SC RC IFC	A A	Rotate A Right thru C Rotate A Left thru C SWAP nibbles of A Set C Reset C IF C	$C \rightarrow A7 \rightarrow \dots \rightarrow A0 \rightarrow C$ $C \leftarrow A7 \leftarrow \dots \leftarrow A0 \leftarrow C$ $A7 \dots A4 \longleftrightarrow A3 \dots A0$ $C \leftarrow 1, HC \leftarrow 1$ $C \leftarrow 0, HC \leftarrow 0$ IF C is true, do next instruction
RRC RLC SWAP SC RC IFC IFNC	A A A	Rotate A Right thru C Rotate A Left thru C SWAP nibbles of A Set C Reset C IF C IF Not C	$C \rightarrow A7 \rightarrow \dots \rightarrow A0 \rightarrow C$ $C \leftarrow A7 \leftarrow \dots \leftarrow A0 \leftarrow C$ $A7 \dots A4 \leftarrow A3 \dots A0$ $C \leftarrow 1, HC \leftarrow 1$ $C \leftarrow 0, HC \leftarrow 0$ IF C is true, do next instruction If C is not true, do next instruction
RRC RLC SWAP SC RC IFC IFNC POP	A A A	Rotate A Right thru C Rotate A Left thru C SWAP nibbles of A Set C Reset C IF C IF Not C POP the stack into A	$C \rightarrow A7 \rightarrow \dots \rightarrow A0 \rightarrow C$ $C \leftarrow A7 \leftarrow \dots \leftarrow A0 \leftarrow C$ $A7 \dots A4 \leftarrow A3 \dots A0$ $C \leftarrow 1, HC \leftarrow 1$ $C \leftarrow 0, HC \leftarrow 0$ IF C is true, do next instruction If C is not true, do next instruction $SP \leftarrow SP + 1, A \leftarrow [SP]$
RRC RLC SWAP SC RC IFC IFNC	A A A	Rotate A Right thru C Rotate A Left thru C SWAP nibbles of A Set C Reset C IF C IF Not C	$C \rightarrow A7 \rightarrow \dots \rightarrow A0 \rightarrow C$ $C \leftarrow A7 \leftarrow \dots \leftarrow A0 \leftarrow C$ $A7 \dots A4 \leftarrow A3 \dots A0$ $C \leftarrow 1, HC \leftarrow 1$ $C \leftarrow 0, HC \leftarrow 0$ IF C is true, do next instruction If C is not true, do next instruction
RRC RLC SWAP SC RC IFC IFNC POP PUSH	A A A	Rotate A Right thru C Rotate A Left thru C SWAP nibbles of A Set C Reset C IF C IF Not C POP the stack into A	$C \rightarrow A7 \rightarrow \dots \rightarrow A0 \rightarrow C$ $C \leftarrow A7 \leftarrow \dots \leftarrow A0 \leftarrow C$ $A7 \dots A4 \leftarrow A3 \dots A0$ $C \leftarrow 1, HC \leftarrow 1$ $C \leftarrow 0, HC \leftarrow 0$ IF C is true, do next instruction If C is not true, do next instruction $SP \leftarrow SP + 1, A \leftarrow [SP]$ $[SP] \leftarrow A, SP \leftarrow SP - 1$ $PU \leftarrow [VU], PL \leftarrow [VL]$
RRC RLC SWAP SC RC IFC IFNC POP PUSH VIS JMPL	A A A A Addr.	Rotate A Right thru C Rotate A Left thru C SWAP nibbles of A Set C Reset C IF C IF Not C POP the stack into A PUSH A onto the stack Vector to Interrupt Service Routine Jump absolute Long	$C \rightarrow A7 \rightarrow \dots \rightarrow A0 \rightarrow C$ $C \leftarrow A7 \leftarrow \dots \leftarrow A0 \leftarrow C$ $A7 \dots A4 \longleftrightarrow A3 \dots A0$ $C \leftarrow 1, HC \leftarrow 1$ $C \leftarrow 0, HC \leftarrow 0$ IF C is true, do next instruction If C is not true, do next instruction $SP \leftarrow SP + 1, A \leftarrow [SP]$ $[SP] \leftarrow A, SP \leftarrow SP - 1$ $PU \leftarrow [VU], PL \leftarrow [VL]$ $PC \leftarrow ii (ii = 15 \text{ bits, 0k to } 32\text{k})$
RRC RLC SWAP SC RC IFC IFNC POP PUSH VIS JMPL JMP	A A A A Addr. Addr.	Rotate A Right thru C Rotate A Left thru C SWAP nibbles of A Set C Reset C IF C IF Not C POP the stack into A PUSH A onto the stack Vector to Interrupt Service Routine Jump absolute Long Jump absolute	$C \rightarrow A7 \rightarrow \dots \rightarrow A0 \rightarrow C$ $C \leftarrow A7 \leftarrow \dots \leftarrow A0 \leftarrow C$ $A7 \dots A4 \longleftrightarrow A3 \dots A0$ $C \leftarrow 1, HC \leftarrow 1$ $C \leftarrow 0, HC \leftarrow 0$ $IF C is true, do next instruction$ $If C is not true, do next instruction$ $SP \leftarrow SP + 1, A \leftarrow [SP]$ $[SP] \leftarrow A, SP \leftarrow SP - 1$ $PU \leftarrow [VU], PL \leftarrow [VL]$ $PC \leftarrow ii (ii = 15 bits, 0k to 32k)$ $PC9 \dots 0 \leftarrow i (i = 12 bits)$
RRC RLC SWAP SC RC IFC IFNC POP PUSH VIS JMPL JMP JP	A A A A Addr. Addr. Disp.	Rotate A Right thru C Rotate A Left thru C SWAP nibbles of A Set C Reset C IF C IF Not C POP the stack into A PUSH A onto the stack Vector to Interrupt Service Routine Jump absolute Jump relative short	$C \rightarrow A7 \rightarrow \dots \rightarrow A0 \rightarrow C$ $C \leftarrow A7 \leftarrow \dots \leftarrow A0 \leftarrow C$ $A7 \dots A4 \longleftrightarrow A3 \dots A0$ $C \leftarrow 1, HC \leftarrow 1$ $C \leftarrow 0, HC \leftarrow 0$ IF C is true, do next instruction If C is not true, do next instruction $SP \leftarrow SP + 1, A \leftarrow [SP]$ $[SP] \leftarrow A, SP \leftarrow SP - 1$ $PU \leftarrow [VU], PL \leftarrow [VL]$ $PC \leftarrow ii (ii = 15 bits, 0k to 32k)$ $PC9 \dots 0 \leftarrow i (i = 12 bits)$ $PC \leftarrow PC + r (r is -31 to +32, except 1)$
RRC RLC SWAP SC RC IFC IFNC POP PUSH VIS JMPL JP JSRL	A A A A A A A A A A A A A A A A A A A	Rotate A Right thru C Rotate A Left thru C SWAP nibbles of A Set C Reset C IF C IF Not C POP the stack into A PUSH A onto the stack Vector to Interrupt Service Routine Jump absolute Jump relative short Jump SubRoutine Long	$C \rightarrow A7 \rightarrow \dots \rightarrow A0 \rightarrow C$ $C \leftarrow A7 \leftarrow \dots \leftarrow A0 \leftarrow C$ $A7 \dots A4 \longleftrightarrow A3 \dots A0$ $C \leftarrow 1, HC \leftarrow 1$ $C \leftarrow 0, HC \leftarrow 0$ IF C is true, do next instruction If C is not true, do next instruction $SP \leftarrow SP + 1, A \leftarrow [SP]$ $[SP] \leftarrow A, SP \leftarrow SP - 1$ $PU \leftarrow [VU], PL \leftarrow [VL]$ $PC \leftarrow ii (ii = 15 bits, 0k to 32k)$ $PC9 \dots 0 \leftarrow i (i = 12 bits)$ $PC \leftarrow PC + r (r is -31 to +32, except 1)$ $[SP] \leftarrow PL, [SP-1] \leftarrow PU, SP -2, PC \leftarrow ii$
RRC RLC SWAP SC RC IFC IFNC POP PUSH VIS JMPL JMP JP JSRL JSR	A A A A Addr. Addr. Disp.	Rotate A Right thru C Rotate A Left thru C SWAP nibbles of A Set C Reset C IF C IF Not C POP the stack into A PUSH A onto the stack Vector to Interrupt Service Routine Jump absolute Long Jump relative short Jump SubRoutine Long Jump SubRoutine	$C \rightarrow A7 \rightarrow \dots \rightarrow A0 \rightarrow C$ $C \leftarrow A7 \leftarrow \dots \leftarrow A0 \leftarrow C$ $A7 \dots A4 \leftrightarrow A3 \dots A0$ $C \leftarrow 1, HC \leftarrow 1$ $C \leftarrow 0, HC \leftarrow 0$ IF C is true, do next instruction If C is not true, do next instruction $SP \leftarrow SP + 1, A \leftarrow [SP]$ $[SP] \leftarrow A, SP \leftarrow SP - 1$ $PU \leftarrow [VU], PL \leftarrow [VL]$ $PC \leftarrow ii (ii = 15 \text{ bits, 0k to } 32\text{k})$ $PC9 \dots 0 \leftarrow i (i = 12 \text{ bits})$ $PC \leftarrow PC + r (r \text{ is } -31 \text{ to } +32, \text{ except } 1)$ $[SP] \leftarrow PL, [SP - 1] \leftarrow PU, SP - 2, PC \leftarrow ii$ $[SP] \leftarrow PL, [SP - 1] \leftarrow PU, SP - 2, PC9 \dots 0 \leftarrow i$
RRC RLC SWAP SC RC IFC IFNC POP PUSH VIS JMPL JMP JSRL JSR JID	A A A A A A A A A A A A A A A A A A A	Rotate A Right thru C Rotate A Left thru C SWAP nibbles of A Set C Reset C IF C IF Not C POP the stack into A PUSH A onto the stack Vector to Interrupt Service Routine Jump absolute Long Jump absolute Jump Fulative short Jump SubRoutine Jump InDirect	$C \rightarrow A7 \rightarrow \dots \rightarrow A0 \rightarrow C$ $C \leftarrow A7 \leftarrow \dots \leftarrow A0 \leftarrow C$ $A7 \dots A4 \leftarrow A3 \dots A0$ $C \leftarrow 1, HC \leftarrow 1$ $C \leftarrow 0, HC \leftarrow 0$ $IF C is true, do next instruction$ $If C is not true, do next instruction$ $SP \leftarrow SP + 1, A \leftarrow [SP]$ $[SP] \leftarrow A, SP \leftarrow SP - 1$ $PU \leftarrow [VU], PL \leftarrow [VL]$ $PC \leftarrow ii (ii = 15 \text{ bits, 0k to 32k})$ $PC9 \dots 0 \leftarrow i (i = 12 \text{ bits})$ $PC \leftarrow PC + r (r \text{ is } -31 \text{ to } +32, \text{ except 1})$ $[SP] \leftarrow PL, [SP - 1] \leftarrow PU, SP - 2, PC \leftarrow ii$ $[SP] \leftarrow PL, [SP - 1] \leftarrow PU, SP - 2, PC9 \dots 0 \leftarrow i$ $PL \leftarrow ROM (PU, A)$
RRC RLC SWAP SC RC IFC IFNC POP PUSH VIS JMPL JMP JP JSRL JSR JID RET	A A A A A A A A A A A A A A A A A A A	Rotate A Right thru C Rotate A Left thru C SWAP nibbles of A Set C Reset C IF C IF Not C POP the stack into A PUSH A onto the stack Vector to Interrupt Service Routine Jump absolute Long Jump absolute Jump Felative short Jump SubRoutine Jump SubRoutine Jump InDirect RETurn from subroutine	$C \rightarrow A7 \rightarrow \dots \rightarrow A0 \rightarrow C$ $C \leftarrow A7 \leftarrow \dots \leftarrow A0 \leftarrow C$ $A7 \dots A4 \longleftrightarrow A3 \dots A0$ $C \leftarrow 1, HC \leftarrow 1$ $C \leftarrow 0, HC \leftarrow 0$ $IF C is true, do next instruction$ $If C is not true, do next instruction$ $SP \leftarrow SP + 1, A \leftarrow [SP]$ $[SP] \leftarrow A, SP \leftarrow SP - 1$ $PU \leftarrow [VU], PL \leftarrow [VL]$ $PC \leftarrow ii (ii = 15 \text{ bits, 0k to 32k})$ $PC9 \dots 0 \leftarrow i (i = 12 \text{ bits})$ $PC \leftarrow PC + r (r \text{ is } -31 \text{ to } +32, \text{ except } 1)$ $[SP] \leftarrow PL, [SP - 1] \leftarrow PU, SP - 2, PC \leftarrow ii$ $[SP] \leftarrow PL, [SP - 1] \leftarrow PU, SP - 2, PC9 \dots 0 \leftarrow i$ $PL \leftarrow ROM (PU, A)$ $SP + 2, PL \leftarrow [SP], PU \leftarrow [SP - 1]$
RRC RLC SWAP SC RC IFC IFNC POP PUSH VIS JMPL JMP JP JSRL JSR JID RET RETSK	A A A A A A A A A A A A A A A A A A A	Rotate A Right thru C Rotate A Left thru C SWAP nibbles of A Set C Reset C IF C IF Not C POP the stack into A PUSH A onto the stack Vector to Interrupt Service Routine Jump absolute Long Jump absolute Jump relative short Jump SubRoutine Long Jump SubRoutine Jump InDirect RETurn from subroutine RETurn and SKip	$C \rightarrow A7 \rightarrow \dots \rightarrow A0 \rightarrow C$ $C \leftarrow A7 \leftarrow \dots \leftarrow A0 \leftarrow C$ $A7 \dots A4 \longleftrightarrow A3 \dots A0$ $C \leftarrow 1, HC \leftarrow 1$ $C \leftarrow 0, HC \leftarrow 0$ $IF C is true, do next instruction$ $If C is not true, do next instruction$ $SP \leftarrow SP + 1, A \leftarrow [SP]$ $[SP] \leftarrow A, SP \leftarrow SP - 1$ $PU \leftarrow [VU], PL \leftarrow [VL]$ $PC \leftarrow ii (ii = 15 \text{ bits, 0k to 32k})$ $PC9 \dots 0 \leftarrow i (i = 12 \text{ bits})$ $PC \leftarrow PC + r (r \text{ is } -31 \text{ to } +32, \text{ except } 1)$ $[SP] \leftarrow PL, [SP - 1] \leftarrow PU, SP - 2, PC \leftarrow ii$ $[SP] \leftarrow PL, [SP - 1] \rightarrow PU, SP - 2, PC9 \dots 0 \leftarrow i$ $PL \leftarrow ROM (PU, A)$ $SP + 2, PL \leftarrow [SP], PU \leftarrow [SP - 1]$ $SP + 2, PL \leftarrow [SP], PU \leftarrow [SP - 1]$
RRC RLC SWAP SC RC IFC IFNC POP PUSH VIS JMP JP JSRL JSR JID RET RETSK RETI	A A A A A A A A A A A A A A A A A A A	Rotate A Right thru C Rotate A Left thru C SWAP nibbles of A Set C Reset C IF C IF Not C POP the stack into A PUSH A onto the stack Vector to Interrupt Service Routine Jump absolute Long Jump absolute Long Jump SubRoutine Jump SubRoutine Jump InDirect RETurn from subroutine RETurn from Interrupt	$C \rightarrow A7 \rightarrow \dots \rightarrow A0 \rightarrow C$ $C \leftarrow A7 \leftarrow \dots \leftarrow A0 \leftarrow C$ $A7 \dots A4 \longleftrightarrow A3 \dots A0$ $C \leftarrow 1, HC \leftarrow 1$ $C \leftarrow 0, HC \leftarrow 0$ $IF C is true, do next instruction$ $If C is not true, do next instruction$ $SP \leftarrow SP + 1, A \leftarrow [SP]$ $[SP] \leftarrow A, SP \leftarrow SP - 1$ $PU \leftarrow [VU], PL \leftarrow [VL]$ $PC \leftarrow ii (ii = 15 bits), 0k to 32k)$ $PC9 \dots 0 \leftarrow i (i = 12 bits)$ $PC \leftarrow PC + r (r is -31 to +32, except 1)$ $[SP] \leftarrow PL, [SP-1] \leftarrow PU, SP-2, PC \leftarrow ii$ $[SP] \leftarrow PL, [SP-1] \leftarrow PU, SP-2, PC9 \dots 0 \leftarrow i$ $PL \leftarrow ROM (PU,A)$ $SP + 2, PL \leftarrow [SP], PU \leftarrow [SP-1]$ $SP + 2, PL \leftarrow [SP], PU \leftarrow [SP-1]$ $SP + 2, PL \leftarrow [SP], PU \leftarrow [SP-1]$
RRC RLC SWAP SC RC IFC IFNC POP PUSH VIS JMPL JMP JP JSRL JSR JID RET RETSK	A A A A A A A A A A A A A A A A A A A	Rotate A Right thru C Rotate A Left thru C SWAP nibbles of A Set C Reset C IF C IF Not C POP the stack into A PUSH A onto the stack Vector to Interrupt Service Routine Jump absolute Long Jump absolute Jump relative short Jump SubRoutine Long Jump SubRoutine Jump InDirect RETurn from subroutine RETurn and SKip	$C \rightarrow A7 \rightarrow \dots \rightarrow A0 \rightarrow C$ $C \leftarrow A7 \leftarrow \dots \leftarrow A0 \leftarrow C$ $A7 \dots A4 \longleftrightarrow A3 \dots A0$ $C \leftarrow 1, HC \leftarrow 1$ $C \leftarrow 0, HC \leftarrow 0$ $IF C is true, do next instruction$ $If C is not true, do next instruction$ $SP \leftarrow SP + 1, A \leftarrow [SP]$ $[SP] \leftarrow A, SP \leftarrow SP - 1$ $PU \leftarrow [VU], PL \leftarrow [VL]$ $PC \leftarrow ii (ii = 15 \text{ bits, 0k to } 32\text{k})$ $PC9 \dots 0 \leftarrow i (i = 12 \text{ bits})$ $PC \leftarrow PC + r (r \text{ is } -31 \text{ to } +32, \text{ except } 1)$ $[SP] \leftarrow PL, [SP - 1] \leftarrow PU, SP - 2, PC \leftarrow ii$ $[SP] \leftarrow PL, [SP - 1] \rightarrow PU, SP - 2, PC9 \dots 0 \leftarrow i$ $PL \leftarrow ROM (PU, A)$ $SP + 2, PL \leftarrow [SP], PU \leftarrow [SP - 1]$ $SP + 2, PL \leftarrow [SP], PU \leftarrow [SP - 1]$

Instruction Execution Time

Most instructions are single byte (with immediate addressing mode instructions taking two bytes).

Most single byte instructions take one cycle time to execute.

Skipped instructions require x number of cycles to be skipped, where x equals the number of bytes in the skipped instruction opcode.

See the BYTES and CYCLES per INSTRUCTION table for details.

Bytes and Cycles per Instruction

The following table shows the number of bytes and cycles for each instruction in the format of byte/cycle.

Logic and Arithmetic Instructions

	[B]	Direct	Immed.			
ADD	1/1	3/4	2/2			
ADC	1/1	3/4	2/2			
SUBC	1/1	3/4	2/2			
AND	1/1	3/4	2/2			
OR	1/1	3/4	2/2			
XOR	1/1	3/4	2/2			
IFEQ	1/1	3/4	2/2			
IFGT	1/1	3/4	2/2			
IFBNE	1/1					
DRSZ		1/3				
SBIT	1/1	3/4				
RBIT	1/1	3/4				
IFBIT	1/1	3/4				
	•	•	•			

Instructions Using A and C

	•
CLRA	1/1
INCA	1/1
DECA	1/1
LAID	1/3
DCORA	1/1
RRCA	1/1
RLCA	1/1
SWAPA	1/1
SC	1/1
RC	1/1
IFC	1/1
IFNC	1/1
PUSHA	1/3
POPA	1/3
ANDSZ	2/2

Transfer of Control Instructions

mon donono							
JMPL	3/4						
JMP	2/3						
JP	1/3						
JSRL	3/5						
JSR	2/5						
JID	1/3						
VIS	1/5						
RET	1/5						
RETSK	1/5						
RETI	1/5						
INTR	1/7						
NOP	1/1						

RPND	1/1

Memory Transfer Instructions

	Register Indirect		Direct	Immed.	Register Indirect Auto Incr. and Decr.			
	[B]	[X]			[B+,B-]	[X +, X -]		
X A,*	1/1	1/3	2/3		1/2	1/3		
LD A,*	1/1	1/3	2/3	2/2	1/2	1/3		
LD B, Imm				1/1				
LD B, Imm				2/3				
LD Mem, Imm	2/2	2/2	3/3		2/2			
LD Reg, Imm			2/3					
IFEQ MD, Imm			3/3					

(IF B < 16) (IF B > 15)

^{* = &}gt; Memory location addressed by B or X or directly.

										LOWE	R NIBI	BLE								
			0	-	2	3	4	5	9	7	8	6	4	В	O e	O 1	Ш	Ц		
		0	INTR	JP + 2	JP + 3	JP + 4	JP + 5	JP + 6	JP + 7	JP + 8	JP + 9	JP + 10	JP + 11	JP + 12	JP + 13	JP + 14	JP + 15	JP + 16		
		1	JP + 17	JP + 18	JP + 19	JP + 20	JP + 21	JP + 22	JP + 23	JP + 24	JP + 25	JP + 26	JP + 27	JP + 28	JP + 29	JP + 30	JP + 31	JP + 32		
		2	JMP x000-x0FF	JMP x100-x1FF	JMP x200-x2FF	JMP x300-x3FF	JMP x400-x4FF	JMP x500-x5FF	JMP x600-x6FF	JMP ×700-×7FF	JMP x800-x8FF	JMP x900-x9FF	JMP xA00-xAFF	JMP xB00-xBFF	JMP xC00-xCFF	JMP xD00-xDFF	JMP xE00-xEFF	JMP xF00-xFFF		
		3	JSR x000-x0FF	JSR x100-x1FF	JSR x200-x2FF	JSR x300-x3FF	JSR x400-x4FF	JSR x500-x5FF	JSR x600-x6FF	JSR x700-x7FF	JSR x800-x8FF	JSR x900-x9FF	JSR xA00-xAFF	JSR xB00-xBFF	JSR xC00-xCFF	JSR xD00-xDFF	JSR xE00-xEFF	JSR xF00-xFFF		
		4	IFBNE 0	IFBNE 1	IFBNE 2	IFBNE 3	IFBNE 4	IFBNE 5	IFBNE 6	IFBNE 7	IFBNE 8	IFBNE 9	IFBNE 0A	IFBNE 0B	IFBNE 0C	IFBNE 0D	IFBNE 0E	IFBNE 0F		
		2	LD B, # 0F	LD B, # 0E	LD B,#0D	LD B, # 0C	LD B, # 0B	LD B, # 0A	LD B, # 09	LD B, # 08	LD B, # 07	LD B, # 06	LD B,#05	LD B, # 04	LD B,#03	LD B,# 02	LD B, # 01	LD B, # 00		
	BLE	9	ANDSZ A, #i	*	*	*	CLRA	SWAPA	DCORA	PUSHA	RBIT 0,[B]	RBIT 1,[B]	RBIT 2,[B]	RBIT 3,[B]	RBIT 4,[B]	RBIT 5,[B]	RBIT 6,[B]	RBIT 7,[B]		
	UPPER NIBBLE	7	IFBIT 0,[B]	IFBIT 1,[B]	IFBIT 2,[B]	IFBIT 3,[B]	IFBIT 4,[B]	IFBIT 5,[B]	IFBIT 6,[B]	IFBIT 7,[B]	SBIT 0,[B]	SBIT 1,[B]	SBIT 2,[B]	SBIT 3,[B]	SBIT 4,[B]	SBIT 5,[B]	SBIT 6,[B]	SBIT 7,[B]		
	UPPE	8	ADC A,[B]	SUB A,[B]	IFEQ A,[B]	IFGT A,[B]	ADD A,[B]	AND A,[B]	XOR A,[B]	OR A,[B]	IFC	IFNC	INCA	DECA	POPA	RETSK	RET	RETI		
Table		6	ADC A, # i	SUBC A, #i	IFEQ A, #i	IFGT A, #i	ADD A, # i	AND A, #i	XOR A, #i	OR A, #i	LD A, # i	IFNE A,#i	LD [B+], #i	LD [B-], #i	X A,Md	LD A,Md	LD [B],#i	LDB,#i		
code.		۷	RC	SC	X A, [B+]	X A, [B-]	LAID	all	X A,[B]	*	RLCA	IFEQ Md,#i	LD A, [B+]	LD A, [B-]	JMPL	JSBL	LD A,[B]	*		
4EG Opcode Table		8	RRCA	*	X A, [X+]	X A, [X-]	NIS	RPND	X A, [X]	*	NOP	IFNE A,[B]	LD A, [X+]	LD A, [X—]	LD Md, # i	DIR	LD A,[X]	*	location	BIT #i,A
JP87L8		ပ	DRSZ 0F0	DRSZ 0F1	DRSZ 0F2	DRSZ 0F3	DRSZ 0F4	DRSZ 0F5	DRSZ 0F6	DRSZ 0F7	DRSZ 0F8	DRSZ 0F9	DRSZ 0FA	DRSZ 0FB	DRSZ 0FC	DRSZ 0FD	DRSZ 0FE	DRSZ 0FF	ssed memory	ne opcode for IF
COP87L88EG/COP87L8		٥	LD 0F0, #i	LD 0F1, #i	LD 0F2, #i	LD 0F3, #i	LD 0F4, #i	LD 0F5, #i	LD 0F6, #i	LD 0F7, #i	LD 0F8, #i	LD 0F9, #i	LD 0FA, #i	LD 0FB, #i	LD 0FC, #i	LD 0FD, #i	LD 0FE, #i	LD 0FF, #i	i is the immediate data Md is a directly addressed memory location * is an unused opcode	Note: The opcode 60 Hex is also the opcode for IFBIT # i,A
P87L8		Ш	JP -31	JP -30	JP -29	JP -28	JP -27	JP -26	JP -25	JP -24	JP -23	JP -22	JP -21	JP -20	JP - 19	JP - 18	JP -17	JP -16	i is the im Md is a d * is an ur	The opcode
000		ш	JP - 15	JP 14	JP - 13	JP - 12	JP - 11	JP 10	9- AL	JP8	JP -7	JP6	JP5	JP4	JP -3	JP -2	JP 1	JP -0	where,	Note:

Development Support

SUMMARY

- iceMASTERTM: IM-COP8/400—Full feature in-circuit emulation for all COP8 products. A full set of COP8 Basic and Feature Family device and package specific probes are available.
- COP8 Debug Module: Moderate cost in-circuit emulation and development programming unit.
- COP8 Evaluation and Programming Unit: EPU-COP888GG—low cost in-circuit simulation and development programming unit.
- Assembler: COP8-DEV-IBMA. A DOS installable cross development Assembler, Linker, Librarian and Utility Software Development Tool Kit.
- C Compiler: COP8C. A DOS installable cross development Software Tool Kit.
- OTP/EPROM Programmer Support: Covering needs from engineering prototype, pilot production to full production environments.

iceMASTER (IM) IN-CIRCUIT EMULATION

The iceMASTER IM-COP8/400 is a full feature, PC based, in-circuit emulation tool developed and marketed by Meta-Link Corporation to support the whole COP8 family of products. National is a resale vendor for these products.

See Figure 19 for configuration.

The iceMASTER IM-COP8/400 with its device specific COP8 Probe provides a rich feature set for developing, testing and maintaining product:

- Real-time in-circuit emulation; full 2.4V-5.5V operation range, full DC-10 MHz clock. Chip options are programmable or jumper selectable.
- Direct connection to application board by package compatible socket or surface mount assembly.
- Full 32 kbytes of loadable programming space that overlays (replaces) the on-chip ROM or EPROM. On-chip RAM and I/O blocks are used directly or recreated on the probe as necessary.
- Full 4k frame synchronous trace memory. Address, instruction, and 8 unspecified, circuit connectable trace lines. Display can be HLL source (e.g., C source), assembly or mixed.

- A full 64k hardware configurable break, trace on, trace off control, and pass count increment events.
- Tool set integrated interactive symbolic debugger—supports both assembler (COFF) and C Compiler (.COD) linked object formats.
- Real time performance profiling analysis; selectable bucket definition.
- Watch windows, content updated automatically at each execution break.
- Instruction by instruction memory/register changes displayed on source window when in single step operation.
- Single base unit and debugger software reconfigurable to support the entire COP8 family; only the probe personality needs to change. Debugger software is processor customized, and reconfigured from a master model file.
- Processor specific symbolic display of registers and bit level assignments, configured from master model file.
- Halt/Idle mode notification.
- On-line HELP customized to specific processor using master model file.
- Includes a copy of COP8-DEV-IBMA assembler and linker SDK.

IM Order Information

Base Unit				
IM-COP8/400-1	iceMASTER Base Unit, 110V Power Supply			
IM-COP8/400-2	iceMASTER Base Unit, 220V Power Supply			
iceMASTER Probe				
MHW-888GG28DWPC	28 DIP			
28 DIP to 28 SO Adapter				
MHW-SOIC28	28 SO			

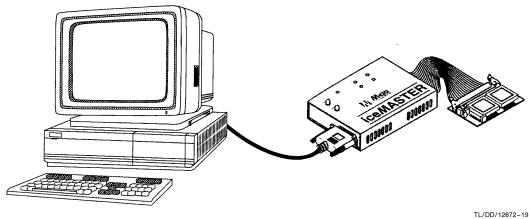


FIGURE 19. COP8 iceMASTER Environment

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iceMASTER DEBUG MODULE (DM)

The iceMASTER Debug Module is a PC based, combination in-circuit emulation tool and COP8 based OTP/EPROM programming tool developed and marketed by MetaLink Corporation to support the whole COP8 family of products. National is a resale vendor for these products.

See Figure 20 for configuration.

The iceMASTER Debug Module is a moderate cost development tool. It has the capability of in-circuit emulation for a specific COP8 microcontroller and in addition serves as a programming tool for COP8 OTP and EPROM product families. Summary of features is as follows:

- Real-time in-circuit emulation; full operating voltage range operation, full DC-10 MHz clock.
- All processor I/O pins can be cabled to an application development board with package compatible cable to socket and surface mount assembly.
- Full 32 kbytes of loadable programming space that overlays (replaces) the on-chip ROM or EPROM. On-chip RAM and I/O blocks are used directly or recreated as necessary.
- 100 frames of synchronous trace memory. The display can be HLL source (C source), assembly or mixed. The most recent history prior to a break is available in the trace memory.
- Configured break points; uses INTR instruction which is modestly intrusive.
- Software—only supported features are selectable.
- Tool set integrated interactive symbolic debugger—supports both assembler (COFF) and C Compiler (.COD)
 SDK linked object formats.

- Instruction by instruction memory/register changes displayed when in single step operation.
- Debugger software is processor customized, and reconfigured from a master model file.
- Processor specific symbolic display of registers and bit level assignments, configured from master model file.
- Halt/Idle mode notification.
- Programming menu supports full product line of programmable OTP and EPROM COP8 products. Program data is taken directly from the overlay RAM.
- Programming of 44 PLCC and 68 PLCC parts requires external programming adapters.
- · Includes wallmount power supply.
- On-board V_{PP} generator from 5V input or connection to external supply supported. Rquires V_{PP} level adjustment per the family programming specification (correct level is provided on an on-screen pop-down display).
- On-line HELP customized to specific processor using master model file.
- Includes a copy of COP8-DEV-IBMA assembler and linker SDK.

DM Order Information

Debug Module Unit			
COP8-DM/888GG			
Cable Adapter			
DM-COP8/28D	28 DIP		
28 DIP to 28 SO Adapter			
DM-COP8/28D-SO	28 SO		

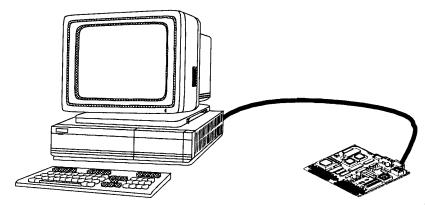


FIGURE 20. COP8-DM Environment

TL/DD/12872-20

iceMASTER EVALUATION PROGRAMMING UNIT (EPU)

The iceMASTER EPU-COP888GG is a PC based, in-circuit simulation tool to support the feature family COP8 products. See *Figure 21* for configuration.

The simulation capability is a very low cost means of evaluating the general COP8 architecture. In addition, the EPU has programming capability, with added adapters, for programming the whole COP8 product family of OTP and EPROM products. The product includes the following features:

- Non-real-time in-circuit simulation. Program overlay memory is PC resident; instructions are downloaded over RS-232 as executed. Approximate performance is 20 kHz.
- Includes a 40-pin DIP cable adapter. Other target packages are not supported. All processor I/O pins are cabled to the application development environment.
- Full 32 kbytes of loadable programming space that overlays (replaces) the on-chip ROM or EPROM. On-chip RAM and I/O blocks are used directly or recreated as necessary.
- On-chip timer and WATCHDOG execution are not well synchronized to the instruction simulation.
- 100 frames of synchronous trace memory. The display can be HLL source (e.g., C source), assembly or mixed.
 The most recent history prior to a break is available in the trace memory.
- Up to eight software configured break points; uses INTR instruction which is modestly intrusive.
- Common look-feel debugger software across all Meta-Link products—only supported features are selectable.

- Tool set integrated interactive symbolic debugger—supports both assembler (COFF) and C Compiler (.COD)
 SDK linked object formats.
- Instruction by instruction memory/register changes displayed when in single step operation.
- Processor specific symbolic display of registers and bit level assignments, configured from master model file.
- Halt/Idle mode notification. Restart requires special handling.
- Programming menu supports full product line of programmable OTP and EPROM COP8 products. Only a 40 ZIF socket is available on the EPU unit. Adapters are available for other part package configurations.
- Integral wall mount power supply provides 5V and develops the required V_{PP} to program parts.
- Includes a copy of COP8-DEV-IBMA assembler, linker SDK.

EPU Order Information

Evaluation Programming Unit				
EPU-COP888GG	Evaluation Programming Unit with debugger and programmer control software with 40 ZIF programming socket.			
General Programming Adapters				
COP8-PGMA-DS44P	28 and 20 DIP and SOIC plus 44 PLCC adapter			

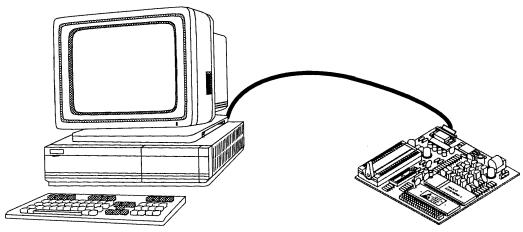


FIGURE 21. EPU-COP8 Tool Environment

TL/DD/12872-21

COP8 ASSEMBLER/LINKER SOFTWARE DEVELOPMENT TOOL KIT

National Semiconductor offers a relocateable COP8 macro cross assembler, linker, librarian and utility software development tool kit. Features are summarized as follows:

- Basic and Feature Family instruction set by "device" type.
- · Nested macro capability.
- Extensive set of assembler directives.
- · Supported on PC/DOS platform.
- · Generates National standard COFF output files.
- Integrated Linker and Librarian.
- Integrated utilities to generate ROM code file outputs.
- DUMPCOFF utility.

This product is integrated as a part of MetaLink tools as a development kit, fully supported by the MetaLink debugger. It may be ordered separately or it is bundled with the MetaLink products at no additional cost.

Order Information

Assembler SDK		
COP8-DEV-IBMA	Assembler SDK on installable 3.5" PC/DOS Floppy Disk Drive format. Periodic upgrades and most recent version is available on National's BBS and Internet.	

COP8 C COMPILER

A C Compiler is developed and marketed by Byte Craft Limited. The COP8C compiler is a fully integrated development tool specifically designed to support the compact embedded configuration of the COP8 family of products.

Features are summarized as follows:

- ANSI C with some restrictions and extensions that optimize development for the COP8 embedded application.
- BITS data type extension. Register declaration #pragma with direct bit level definitions.
- C language support for interrupt routines.
- Expert system, rule based code generation and optimization.
- Performs consistency checks against the architectural definitions of the target COP8 device.
- Generates program memory code.
- Supports linking of compiled object or COP8 assembled object formats.
- · Global optimization of linked code.
- Symbolic debug load format fully sourced level supported by the MetaLink debugger.

INDUSTRY WIDE OTP/EPROM PROGRAMMING SUPPORT

Programming support, in addition to the MetaLink development tools, is provided by a full range of independent approved vendors to meet the needs from the engineering laboratory to full production.

Approved List

Manufacturer	North America	Europe	Asia
BP Microsystems	(800) 225-2102 (713) 688-4600 Fax: (713) 688-0920	+ 49-8152-4183 + 49-8856-932616	+ 852-234-16611 + 852-2710-8121
Data I/O	(800) 426-1045 (206) 881-6444 Fax: (206) 882-1043	+44-0734-440011	Call North America
HI-LO	(510) 623-8860	Call Asia	+886-2-764-0215 Fax: +886-2-756-6403
ICE Technology	(800) 624-8949 (919) 430-7915	+ 44-1226-767404 Fax: 0-1226-370-434	
MetaLink	(800) 638-2423 (602) 926-0797 Fax: (602) 693-0681	+49-80 9156 96-0 Fax: +49-80 9123 86	+852-737-1800
Systems General	(408) 263-6667	+41-1-9450300	+886-2-917-3005 Fax: +886-2-911-1283
Needhams	(916) 924-8037 Fax: (916) 924-8065		

AVAILABLE LITERATURE

For more information, please see the COP8 Basic Family User's Manual, Literature Number 620895, COP8 Feature Family User's Manual, Literature Number 620897 and National's Family of 8-bit Microcontrollers COP8 Selection Guide, Literature Number 630009.

DIAL-A-HELPER SERVICE

Dial-A-Helper is a service provided by the Microcontroller Applications group. The Dial-A-Helper is an Electronic Information System that may be accessed as a Bulletin Board System (BBS) via data modem, as an FTP site on the Internet via standard FTP client application or as an FTP site on the Internet using a standard Internet browser such as Netscape or Mosaic.

The Dial-A-Helper system provides access to an automated information storage and retrieval system . The system capabilities include a MESSAGE SECTION (electronic mail, when accessed as a BBS) for communications to and from the Microcontroller Applications Group and a FILE SECTION which consists of several file areas where valuable application software and utilities could be found.

DIAL-A-HELPER BBS via Standard Modem

Modem: CANADA/U.S.: (800) NSC-MICRO

(800) 672-6427

EUROPE: (+49) 0-8141-351332

Baud: 14.4k

Set-Up: Length: 8-Bit

Parity: None Stop Bit: 1

Operation: 24 Hours, 7 Days

DIAL-A-HELPER via FTP

ftp nscmicro.nsc.com user: anonymous

password: username@yourhost.site.domain

DIAL-A-HELPER via WorldWide Web Browser

ftp://nscmicro.nsc.com

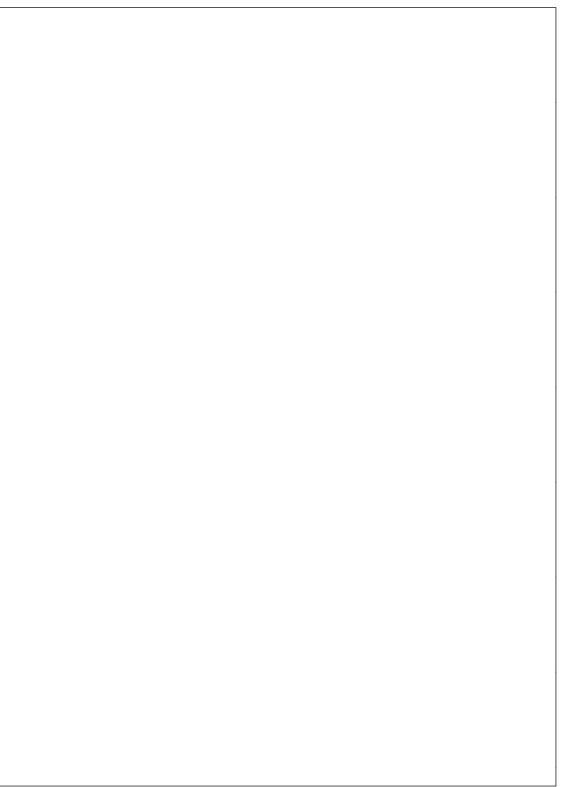
National Semiconductor on the WorldWide Web

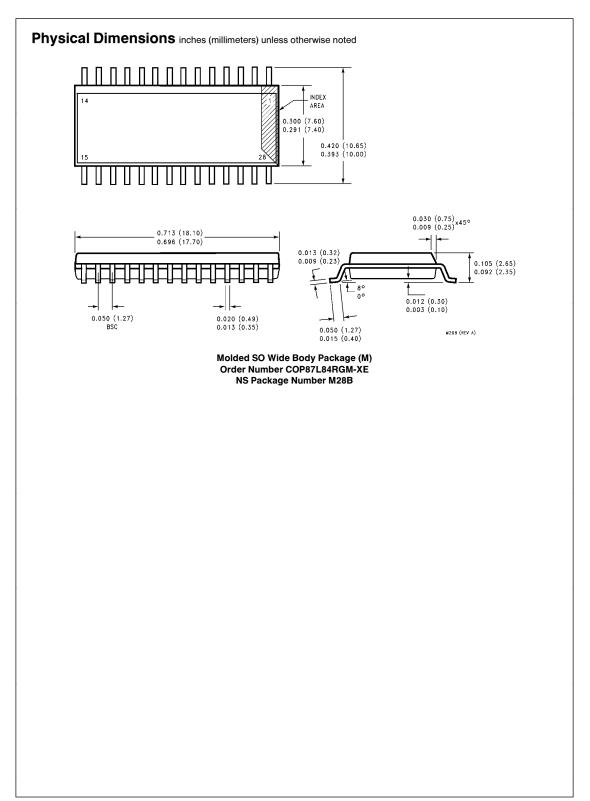
See us on the WorldWide Web at: http://www.natsemi.com

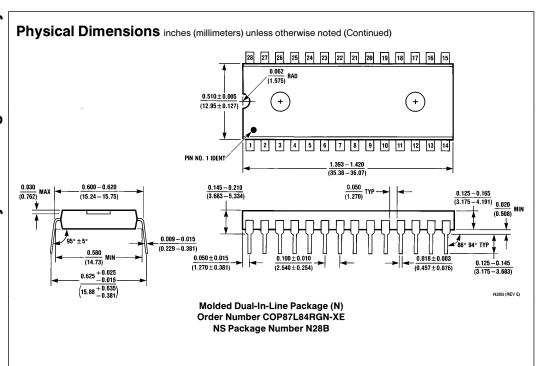
CUSTOMER RESPONSE CENTER

Complete product information and technical support is available from National's customer response centers.

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	Shanghai Tel:	(+86) 21-6415-4092
	Hong Kong Tel:	(+852) 2737-1600
	Korea Tel:	(+82) 2-3771-6909
	Malaysia Tel:	(+60-4) 644-9061
	Singapore Tel:	(+65) 255-2226
	Taiwan Tel:	+886-2-521-3288
AUSTRALIA:	Tel:	(+61) 3-9558-9999
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- 2. A critical component is any component of a life support device or system whose failure to perform can be reasonably expected to cause the failure of the life support device or system, or to affect its safety or effectiveness.



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