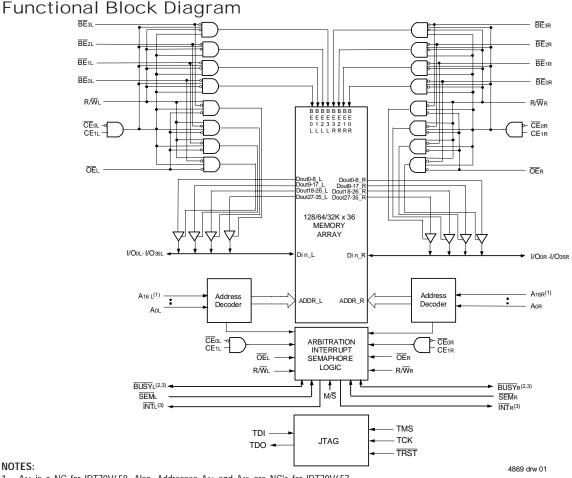
HIGH-SPEED 3.3V 128/64/32K x 36 **ASYNCHRONOUS DUAL-PORT** STATIC RAM

IDT70V659/58/57S

Features

- True Dual-Port memory cells which allow simultaneous access of the same memory location
- High-speed access
 - Commercial: 10/12/15ns (max.)
 - Industrial: 12/15ns (max.)
- Dual chip enables allow for depth expansion without external logic
- IDT70V659/58/57 easily expands data bus width to 72 bits or more using the Master/Slave select when cascading more than one device
- M/S = VIH for BUSY output flag on Master, $M/\overline{S} = VIL$ for \overline{BUSY} input on Slave
- **Busy and Interrupt Flags**
- On-chip port arbitration logic

- ٠ Full on-chip hardware support of semaphore signaling between ports
- ٠ Fully asynchronous operation from either port
- ٠ Separate byte controls for multiplexed bus and bus matching compatibility
- Supports JTAG features compliant to IEEE 1149.1
- LVTTL-compatible, single 3.3V (±150mV) power supply for core
- ٠ LVTTL-compatible, selectable 3.3V (±150mV)/2.5V (±100mV) power supply for I/Os and control signals on each port
- Available in a 208-pin Plastic Quad Flatpack, 208-ball fine pitch Ball Grid Array, and 256-ball Ball Grid Array
- Industrial temperature range (-40°C to +85°C) is available for selected speeds
- Green parts available, see ordering information



NOTES:

- 1 A16 is a NC for IDT70V658. Also, Addresses A16 and A15 are NC's for IDT70V657.
- BUSY is an input as a Slave (M/S=VIL) and an output when it is a Master (M/S=VIH). 2.
- BUSY and INT are non-tri-state totem-pole outputs (push-pull). 3.

OCTOBER 2008

High-Speed 3.3V 128/64/32K x 36 Asynchronous Dual-Port Static RAM

Description

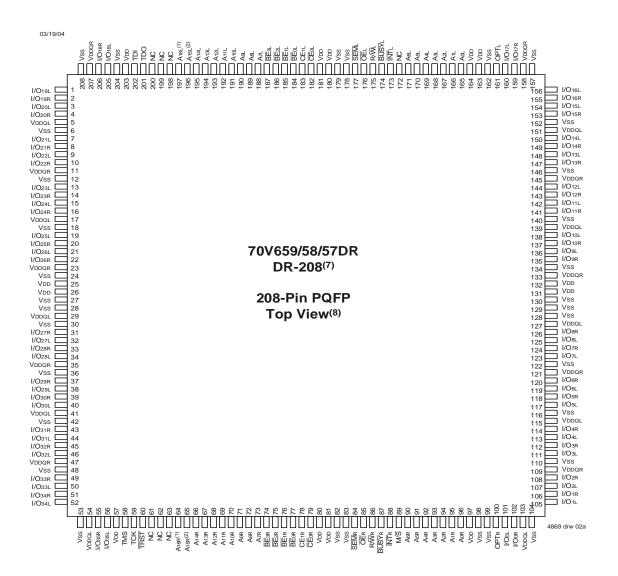
The IDT70V659/58/57 is a high-speed 128/64/32K x 36 Asynchronous Dual-Port Static RAM. The IDT70V659/58/57 is designed to be used as a stand-alone 4/2/1Mbit Dual-Port RAM or as a combination MASTER/ SLAVE Dual-Port RAM for 72-bit-or-more word system. Using the IDT MASTER/SLAVE Dual-Port RAM approach in 72-bit or wider memory system applications results in full-speed, error-free operation without the need for additional discrete logic.

This device provides two independent ports with separate control,

address, and I/O pins that permit independent, asynchronous access for reads or writes to any location in memory. An automatic power down feature controlled by the chip enables (either \overline{CE} 0 or CE1) permit the on-chip circuitry of each port to enter a very low standby power mode.

The 70V659/58/57 can support an operating voltage of either 3.3V or 2.5V on one or both ports, controlled by the OPT pins. The power supply for the core of the device (VDD) remains at 3.3V.

PinConfigurations^(3,4,5,6,7,8)



- 1. Pin is a NC for IDT70V658 and IDT70V657.
- 2. Pin is a NC for IDT70V657.
- 3. All VDD pins must be connected to 3.3V power supply.
- 4. All VDDD pins must be connected to appropriate power supply: 3.3V if OPT pin for that port is set to VDD (3.3V) and 2.5V if OPT pin for that port is set to VSS (0V).
- 5. All Vss pins must be connected to ground.
- 6. Package body is approximately 28mm x 28mm x 3.5mm.
- 7. This package code is used to reference the package diagram.
- 8. This text does not indicate orientation of the actual part-marking.

Pin Configurations^(3,4,5,6,7,8) (con't.)

70V659**/**58/5**7**BC BC-256⁽⁷⁾

256-Pin BGA Top View⁽⁸⁾

03/19/04															
A1	^{A2}	A3	A4	A5	A6	A7	A8	A9	A10	a11	A12	A13	A14	A15	A16
NC	TDI	NC	NC	A14L	A11L	A8L	BE2L	CE1L	OEL	ĪNT∟	A5L	A2L	Aol	NC	NC
b1	^{B2}	^{B3}	^{B4}	B5	B6	B7	B8	B9	^{B10}	B11	B12	B13	^{B14}	в15	B16
I/O18L	NC	TDO	NC	A15L ⁽²⁾	A12L	A9L	BE3L	CEOL	R/WL	NC	A4L	A1L	NC	I/O17L	NC
C1	C2	C3	C4	C5	C6	C7	C8	C9	C10	C11	C12	C13	C14	C15	C16
I/O18R	I/O19L	Vss	A16L ⁽¹⁾	A13L	A10L	A7L	BE1L	BEOL	SEM∟	BUSYL	A6L	A3L	OPTL	I/O17R	I/O16L
D1	d2	d3	d4	d5	d6	d7	d8	d9	d10	d11	d12	D13	D14	d15	d16
I/O20R	I/O19r	I/O20L	Vdd	Vddql	Vddql	Vddqr	Vddqr	Vddql	Vddql	Vddqr	Vddqr	Vdd	I/O15R	I/O15L	I/O16R
e1	e2	e3	e4	e5	e6	e7	^{E8}	^{E9}	E10	e11	e12	e13	e14	e15	e16
I/O21r	I/O21l	I/O22L	Vddql	Vdd	Vdd	Vss	Vss	Vss	Vss	Vdd	Vdd	Vddqr	I/O13L	I/O14L	I/O14r
f1	f2	f3	f4	f5	^{F6}	F7	^{F8}	^{F9}	^{F10}	F11	f12	f13	F14	f15	F16
I/O23L	I/O22R	I/O23R	Vddql	Vdd	Vss	Vss	Vss	Vss	Vss	Vss	Vdd	Vddqr	I/O12R	I/O13R	I/O12L
G1	G2	G3	g4	G5	_{G6}	G7	G8	^{G9}	G10	G11	G12	g13	G14	G15	G16
I/O24R	I/O24L	I/O25L	Vddqr	Vss	Vss	Vss	Vss	Vss	Vss	Vss	Vss	Vddql	I/O10L	I/O11L	I/O11R
h1	h2	H3	h4	H5	H6	нт	H8	H9	H10	H11	^{H12}	h13	h14		h16
I/O26l	I/O25r	I/O26R	Vddqr	Vss	Vss	Vss	Vss	Vss	Vss	Vss	Vss	Vddql	I/O9r		I/O10r
J1	j2	j3	j4	J5	_{J6}	J7	_{J8}	^{J9}	^{J10}	J11	^{J12}	j13	J14	j15	J16
I∕O27L	I/O28R	I/O27R	Vddql	Vss	Vss	Vss	Vss	Vss	Vss	Vss	Vss	Vddqr	I/O8r	I/O7r	I/O8∟
k1	k2	k3	k4	^{K5}	K6	кт	ка	к9	K10	к11	к12	k13	k14	к15	к16
I/O29r	I/O29l	I/O28l	Vddql	Vss	Vss	Vss	Vss	Vss	Vss	Vss	Vss	Vddqr	I/O6r	I/O6L	I/O7l
l1	l2	l3	l4	l5	L6	L7	L8	L9	L10	L11	l12	l13	l14	l15	l16
I/O30l	I/O31R	I/O30r	Vddqr	Vdd	Vss	Vss	Vss	Vss	Vss	Vss	Vdd	Vddql	I/O5l	I/O4r	I/O5r
m1	m2	m3	^{m4}	^{M5}	M6	M7	M8	M9	M10	M11	M12	m13	м14	m15	^{M16}
I/O32r	I/O32l	I/O31l	Vddqr	Vdd	Vdd	Vss	Vss	Vss	Vss	Vdd	Vdd	Vddql	I/Oзr	I/O3l	I/O4L
n1	n2	n3	N4	n5	n6	n7	n8	n9	n10	n11	n12	N13	n14	n15	n16
I/O33l	I/O34r	I/O33r	Vdd	Vddqr	Vddqr	Vddql	Vddql	Vddqr	Vddqr	Vddql	Vddql	Vdd	I/O2l	I/O1r	I/O2r
Р1	p2	P3	Р4	P5	P6	P7	P8	P9	^{P10}	^{p11}	P12	Р13	P14	P15	P16
I/O35R	I/O34L	TMS	А16R ⁽¹⁾	A13R	A10R	A7r	BE1R	BEOR	SEMr	BUSYr	A6R	Азк	I/Ool	I/O0R	I/O1L
r1	R2	^{R3}	^{R4}	R5	-	r7	r8	R9	^{R10}	R11	R12	R13	^{R14}	R15	R16
I/O35l	NC	TRST	NC	A15R ⁽²⁾		A9r	BE3r	CE0R	R/WR	M/S	A4R	A1R	OPTr	NC	NC
T1	T2	тз	T4	t5	t6	t7	t8	t9	T10	t11	t12	t13	t14	T15	^{т16}
NC	TCK	NC	NC	A14r	A11r	A8r	BE2r	CE1r	OEr	INTr	A5r	A2r	Aor	NC	NC

03/19/04

4869 drw 02c

- 1. Pin is a NC for IDT70V658 and IDT70V657.
- 2. Pin is a NC for IDT70V657.
- 3. All VDD pins must be connected to 3.3V power supply.
- 4. All VDDD pins must be connected to appropriate power supply: 3.3V if OPT pin for that port is set to VDD (3.3V), and 2.5V if OPT pin for that port is set to Vss (0V).
- 5. All Vss pins must be connected to ground supply.
- 6. Package body is approximately 17mm x 17mm x 1.4mm, with 1.0mm ball-pitch.
- 7. This package code is used to reference the package diagram.
- 8. This text does not indicate orientation of the actual part-marking.

Pin Configuration^(3,4,5,6,7,8) (con't.)

	03/19/0 1	⁰⁴	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	
А	I/O19L	I/O18L	Vss	TDO	NC	A16L ⁽¹⁾	A12L	AsL	BE 1L	VDD	SEML	ĪNTL	A4L	Aol	OPT∟	I/O17L	Vss	А
В	I/O20R	Vss	I/O18R	TDI	NC	A13L	A9L	BE2L	CE OL	Vss	BUSYL	A5∟	Aıl	Vss	Vddqr	I/O16L	I/O15R	В
С	Vddql	I/O19R	Vddqr	Vdd	NC	A14L	A10L	BE 3L	CE1L	Vss	R/₩L	A6L	A2L	VDD	I/O16R	I/O15L	Vss	С
D	I/O22L	Vss	I/O21L	I/O20L	A15L ⁽²⁾	A11L	A7L	BEOL	Vdd	ŌĒL	NC	Аз∟	Vdd	I/O17R	Vddql	I/O14L	I/O14R	D
Е	I/O23L	I/O22R	Vddqr	I/O21R										I/O _{12L}	I/O13R	Vss	I/O _{13L}	Е
F	Vddql	I/O23R	I/O24L	Vss										Vss	I/O12R	I/O11L	Vddqr	F
G	I/O26L	Vss	I/O25L	I/O24R										I/O9L	Vddql	I/O10L	I/O11r	G
Н	Vdd	I/O26R	Vddqr	I/O25R					59/58 F-20		3F			Vdd	I/O9R	Vss	I/O10R	Н
J	VDDQL	Vdd	Vss	Vss				208	-Ball	BGA	٨			Vss	Vdd	Vss	Vddqr	J
Κ	I/O28R	Vss	I/O27r	Vss				Τc	p Vie	ew ⁽⁸⁾				I/O7R	VDDQL	I/08R	Vss	К
L	I/O29R	VO28L	Vddqr	I/O27L										I/O6R	I/O7∟	Vss	I/O _{8L}	L
М	VDDQL	I/O29L	I/O30R	Vss										Vss	I/O _{6L}	I/O5r	Vddqr	Μ
Ν	I/O31L	Vss	I/O31r	I/O30L										I/O3r	Vddql	I/O4R	I/O5L	Ν
Ρ	I/O32R	I/O32L	Vddqr	I/O35R	TRST	A16R ⁽¹⁾	A12R	Asr	₿Ē1r	Vdd	SEMR	ĪNTr	A4R	I/O2L	I/O3L	Vss	I/O4L	Ρ
R	Vss	I/O33L	I/O34 R	тск	NC	A13R	A9R	BE2R	CEor	Vss	BUSYR	A5R	A1R	Vss	VDDQL	I/O1R	VDDQR	R
Т	I/O33R	I/O34L	Vddql	тмз	NC	A14R	A10R	БЕзк	CE1R	Vss	R/WR	Agr	A2R	Vss	I/Oor	Vss	I/O2R	Т
U	Vss	I/O35L	Vdd	NC	A15R ⁽²⁾	A11R	A7R	BEOR	VDD	ŌĒr	M/S	Aзr	Aor	Vdd	OPTr	I/Ool	I/O1L	U
																4869	drw 02b	

- 1. Pin is a NC for IDT70V658 and IDT70V657.
- 2. Pin is a NC for IDT70V657.
- 3. All VDD pins must be connected to 3.3V power supply.
- 4. All VDDO pins must be connected to appropriate power supply: 3.3V if OPT pin for that port is set to VDD (3.3V) and 2.5V if OPT pin for that port is set to Vss (OV).
- 5. All Vss pins must be connected to ground.
- 6. Package body is approximately 15mm x 15mm x 1.4mm with 0.8mm ball pitch.
- 7. This package code is used to reference the package diagram.
- 8. This text does not indicate orientation of the actual part-marking.

High-Speed 3.3V 128/64/32K x 36 Asynchronous Dual-Port Static RAM

Industrial and Commercial Temperature Ranges

Pin Names

Left Port	Right Port	Names
CEOL, CE1L	$\overline{\text{CE}}$ OR, CE1R	Chip Enables - (Input)
R/WL	R/Wr	Read/Write Enable - (Input)
ŌĒL	ŌĒr	Output Enable - (Input)
Aol - A16l ⁽³⁾	Aor - A16r ⁽³⁾	Address - (Input)
1/Ool - 1/O35L	I/Oor - I/O35r	Data Input/Output
SEML	SEM R	Semaphore Enable - (Input)
ĪNTL	INT R	Interrupt Flag - (Output)
BUSYL	BUSYR	Busy Flag - (Output) ⁽⁴⁾
BEOL - BE3L	BEOR - BE3R	Byte Enables (9-bit bytes) - (Input)
VDDQL	Vddqr	Power (I/O Bus) (3.3V or 2.5V) - (Input) ⁽¹⁾
OPTL	OPTr	Option for selecting VDDax - (Input) ^(1,2)
N	NS	Master or Slave Select - (Input)
V	DD	Power (3.3V) - (Input) ⁽¹⁾
V	SS	Ground (0V) - (Input)
Т	DI	Test Data Input
П	DO	Test Data Output
Ţ	СК	Test Logic Clock (10MHz)
יד	MS	Test Mode Select
ना	RST	Reset (Initialize TAP Controller)

4869 tbl 01

- 1. VDD, OPTx, and VDDox must be set to appropriate operating levels prior to applying inputs on I/Ox.
- 2. OPTx selects the operating voltage levels for the I/Os and controls on that port. If OPTx is set to VIH (3.3V), then that port's I/Os and controls will operate at 3.3V levels and VDDox must be supplied at 3.3V. If OPTx is set to VIL (0V), then that port's I/Os and controls will operate at 2.5V levels and VDDox must be supplied at 2.5V. The OPT pins are independent of one another—both ports can operate at 3.3V levels, both can operate at 2.5V levels, or either can operate at 3.3V with the other at 2.5V.
- Addresses A16x is a NC for IDT70V658. Also, Addresses A16x and A15x are NC's for IDT70V657.
- 4. $\overline{\text{BUSY}}$ is an input as a slave (M/ $\overline{\text{S}}$ = VIL).

	0V659 -Speed			/32K x	36 As	ynchro	onous	Dual-P	ort Static	RAM	Inc	lustrial an	d Commercial Temperature Ranges
Tru	th T	abl	e I–	-Re	ad/\	Nrit	eai	nd E	nable	e Cont	rol ^(1,2))	
ŌĒ	SEM	Ē€	CE1	BE 3	BE 2	BE 1	BE 0	R/W	Byte 3 I/O27-35	Byte 2 I/O18-26	Byte 1 I/O9-17	Byte 0 I/Oo-8	MODE
Х	Н	Н	Х	Х	Х	Х	Х	Х	High-Z	High-Z	High-Z	High-Z	Deselected-Power Down
Х	Н	Х	L	Х	Х	Х	Х	Х	High-Z	High-Z	High-Z	High-Z	Deselected–Power Down
Х	Н	L	Н	Н	Н	Н	Н	Х	High-Z	High-Z	High-Z	High-Z	All Bytes Deselected
Х	Н	L	Н	Н	Н	Н	L	L	High-Z	High-Z	High-Z	Din	Write to Byte 0 Only
Х	Н	L	Н	Н	Н	L	Н	L	High-Z	High-Z	Din	High-Z	Write to Byte 1 Only
Х	Н	L	Н	Н	L	Н	Н	L	High-Z	Din	High-Z	High-Z	Write to Byte 2 Only
Х	Н	L	Н	L	Н	Н	Н	L	Din	High-Z	High-Z	High-Z	Write to Byte 3 Only
Х	Н	L	Н	Н	Н	L	L	L	High-Z	High-Z	Din	Din	Write to Lower 2 Bytes Only
Х	Н	L	Н	L	L	Н	Н	L	Din	Din	High-Z	High-Z	Write to Upper 2 bytes Only
Х	Н	L	Н	L	L	L	L	L	Din	Din	Din	Din	Write to All Bytes
L	Н	L	Н	Н	Н	Н	L	Н	High-Z	High-Z	High-Z	Dout	Read Byte 0 Only
L	Н	L	Н	Н	Н	L	Н	Н	High-Z	High-Z	Dout	High-Z	Read Byte 1 Only
L	Н	L	Н	Н	L	Н	Н	Н	High-Z	Dout	High-Z	High-Z	Read Byte 2 Only
L	Н	L	Н	L	Н	Н	Н	Н	Dout	High-Z	High-Z	High-Z	Read Byte 3 Only
L	Н	L	Н	Н	Н	L	L	Н	High-Z	High-Z	Dout	Dout	Read Lower 2 Bytes Only
L	Н	L	Н	L	L	Н	Н	Н	Dout	Dout	High-Z	High-Z	Read Upper 2 Bytes Only
L	Н	L	Н	L	L	L	L	Н	Dout	Dout	Dout	Dout	Read All Bytes
Н	Η	L	Н	L	L	L	L	Х	High-Z	High-Z	High-Z	High-Z	Outputs Disabled

NOTES:

1. "H" = VIH, "L" = VIL, "X" = Don't Care.

2. It is possible to read or write any combination of bytes during a given access. A few representative samples have been illustrated here.

Truth Table II - Semaphore Read/Write Control⁽¹⁾

			Inp	uts ⁽¹⁾				Out	puts	
CE ⁽²⁾	R/W OE BE3 BE2 BE1 BE5				BĒ₀	SEM	I/O 1-35	I/Oo	Mode	
Н	Н	L	L	L	L	L L DATAOUT [Read Data in Semaphore Flag ⁽³⁾
Н	Ŷ	Х	Х	Х	Х	L	L	Х	DATAIN	Write I/Oo into Semaphore Flag
Ĺ	Х	Х	Х	Х	Х	Х	L			Not Allowed

NOTES:

1. There are eight semaphore flags written to I/Oo and read from all the I/Os (I/Oo-I/O35). These eight semaphore flags are addressed by Ao-A2.

2. $\overline{CE} = L$ occurs when $\overline{CE}_0 = V_{IL}$ and $CE_1 = V_{IH}$.

3. Each byte is controlled by the respective \overline{BEn} . To read data $\overline{BEn} = VIL$.

4869 tbl 03

4869 tbl 02

IDT70V659/58/57S High-Speed 3.3V 128/64/32K x 36 Asynchronous Dual-Port Static RAM

Industrial and Commercial Temperature Ranges

4869 tbl 07

Recommended DC Operating Conditions with VDDQ at 2.5V

Symbol	Parameter	Min.	Тур.	Max.	Unit
Vdd	Core Supply Voltage	3.15	3.3	3.45	V
VDDQ	I/O Supply Voltage ⁽³⁾	2.4	2.5	2.6	V
Vss	Ground	0	0	0	V
Vін	Input High Voltage ⁽³⁾ (Address & Control Inputs)	1.7	-	Vddq + 100mV ⁽²⁾	V
Vih	Input High Voltage - I/O ⁽³⁾	1.7	-	VDDQ + 100mV ⁽²⁾	V
VIL	Input Low Voltage	-0.5 ⁽¹⁾	_	0.7	V
				4	869 tbl 06

NOTES:

1. VIL \geq -1.5V for pulse width less than 10 ns.

2. VTERM must not exceed VDDQ + 100mV.

To select operation at 2.5V levels on the I/Os and controls of a given port, the OPT pin for that port must be set to Vss (0V), and VDDOX for that port must be supplied as indicated above.

Capacitance⁽¹⁾

(TA = +25°C, F = 1.0MHz) PQFP ONLY

Symbol	Parameter	Conditions	Max.	Unit
Cin	Input Capacitance	Vin = 0V	8	рF
Cout ⁽²⁾	Output Capacitance	Vout = 0V	10.5	pF
				4869 tbl 08

NOTES:

 These parameters are determined by device characterization, but are not production tested.

2. COUT also references CI/O.

Absolute Maximum Ratings⁽¹⁾

Symbol	Rating	Commercial & Industrial	Unit	
Vterm ⁽²⁾ (Vdd)	Vod Terminal Voltage with Respect to GND	-0.5 to + 4.6	V	N
TBIAS ⁽³⁾	Temperature Under Bias	-55 to +125	٥C	1
Tstg	Storage Temperature	-65 to +150	٥C	
Тли	Junction Temperature	+150	٥C	
IOUT(For VDDQ = 3.3V)	DC Output Current	50	mA	2
IOUT(For VDDQ = 2.5V)	DC Output Current	40	mA	3

4869 tbl 05

Recommended DC Operating Conditions with VDDQ at 3.3V

Symbol	Parameter	Min.	Тур.	Мах.	Unit
Vdd	Core Supply Voltage	3.15	3.3	3.45	۷
VDDQ	I/O Supply Voltage ⁽³⁾	3.15	3.3	3.45	۷
Vss	Ground	0	0	0	۷
Vih	Input High Voltage (Address & Control Inputs) ⁽³⁾	2.0		Vddq + 150mV ⁽²⁾	V
Vн	Input High Voltage - I/O ⁽³⁾	2.0	_	Vddq + 150mV ⁽²⁾	۷
VIL	Input Low Voltage	-0.3(1)	_	0.8	۷

NOTES:

1. VIL \geq -1.5V for pulse width less than 10 ns.

2. VTERM must not exceed VDDQ + 150mV.

 To select operation at 3.3V levels on the I/Os and controls of a given port, the OPT pin for that port must be set to Vob (3.3V), and Vobox for that port must be supplied as indicated above.

Maximum Operating Temperature and Supply Voltage⁽¹⁾

Grade	Ambient Temperature	GND	Vdd
Commercial	0°C to +70°C	0V	3.3V <u>+</u> 150mV
Industrial	-40°C to +85°C	0V	3.3V <u>+</u> 150mV
			4869 tbl 04

NOTE:

1. This is the parameter TA. This is the "instant on" case temperature.

- Stresses greater than those listed under ABSOLUTE MAXIMUM RATINGS may cause permanent damage to the device. This is a stress rating only and functional operation of the device at these or any other conditions above those indicated in the operational sections of this specification is not implied. Exposure to absolute maximum rating conditions for extended periods may affect reliability.
- VTERM must not exceed VDD + 150mV for more than 25% of the cycle time or 4ns maximum, and is limited to ≤ 20mA for the period of VTERM ≥ VDD + 150mV.
- Ambient Temperature under DC Bias. No AC Conditions. Chip Deselected.

High-Speed 3.3V 128/64/32K x 36 Asynchronous Dual-Port Static RAM

Industrial and Commercial Temperature Ranges

4869 tbl 09

DC Electrical Characteristics Over the Operating Temperature and Supply Voltage Range (VDD = 3.3V ± 150mV)

			70V659	/58/57S	
Symbol	Parameter	Test Conditions	Min.	Max.	Unit
Lu	Input Leakage Current ⁽¹⁾	$V_{DDQ} = Max., V_{IN} = 0V$ to V_{DDQ}	_	10	μA
llo	Output Leakage Current	\overline{CE}_0 = Vih or CE1 = Vil, Vout = 0V to VDDQ	_	10	μA
Vol (3.3V)	Output Low Voltage ⁽²⁾	IOL = +4mA, $VDDQ = Min$.		0.4	V
Voн (3.3V)	Output High Voltage ⁽²⁾	Ioh = -4mA, VDDQ = Min.	2.4		V
Vol (2.5V)	Output Low Voltage ⁽²⁾	IOL = +2mA, $VDDQ = Min$.		0.4	V
Voн (2.5V)	Output High Voltage ⁽²⁾	Ioh = -2mA, VDDQ = Min.	2.0		V

NOTE:

1. At VDD \leq - 2.0V input leakages are undefined.

2. VDDQ is selectable (3.3V/2.5V) via OPT pins. Refer to p.6 for details.

DC Electrical Characteristics Over the Operating Temperature and Supply Voltage Range⁽³⁾ (V_{DD} = $3.3V \pm 150mV$)

						58/57S10 Only	Co	58/57S12 m'l Ind	Co	58/57S15 m'l Ind	
ymbol	Parameter	Test Condition	Versio	Version		Typ. ⁽⁴⁾ Max.		Мах.	Тур. ⁽⁴⁾	Мах.	Uni
IDD	Dynamic Operating Current (Both	CEL and CER= VIL, Outputs Disabled	COM'L	S	340	500	315	465	300	440	m
	Ports Active)	$f = fMAX^{(1)}$	IND	S			365	515	350	490	
ISB1	Standby Current (Both Ports - TTL	$\overline{CE}L = \overline{CE}R = VIH$ f = fMAX ⁽¹⁾	COM'L	S	115	165	90	125	75	100	m
	Level Inputs)	I = IMAX''	IND	S			115	150	100	125	
ISB2	Standby Current (One Port - TTL	$\overline{CE}^{"}A^{"} = VIL \text{ and } \overline{CE}^{"}B^{"} = VIH^{(5)}$	COM'L	S	225	340	200	325	175	315	m
	Level Inputs)	Active Port Outputs Disabled, f=fMAX ⁽¹⁾	IND	S			225	365	200	350	
ISB3	Full Standby Current (Both Ports - CMOS	Both Ports CEL and CER ≥ VDDQ - 0.2V,	COM'L	S	3	15	3	15	3	15	m
	Level Inputs)	$\label{eq:VIN_solution} \begin{array}{l} \text{VIN} \geq \text{VDDQ} \text{ - } 0.2\text{V} \text{ or } \text{VIN} \leq 0.2\text{V}, \\ f = 0^{(2)} \end{array}$	IND	S	—		6	15	6	15	
ISB4	Full Standby Current (One Port - CMOS	$\overline{CE}^{"}A^{"} \leq 0.2V$ and $\overline{CE}^{"}B^{"} \geq VDDQ - 0.2V^{(5)}$	COM'L	S	220	335	195	320	170	310	m
	evel Inputs) $ \begin{array}{l} \mbox{VIN} \geq \mbox{VDDQ} \ - \ 0.2 \mbox{V or VIN} \leq \ 0.2 \mbox{V}, \\ \mbox{Active Port, Outputs Disabled,} \\ \mbox{f} \ = \ \mbox{fMAX}^{(1)} \end{array} $		S	_		220	360	195	345		

NOTES:

1. At f = fMAX, address and control lines (except Output Enable) are cycling at the maximum frequency read cycle of 1/tRc, using "AC TEST CONDITIONS" at input levels of GND to 3V.

2. f = 0 means no address or control lines change. Applies only to input at CMOS level standby.

3. Port "A" may be either left or right port. Port "B" is the opposite from port "A".

4. VDD = 3.3V, TA = 25°C for Typ, and are not production tested. IDD DC(f=0) = 120mA (Typ).

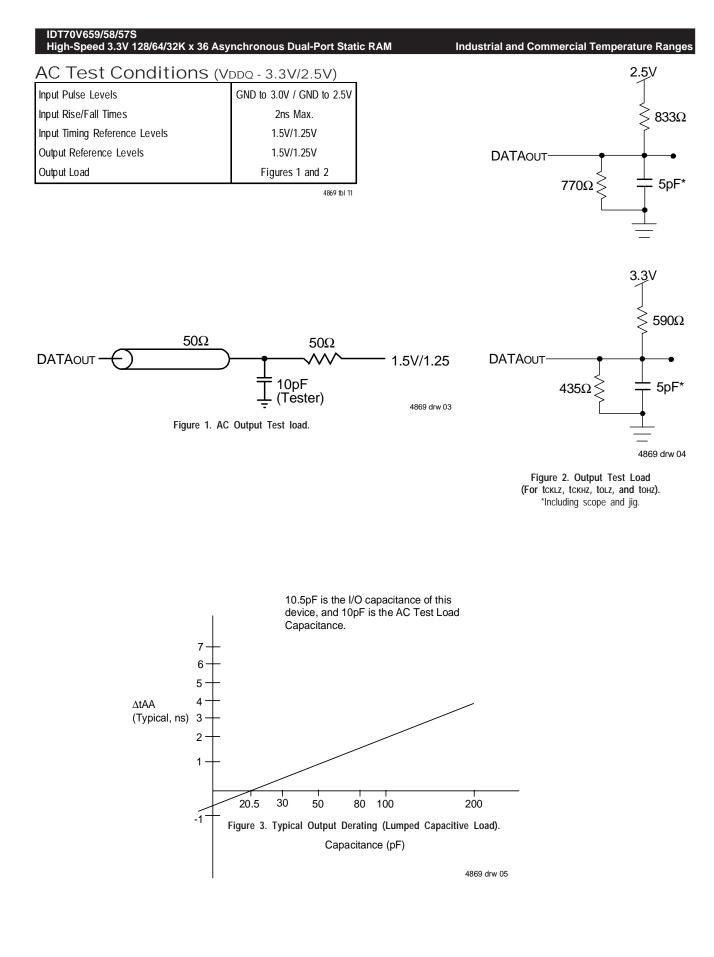
5. $\overline{CE}x = VIL$ means $\overline{CE}0x = VIL$ and CE1x = VIH

 $\overline{CE}x = VIH$ means $\overline{CE}ox = VIH$ or CE1x = VIL

 $\overline{\text{CE}}x \leq 0.2V$ means $\overline{\text{CE}}\text{ox} \leq 0.2V$ and $\text{CE}\text{1x} \geq V\text{DDQ}$ - 0.2V

 $\overline{\text{CE}}\text{x} \geq \text{V}\text{DDQ}$ - 0.2V means $\overline{\text{CE}}\text{ox} \geq \text{V}\text{DDQ}$ - 0.2V or CE1x - 0.2V

"X" represents "L" for left port or "R" for right port.



IDT70V659/58/57S High-Speed 3.3V 128/64/32K x 36 Asynchronous Dual-Port Static RAM

AC Electrical Characteristics Over the

Operating Temperature and Supply Voltage Range⁽⁵⁾

			70V659/58/57S10 Com'l Only		70V659/58/57S12 Com'l & Ind		70V659/58/57S15 Com'l & Ind	
Symbol	Parameter	Min.	Max.	Min.	Max.	Min.	Max.	Unit
READ CYCLE	•							
trc	Read Cycle Time	10		12		15		ns
taa	Address Access Time	-	10	—	12		15	ns
t ACE	Chip Enable Access Time ⁽³⁾	-	10	-	12		15	ns
t ABE	Byte Enable Access Time ⁽³⁾	-	5	_	6	-	7	ns
taoe	Output Enable Access Time	-	5	—	6	_	7	ns
toн	Output Hold from Address Change	3		3		3	_	ns
١LZ	Output Low-Z Time ^(1,2)	0		0		0	_	ns
tHZ	Output High-Z Time ^(1,2)	0	4	0	6	0	8	ns
tPU	Chip Enable to Power Up Time ⁽²⁾	0		0		0	_	ns
tpd	Chip Disable to Power Down Time ⁽²⁾	_	10	—	10		15	ns
tsop	Semaphore Flag Update Pulse (OE or SEM)		4		6	_	8	ns
t SAA	Semaphore Address Access Time	3	10	3	12	3	20	ns

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AC Electrical Characteristics Over the Operating Temperature and Supply Voltage⁽⁵⁾

		70V659/58/57 Com'l Onl		Co	58/57S12 m'l Ind	70V659/58/57S15 Com'l & Ind		
Symbol	Parameter	Min.	Мах.	Min.	Мах.	Min.	Мах.	Unit
WRITE CYCLE								
twc	Write Cycle Time	10		12		15	_	ns
tew	Chip Enable to End-of-Write ⁽³⁾	8		10		12	_	ns
taw	Address Valid to End-of-Write	8		10		12	_	ns
tas	Address Set-up Time ⁽³⁾	0		0		0	_	ns
twp	Write Pulse Width	8	-	10		12	_	ns
twr	Write Recovery Time	0		0	_	0	_	ns
tow	Data Valid to End-of-Write	6		8		10	_	ns
tDH	Data Hold Time ⁽⁴⁾	0	_	0		0	_	ns
twz	Write Enable to Output in High-Z ^(1,2)	-	4		4	-	4	ns
tow	Output Active from End-of-Write ^(1,2,4)	0		0		0		ns
tswrd	SEM Flag Write to Read Time	5		5		5	_	ns
tsps	SEM Flag Contention Window	5		5		5		ns

NOTES:

1. Transition is measured 0mV from Low or High-impedance voltage with Output Test Load (Figure 2).

2. This parameter is guaranted by device characterization, but is not production tested. 3. To access RAM, \overline{CE} = VIL and \overline{SEM} = VIH. To access semaphore, \overline{CE} = VIH and \overline{SEM} = VIL. Either condition must be valid for the entire tew time.

4. The specification for tDH must be met by the device supplying write data to the RAM under all operating conditions. Although tDH and tow values will vary over voltage and temperature, the actual tDH will always be smaller than the actual tow.

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5. These values are valid regardless of the power supply level selected for I/O and control signals (3.3V/2.5V). See page 6 for details.

High-Speed 3.3V 128/64/32K x 36 Asynchronous Dual-Port Static RAM

Industrial and Commercial Temperature Ranges

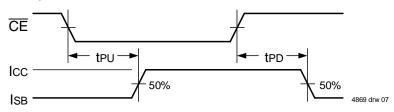
tRC ADDR taa⁽⁴⁾ $\mathsf{tACE}^{\,(4)}$ CE taoe⁽⁴⁾ ŌĒ tabe⁽⁴⁾ BEn R/W tон tLZ (1) VALID DATA⁽⁴⁾ DATAOUT tHZ⁽²⁾ BUSYOUT tbd^(3,4) 4869 drw 06

Waveform of Read Cycles⁽⁵⁾

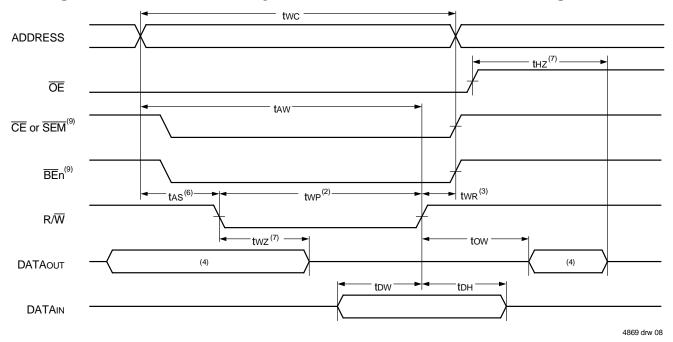
NOTES:

- 1. Timing depends on which signal is asserted last, $\overline{\text{OE}}$, $\overline{\text{CE}}$ or $\overline{\text{BEn}}$.
- 2. Timing depends on which signal is de-asserted first \overline{CE} , \overline{OE} or \overline{BEn} .
- 3. tbpD delay is required only in cases where the opposite port is completing a write operation to the same address location. For simultaneous read operations BUSY has no relation to valid output data.
- 4. Start of valid data depends on which timing becomes effective last tAOE, tACE, tAA or tBDD.
- 5. $\overline{\text{SEM}} = \text{VIH}.$

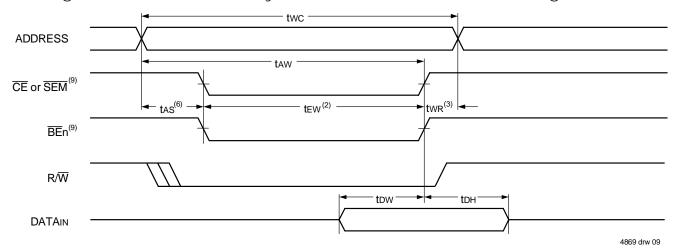
Timing of Power-Up Power-Down



Timing Waveform of Write Cycle No. 1, R/W Controlled Timing^(1,5,8)



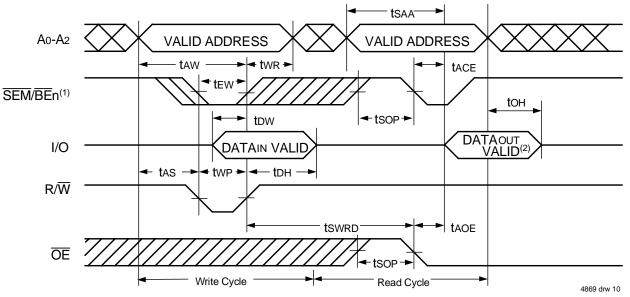
Timing Waveform of Write Cycle No. 2, CE Controlled Timing^(1,5)



- 1. $R\overline{W}$ or \overline{CE} or $\overline{BEn} = V_{H}$ during all address transitions.
- 2. A write occurs during the overlap (tew or twp) of a CE = VIL and a R/W = VIL for memory array writing cycle.
- 3. two is measured from the earlier of CE or R/W (or SEM or R/W) going HIGH to the end of write cycle.
- 4. During this period, the I/O pins are in the output state and input signals must not be applied.
- 5. If the CE or SEM = VIL transition occurs simultaneously with or after the R/W = VIL transition, the outputs remain in the High-impedance state.
- 6. Timing depends on which enable signal is asserted last, \overline{CE} or R/\overline{W} .
- 7. This parameter is guaranteed by device characterization, but is not production tested. Transition is measured 0mV from steady state with the Output Test Load (Figure 2).
- If \overline{OE} = ViL during R/W controlled write cycle, the write pulse width must be the larger of twp or (twz + tow) to allow the I/O drivers to turn off and data to be 8. placed on the bus for the required tow. If \overrightarrow{OE} = VIH during an R/W controlled write cycle, this requirement does not apply and the write pulse can be as short as the specified twp
- 9. To access RAM, CE = VIL and SEM = VIH. To access semaphore, CE = VIH and SEM = VIL. tew must be met for either condition.

Industrial and Commercial Temperature Ranges

Timing Waveform of Semaphore Read after Write Timing, Either Side⁽¹⁾

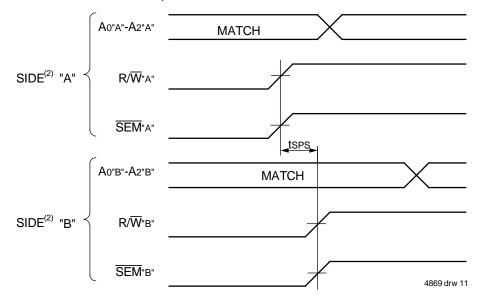


NOTES:

1. $\overline{CE} = V_{IH}$ for the duration of the above timing (both write and read cycle) (Refer to Chip Enable Truth Table). Refer also to Truth Table II for appropriate \overline{BE} controls.

2. "DATAOUT VALID" represents all I/O's (I/O0 - I/O35) equal to the semaphore value.

Timing Waveform of Semaphore Write Contention^(1,3,4)



- 1. DOR = DOL = VIL, $\overline{CE}L = \overline{CE}R = VIH$. Refer to Truth Table II for appropriate \overline{BE} controls.
- 2. All timing is the same for left and right ports. Port "A" may be either left or right port. "B" is the opposite from port "A".
- 3. This parameter is measured from R/W*A* or SEM*A* going HIGH to R/W*B* or SEM*B* going HIGH.
- 4. If tsps is not satisfied, the semaphore will fall positively to one side or the other, but there is no guarantee which side will be granted the semaphore flag.

IDT70V659/58/57S High-Speed 3.3V 128/64/32K x 36 Asynchronous Dual-Port Static RAM

AC Electrical Characteristics Over the Operating Temperature and Supply Voltage Range

Cumhal	Duran dar		70V659/58/57S10 Com'l Only		70V659/58/57S12 Com'l & Ind		70V659/58/57S15 Com'l & Ind	
Symbol	Parameter	Min.	Max.	Min.	Max.	Min.	Max.	Unit
BUSY TIMING	э G (M/S=Viн)		•	•	•		•	
t BAA	BUSY Access Time from Address Match		10		12		15	ns
tBDA	BUSY Disable Time from Address Not Matched		10		12		15	ns
t BAC	BUSY Access Time from Chip Enable Low		10		12		15	ns
tBDC	BUSY Disable Time from Chip Enable High		10		12		15	ns
taps	Arbitration Priority Set-up Time ⁽²⁾	5	—	5		5		ns
tBDD	BUSY Disable to Valid Data ⁽³⁾		10		12		15	ns
twн	Write Hold After BUSY ⁽⁵⁾	8	—	10		12		ns
	G (M/S=VIL)	•						
twв	BUSY Input to Write ⁽⁴⁾	0	—	0		0		ns
twн	Write Hold After BUSY ⁽⁵⁾	8	—	10		12		ns
PORT-TO-POF	RT DELAY TIMING							
twdd	Write Pulse to Data Delay ⁽¹⁾		22		25		30	ns
todd	Write Data Valid to Read Data Delay ⁽¹⁾		20		22		25	ns
		-	-	-	-	-	-	4869 tbl 14

NOTES:

1. Port-to-port delay through RAM cells from writing port to reading port, refer to "Timing Waveform of Write with Port-to-Port Read and BUSY (M/S = VI-)".

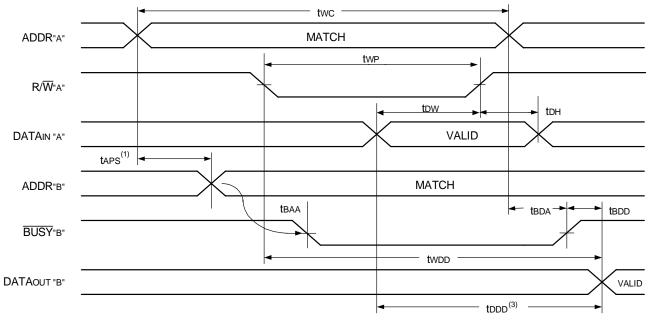
2. To ensure that the earlier of the two ports wins.

3. tBDD is a calculated parameter and is the greater of the Max. spec, twDD - twp (actual), or tDDD - tDw (actual).

4. To ensure that the write cycle is inhibited on port "B" during contention on port "A".

5. To ensure that a write cycle is completed on port "B" after contention on port "A".

Timing Waveform of Write with Port-to-Port Read and **BUSY** $(M/S = VIH)^{(2,4,5)}$

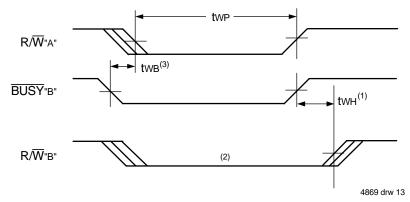


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NOTES:

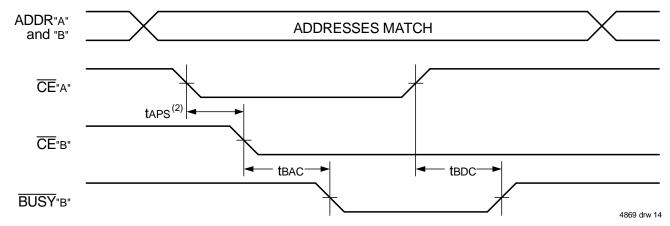
- 1. To ensure that the earlier of the two ports wins. taps is ignored for $M/\overline{S} = VIL$ (SLAVE).
- $2. \quad \overline{CE}_L = \overline{CE}_R = VIL.$
- 3. $\overline{OE} = V_{IL}$ for the reading port.
- 4. If M/S = VIL (slave), BUSY is an input. Then for this example BUSY A" = VIH and BUSY B" input is shown above.
- 5. All timing is the same for left and right ports. Port "A" may be either the left or right port. Port "B" is the port opposite from port "A".

Timing Waveform of Write with **BUSY** (M/S = VIL)

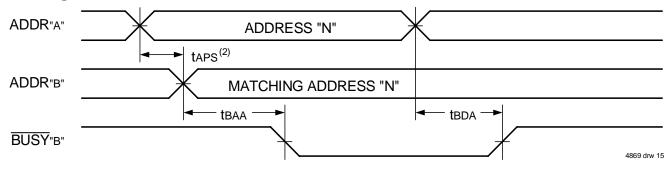


- 1. two must be met for both $\overline{\text{BUSY}}$ input (SLAVE) and output (MASTER).
- 2. BUSY is asserted on port "B" blocking R/W"B", until BUSY"B" goes HIGH.
- 3. twb is only for the 'slave' version.

Waveform of **BUSY** Arbitration Controlled by \overline{CE} Timing (M/ \overline{S} = VIH)⁽¹⁾



Waveform of **BUSY** Arbitration Cycle Controlled by Address Match Timing $(M/\mathbf{S} = VIH)^{(1)}$



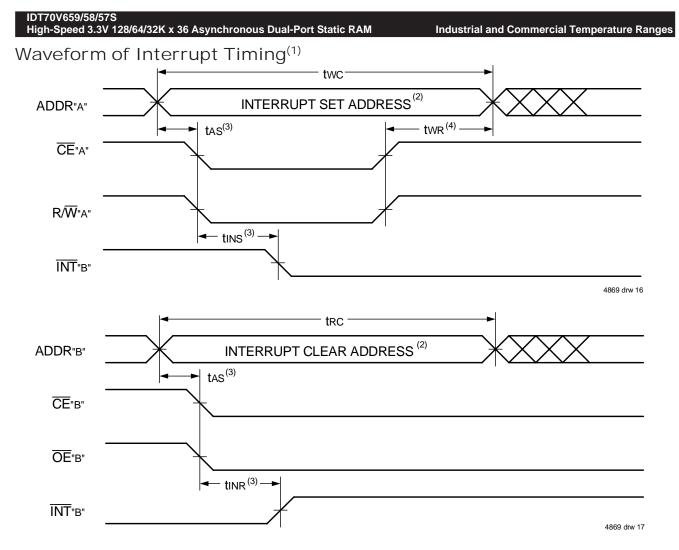
NOTES:

1. All timing is the same for left and right ports. Port "A" may be either the left or right port. Port "B" is the port opposite from port "A".

2. If taps is not satisfied, the BUSY signal will be asserted on one side or another but there is no guarantee on which side BUSY will be asserted.

AC Electrical Characteristics Over the Operating Temperature and Supply Voltage Range

			58/57S10 'I Only	Co	58/57S12 m'l Ind	Co	58/57S15 m'l Ind	
Symbol	Parameter	Min.	Мах.	Min.	Мах.	Min.	Max.	Unit
INTERRUPT T	IMING							
tas	Address Set-up Time	0		0		0	-	ns
twr	Write Recovery Time	0		0	_	0	_	ns
tins	Interrupt Set Time		10		12	_	15	ns
tinr	Interrupt Reset Time		10		12	_	15	ns



NOTES:

1. All timing is the same for left and right ports. Port "A" may be either the left or right port. Port "B" is the port opposite from port "A".

2. Refer to Interrupt Truth Table.

3. Timing depends on which enable signal (\overline{CE} or R/\overline{W}) is asserted last. 4. Timing depends on which enable signal (\overline{CE} or R/\overline{W}) is de-asserted first.

Left Port							Right Por	t		
R/₩L	Ē ⊾	ŌE∟	A16L-A0L ^(5,6)	ĪNTL	R/WR	ĈĒ	OE R	A16R-A0R ^(5,6)	ĪNTR	Function
L	L	Х	1FFFF	Х	Х	Х	Х	Х	L ⁽²⁾	Set Right INTR Flag
Х	Х	Х	Х	Х	Х	L	L	1FFFF	H ⁽³⁾	Reset Right INTR Flag
Х	Х	Х	Х	L ⁽³⁾	L	L	Х	1FFFE	Х	Set Left INTL Flag
Х	L	L	1FFFE	H ⁽²⁾	Х	Х	Х	Х	Х	Reset Left INTL Flag

Truth Table III — Interrupt Flag^(1,4)

NOTES:

1. Assumes $\overline{\text{BUSY}}_{L} = \overline{\text{BUSY}}_{R} = VIH$.

2. If $\overline{\text{BUSY}}_{L} = \text{VIL}$, then no change.

3. If $\overline{\text{BUSY}}_{R} = \text{VIL}$, then no change.

4. INTL and INTR must be initialized at power-up.

A16x is a NC for IDT70V658, therefore Interrupt Addresses are FFFF and FFFE. 5.

6. A16x and A15x are NC's for IDT70V657, therefore Interrupt Addresses are 7FFF and 7FFE. 4869 tbl 16

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High-Speed 3.3V 128/64/32K x 36 Asynchronous Dual-Port Static RAM

Industrial and Commercial Temperature Ranges

Truth Table IV — Address **BUSY** Arbitration

	In	puts	Out	puts	
		Aol-A16l ⁽⁴⁾ Aor-A16r	BUSYL ⁽¹⁾	BUSY R ⁽¹⁾	Function
Х	Х	NO MATCH	Н	Н	Normal
Н	Х	MATCH	Н	Н	Normal
Х	Н	MATCH	Н	Н	Normal
L	L	MATCH	(2)	(2)	Write Inhibit ⁽³⁾

NOTES:

4869 tbl 17

1. Pins BUSYL and BUSYR are both outputs when the part is configured as a master. Both are inputs when configured as a slave. BUSY outputs on the IDT70V659/58/57 are push-pull, not open drain outputs. On slaves the BUSY input internally inhibits writes.

2. "L" if the inputs to the opposite port were stable prior to the address and enable inputs of this port. "H" if the inputs to the opposite port became stable after the address and enable inputs of this port. If taps is not met, either BUSYL or BUSYR = LOW will result. BUSYL and BUSYR outputs can not be LOW simultaneously.

Writes to the left port are internally ignored when BUSYL outputs are driving LOW regardless of actual logic level on the pin. Writes to the right port are internally ignored when BUSYR outputs are driving LOW regardless of actual logic level on the pin.

A16x is a NC for IDT70V658, therefore Address comparison will be for A0 - A15. Also, A16x and A15x are NC's for IDT70V657, therefore Address comparison will be for A0 - A14.

Truth Table V — Example of Semaphore Procurement Sequence^(1,2,3)

Functions	Do - D35 Left	Do - D35 Right	Status
No Action	1	1	Semaphore free
Left Port Writes "0" to Semaphore	0	1	Left port has semaphore token
Right Port Writes "0" to Semaphore	0	1	No change. Right side has no write access to semaphore
Left Port Writes "1" to Semaphore	1	0	Right port obtains semaphore token
Left Port Writes "0" to Semaphore	1	0	No change. Left port has no write access to semaphore
Right Port Writes "1" to Semaphore	0	1	Left port obtains semaphore token
Left Port Writes "1" to Semaphore	1	1	Semaphore free
Right Port Writes "0" to Semaphore	1	0	Right port has semaphore token
Right Port Writes "1" to Semaphore	1	1	Semaphore free
Left Port Writes "0" to Semaphore	0	1	Left port has semaphore token
Left Port Writes "1" to Semaphore	1	1	Semaphore free

NOTES:

1. This table denotes a sequence of events for only one of the eight semaphores on the IDT70V659/58/57.

2. There are eight semaphore flags written to via I/Oo and read from all I/O's (I/Oo-I/O35). These eight semaphores are addressed by Ao - A2.

3. TE = VIH, SEM = VIL to access the semaphores. Refer to the Semaphore Read/Write Control Truth Table.

Functional Description

The IDT70V659/58/57 provides two ports with separate control, address and I/O pins that permit independent access for reads or writes to any location in memory. The IDT70V659/58/57 has an automatic power down feature controlled by \overline{CE} . The $\overline{CE}o$ and CE1 control the on-chip power down circuitry that permits the respective port to go into a standby mode when not selected ($\overline{CE} = HIGH$). When a port is enabled, access to the entire memory array is permitted.

Interrupts

If the user chooses the interrupt function, a memory location (mail box or message center) is assigned to each port. The left port interrupt flag (\overline{INTL}) is asserted when the right port writes to memory location 1FFFE (HEX) (FFFE for IDT70V658 and 7FFE for IDT70V657), where a write is defined as $\overline{CER} = R/\overline{WR} = VIL$ per the Truth Table III. The left port clears the interrupt through access of address location 1FFFE (FFFE for IDT70V658 and 7FFE for IDT70V657) when $\overline{CEL} = \overline{OEL} = VIL$, R/\overline{W} is

High-Speed 3.3V 128/64/32K x 36 Asynchronous Dual-Port Static RAM

a "don't care". Likewise, the right port interrupt flag (INTR) is asserted when the left port writes to memory location 1FFFF (HEX) (FFFF for IDT70V658 and 7FFF for IDT70V657) and to clear the interrupt flag (INTR), the right port must read the memory location 1FFFF (FFFF for IDT70V658 and 7FFF for IDT70V657). The message (36 bits) at 1FFFE (FFFE for IDT70V658 and 7FFE for IDT70V657) or 1FFFF (FFFF for IDT70V658 and 7FFF for IDT70V657) is user-defined since it is an addressable SRAM location. If the interrupt function is not used, address locations 1FFFE (FFFE for IDT70V658 and 7FFF for IDT70V657) and 1FFFF (FFFF for IDT70V658 and 7FFF for IDT70V657) are not used as mail boxes, but as part of the random access memory. Refer to Truth Table III for the interrupt operation.

Busy Logic

Busy Logic provides a hardware indication that both ports of the RAM have accessed the same location at the same time. It also allows one of the two accesses to proceed and signals the other side that the RAM is "Busy". The $\overline{\text{BUSY}}$ pin can then be used to stall the access until the operation on the other side is completed. If a write operation has been attempted from the side that receives a $\overline{\text{BUSY}}$ indication, the write signal is gated internally to prevent the write from proceeding.

The use of $\overline{\text{BUSY}}$ logic is not required or desirable for all applications. In some cases it may be useful to logically OR the $\overline{\text{BUSY}}$ outputs together and use any $\overline{\text{BUSY}}$ indication as an interrupt source to flag the event of an illegal or illogical operation. If the write inhibit function of $\overline{\text{BUSY}}$ logic is not desirable, the $\overline{\text{BUSY}}$ logic can be disabled by placing the part in slave mode with the M/\overline{S} pin. Once in slave mode the $\overline{\text{BUSY}}$ pin operates solely as a write inhibit input pin. Normal operation can be programmed by tying the $\overline{\text{BUSY}}$ pins HIGH. If desired, unintended write operations can be prevented to a port by tying the $\overline{\text{BUSY}}$ pin for that port LOW.

The BUSY outputs on the IDT70V659/58/57 RAM in master mode, are push-pull type outputs and do not require pull up resistors to operate. If these RAMs are being expanded in depth, then the BUSY indication for the resulting array requires the use of an external AND gate.

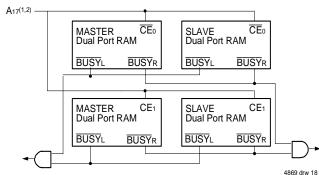


Figure 3. Busy and chip enable routing for both width and depth expansion with IDT70V659/58/57 RAMs.

NOTES:

1. A16 for IDT70V658.

2. A15 for IDT70V657.

Width Expansion with Busy Logic Master/Slave Arrays

When expanding an IDT70V659/58/57 RAM array in width while using BUSY logic, one master part is used to decide which side of the RAMs array will receive a BUSY indication, and to output that indication. Any

number of slaves to be addressed in the same address range as the master use the \overline{BUSY} signal as a write inhibit signal. Thus on the IDT70V659/58/ 57 RAM the \overline{BUSY} pin is an output if the part is used as a master (M/ \overline{S} pin = VIH), and the \overline{BUSY} pin is an input if the part used as a slave (M/ \overline{S} pin = VIL) as shown in Figure 3.

If two or more master parts were used when expanding in width, a split decision could result with one master indicating $\overline{\text{BUSY}}$ on one side of the array and another master indicating $\overline{\text{BUSY}}$ on one other side of the array. This would inhibit the write operations from one port for part of a word and inhibit the write operations from the other port for the other part of the word.

The BUSY arbitration on a master is based on the chip enable and address signals only. It ignores whether an access is a read or write. In a master/slave array, both address and chip enable must be valid long enough for a BUSY flag to be output from the master before the actual write pulse can be initiated with the R/W signal. Failure to observe this timing can result in a glitched internal write inhibit signal and corrupted data in the slave.

Semaphores

The IDT70V659/58/57 is an extremely fast Dual-Port 128/64/32K x 36 CMOS Static RAM with an additional 8 address locations dedicated to binary semaphore flags. These flags allow either processor on the left or right side of the Dual-Port RAM to claim a privilege over the other processor for functions defined by the system designer's software. As an example, the semaphore can be used by one processor to inhibit the other from accessing a portion of the Dual-Port RAM or any other shared resource.

The Dual-Port RAM features a fast access time, with both ports being completely independent of each other. This means that the activity on the left port in no way slows the access time of the right port. Both ports are identical in function to standard CMOS Static RAM and can be read from or written to at the same time with the only possible conflict arising from the simultaneous writing of, or a simultaneous READ/WRITE of, a non-semaphore location. Semaphores are protected against such ambiguous situations and may be used by the system program to avoid any conflicts in the non-semaphore portion of the Dual-Port RAM. These devices have an automatic power-down feature controlled by \overline{CE} , the Dual-Port RAM enable, and \overline{SEM} , the semaphore enable. The \overline{CE} and \overline{SEM} pins control on-chip power down circuitry that permits the respective port to go into standby mode when not selected.

Systems which can best use the IDT70V659/58/57 contain multiple processors or controllers and are typically very high-speed systems which are software controlled or software intensive. These systems can benefit from a performance increase offered by the IDT70V659/58/ 57s hardware semaphores, which provide a lockout mechanism without requiring complex programming.

Software handshaking between processors offers the maximum in system flexibility by permitting shared resources to be allocated in varying configurations. The IDT70V659/58/57 does not use its semaphore flags to control any resources through hardware, thus allowing the system designer total flexibility in system architecture.

An advantage of using semaphores rather than the more common methods of hardware arbitration is that wait states are never incurred in either processor. This can prove to be a major advantage in very highspeed systems.

High-Speed 3.3V 128/64/32K x 36 Asynchronous Dual-Port Static RAM

Industrial and Commercial Temperature Ranges

How the Semaphore Flags Work

The semaphore logic is a set of eight latches which are independent of the Dual-Port RAM. These latches can be used to pass a flag, or token, from one port to the other to indicate that a shared resource is in use. The semaphores provide a hardware assist for a use assignment method called "Token Passing Allocation." In this method, the state of a semaphore latch is used as a token indicating that a shared resource is in use. If the left processor wants to use this resource, it requests the token by setting the latch. This processor then verifies its success in setting the latch by reading it. If it was successful, it proceeds to assume control over the shared resource. If it was not successful in setting the latch, it determines that the right side processor has set the latch first, has the token and is using the shared resource. The left processor can then either repeatedly request that semaphore's status or remove its request for that semaphore to perform another task and occasionally attempt again to gain control of the token via the set and test sequence. Once the right side has relinquished the token, the left side should succeed in gaining control.

The semaphore flags are active LOW. A token is requested by writing a zero into a semaphore latch and is released when the same side writes a one to that latch.

The eight semaphore flags reside within the IDT70V659/58/57 in a separate memory space from the Dual-Port RAM. This address space is accessed by placing a low input on the SEM pin (which acts as a chip select for the semaphore flags) and using the other control pins (Address, \overline{CE} , R/W and \overline{BE} o) as they would be used in accessing a standard Static RAM. Each of the flags has a unique address which can be accessed by either side through address pins A0-A2. When accessing the semaphores, none of the other address pins has any effect.

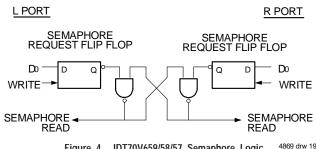
When writing to a semaphore, only data pin Do is used. If a low level is written into an unused semaphore location, that flag will be set to a zero on that side and a one on the other side (see Truth Table V). That semaphore can now only be modified by the side showing the zero. When a one is written into the same location from the same side, the flag will be set to a one for both sides (unless a semaphore request from the other side is pending) and then can be written to by both sides. The fact that the side which is able to write a zero into a semaphore subsequently locks out writes from the other side is what makes semaphore flags useful in interprocessor communications. (A thorough discussion on the use of this feature follows shortly.) A zero written into the same location from the other side will be stored in the semaphore request latch for that side until the semaphore is freed by the first side.

When a semaphore flag is read, its value is spread into all data bits so that a flag that is a one reads as a one in all data bits and a flag containing a zero reads as all zeros. The read value is latched into one side's output register when that side's semaphore select (SEM, BEn) and output enable (OE) signals go active. This serves to disallow the semaphore from changing state in the middle of a read cycle due to a write cycle from the other side. Because of this latch, a repeated read of a semaphore in a test loop must cause either signal (SEM or OE) to go inactive or the output will never change. However, during reads BEn functions only as an output for semaphore. It does not have any influence on the semaphore control logic.

A sequence WRITE/READ must be used by the semaphore in order to guarantee that no system level contention will occur. A processor requests access to shared resources by attempting to write a zero into a semaphore location. If the semaphore is already in use, the semaphore

request latch will contain a zero, yet the semaphore flag will appear as one, a fact which the processor will verify by the subsequent read (see Table V). As an example, assume a processor writes a zero to the left port at a free semaphore location. On used instead, system contention problems could have occurred during the gap between the read and write cycles.

It is important to note that a failed semaphore request must be followed by either repeated reads or by writing a one into the same location. The reason for this is easily understood by looking at the simple logic diagram of the semaphore flag in Figure 4. Two semaphore request latches feed into a semaphore flag. Whichever latch is first to present a zero to the semaphore flag will force its side of the semaphore flag LOW and the other side HIGH. This condition will continue until a one is written to the same semaphore request latch. Should the other side's semaphore request latch have been written to a zero in the meantime, the semaphore flag will flip over to the other side as soon as a one is written into the first side's request latch. The second side's flag will now stay LOW until its semaphore request latch is written to a one. From this it is easy to understand that, if a semaphore is requested and the processor which requested it no longer needs the resource, the entire system can hang up until a one is written into that semaphore request latch.





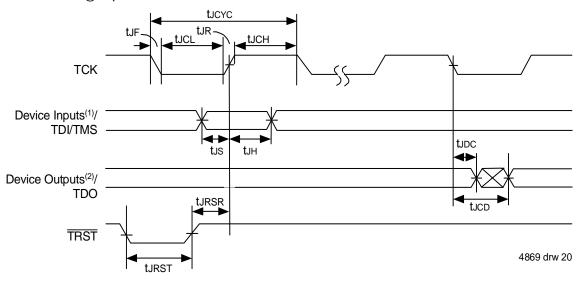
The critical case of semaphore timing is when both sides request a single token by attempting to write a zero into it at the same time. The semaphore logic is specially designed to resolve this problem. If simultaneous requests are made, the logic guarantees that only one side receives the token. If one side is earlier than the other in making the request, the first side to make the request will receive the token. If both requests arrive at the same time, the assignment will be arbitrarily made to one port or the other.

One caution that should be noted when using semaphores is that semaphores alone do not guarantee that access to a resource is secure. As with any powerful programming technique, if semaphores are misused or misinterpreted, a software error can easily happen.

Initialization of the semaphores is not automatic and must be handled via the initialization program at power-up. Since any semaphore request flag which contains a zero must be reset to a one, all semaphores on both sides should have a one written into them at initialization from both sides to assure that they will be free when needed.

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JTAG Timing Specifications



NOTES:

- 1. Device inputs = All device inputs except TDI, TMS, and TRST.
- 2. Device outputs = All device outputs except TDO.

Symbol	Parameter	Min.	Мах.	Units
tucyc	JTAG Clock Input Period	100		ns
tлсн	JTAG Clock HIGH	40		ns
ticl	JTAG Clock Low	40		ns
tır	JTAG Clock Rise Time	_	3(1)	ns
tJF	JTAG Clock Fall Time		3(1)	ns
U RST	JTAG Reset	50		ns
tursr	JTAG Reset Recovery	50		ns
tucd	JTAG Data Output		25	ns
tudc	JTAG Data Output Hold	0		ns
tıs	JTAG Setup	15		ns
tн	JTAG Hold	15		ns

JTAG AC Electrical Characteristics^(1,2,3,4)

NOTES:

- 1. Guaranteed by design.
- 2. 30pF loading on external output signals.
- 3. Refer to AC Electrical Test Conditions stated earlier in this document.
- 4. JTAG operations occur at one speed (10MHz). The base device may run at any speed specified in this datasheet.

IDT70V659/58/57S High-Speed 3.3V 128/64/32K x 36 Asynchronous Dual-Port Static RAM

Industrial and Commercial Temperature Ranges

Identification Register Definitions

Instruction Field	Value	Description
Revision Number (31:28)	0x0	Reserved for version number
IDT Device ID (27:12)	0x303 ⁽¹⁾	Defines IDT part number
IDT JEDEC ID (11:1)	0x33	Allows unique identification of device vendor as IDT
ID Register Indicator Bit (Bit 0)	1	Indicates the presence of an ID register

NOTE:

1. Device ID for IDT70V658 is 0x30B. Device ID for IDT70V657 is 0x323.

Scan Register Sizes

Register Name	Bit Size
Instruction (IR)	4
Bypass (BYR)	1
Identification (IDR)	32
Boundary Scan (BSR)	Note (3)

4869 tbl 21

System Interface Parameters

Instruction	Code	Description
EXTEST	0000	Forces contents of the boundary scan cells onto the device outputs ⁽¹⁾ . Places the boundary scan register (BSR) between TDI and TDO.
BYPASS	1111	Places the bypass register (BYR) between TDI and TDO.
IDCODE	0010	Loads the ID register (IDR) with the vendor ID code and places the register between TDI and TDO.
HIGHZ	0100	Places the bypass register (BYR) between TDI and TDO. Forces all device output drivers to a High-Z state.
CLAMP	0011	Uses BYR. Forces contents of the boundary scan cells onto the device outputs. Places the bypass register (BYR) between TDI and TDO.
SAMPLE/PRELOAD	0001	Places the boundary scan register (BSR) between TDI and TDO. SAMPLE allows data from device inputs ⁽²⁾ and outputs ⁽¹⁾ to be captured in the boundary scan cells and shifted serially through TDO. PRELOAD allows data to be input serially into the boundary scan cells via the TDI.
RESERVED	All other codes	Several combinations are reserved. Do not use codes other than those identified above.

NOTES:

1. Device outputs = All device outputs except TDO.

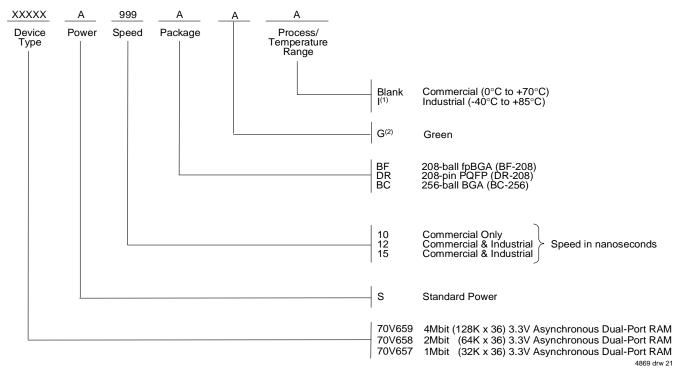
2. Device inputs = All device inputs except TDI, TMS, and TRST.

3. The Boundary Scan Descriptive Language (BSDL) file for this device is available on the IDT website (www.idt.com), or by contacting your local IDT sales representative.

4869 tbl 20

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Ordering Information



Notes:

1. Contact your local sales office for Industrial temp range in other speeds, packages and powers.

2. Green parts available. For specific speeds, packages and powers contact your local sales office.

Datasheet Document History:

6/2/00:	Initial Public Offering
8/11/00:	Page 6, 13 & 20 Inserted additional BEn information
6/20/01:	Page 14 Increased BUSY TIMING parameters tBDA, tBAC, tBDC and tBDD for all speeds
	Page 21 Changed maximum value for JTAG AC Electrical Characteristics for tuch from 20ns to 25ns
12/17/01:	Page 2, 3 & 4 Added date revision for pin configurations
	Page 8, 10, 14 & 16 Removed I-temp 15ns speed from DC & AC Electrical Characteristics
	Page 23 Removed I-temp 15ns speed from ordering information
	Added I-temp footnote
	Page 1 & 23 Replaced ™ logo with ® logo
03/19/04:	Consolidated multiple devices into one data sheet
	Removed "Preliminary" Status
03/22/05:	Page 1 Added green availability to features
	Page 24 Added green indicator to ordering information
	Page 1 & 24 Replaced old IDT TM with new IDT TM logo
07/25/08:	Page 9 Corrected a typo in the DC Chars table
10/23/08:	Page 24 Removed "IDT" from orderable part number



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